

File System Benchmarks Then, Now, and Tomorrow

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Presented by

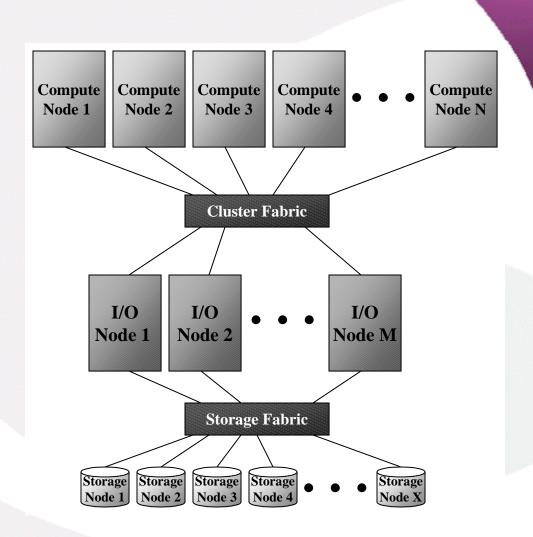
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Overview

- If measuring the performance of I/O subsystems was not complicated enough, it is further complicated by SANs and Clusters
- SANs and emerging clustering technologies add a distributed aspect to the file systems themselves
- As the cluster/SAN grows in size, so does the task of performance measurement
- The objective of this study is to identify some of the more significant issues involved with file system benchmarking in a highly scalable clustered environment
- This research is based on work being done at Los Alamos National Labs on the ASCI 30T machine

: The ASCI 30TeraOp Machine

- ~300 Compute nodes
- ~64 I/O Nodes
- 32 processors per node
- 8 Cluster Fabric connections per compute/I/O node
- 32 FC connections per I/O node into Storage Fabric – ~2048 2.5Gb FC connections into SAN
- ~700 TB disk storage



: Hardware Infrastructure Issues

- Compute nodes How many, how large, and how many connections into the cluster fabric
- System Area Network or Cluster Fabric Latency, bandwidth, and overhead per link
- <u>I/O Nodes</u> How many, how large, and how many connections into the Storage Area Network fabric
- Storage Area Network or Storage Fabric Latency, bandwidth, and overhead per link
- Storage Devices How many, latency, bandwidth, command overhead, number of connections into the SAN per "device"

: System level issues

- The number of measurement points has increased from one computer system to many computer systems
- All the computer systems <u>share</u> access to the disk subsystem, or more importantly the data
- Sharing occurs at many levels
 - File data
 - Metadata (I.e. directories)
 - Host bus adapters
 - Switches
 - Disk controllers
 - Disk media
- Important to separate the performance of the underlying hardware from the file system software

: Other effects

- Caching Effects
 - Distributed File System data and metadata caching
 - Local file system caching
 - Device data caching
 - Caching policies
 - Read versus write
 - Temporal (LRU, ...etc)
 - Data size (l.e. don't cache large files)
- File System Aging effects
 - Fragmentation effects on performance
 - Monitoring and defragmentation impact on performance

Benchmarking versus Characterization

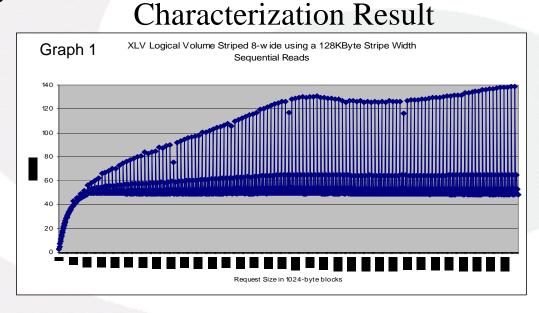
- Benchmarking generally yields a limited set of values that represent the performance of a file system under a specific set of operational parameters
- Characterization provides detailed graphs that describe the performance of a file system under a continuum of operational parameters

Benchmark Result

120 MB/sec

Or

400 I/O Ops per sec



Benchmark I/O Permutations

- Given N Compute Nodes and M I/O nodes there are three permutations of concurrent (parallel) access at the extremes
 - 1 to M a single process accessing M compute nodes
 - N to 1 N process threads (across N compute nodes) accessing a single I/O node (file)
 - N to M N compute nodes accessing M I/O nodes
- Measure the performance of each of the file system levels in isolation (if possible) and then as a composite of all file system levels
- Report results from different <u>perspectives</u>

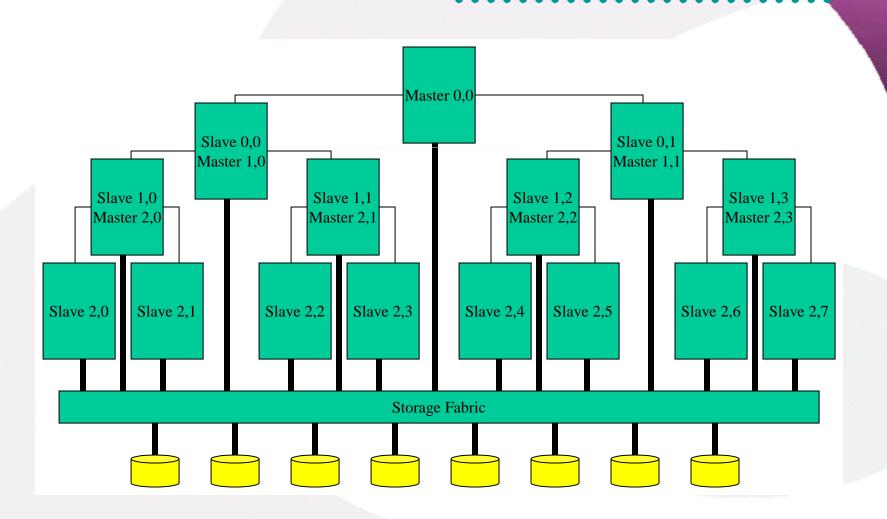
: Perspectives

- Application perspective Meta Data and User Data
 - On a single node
 - Distribute application across all compute nodes
- System perspective- Composite of all applications
 - Single compute node
 - Cluster
 - File System on a single compute node
 - File system distributed across multiple nodes
- Device perspective Composite off all applications and all systems
 - Host bus adapter
 - Storage Area Network
 - Disk Array
 - Disk Drive

: Workload generators

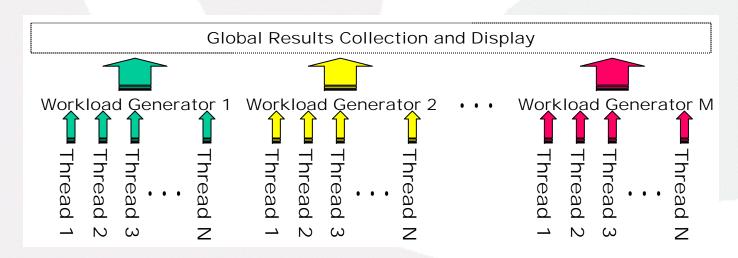
- Control mechanisms
 - Flat versus Hierarchical
- Operational Parameters
 - Request size
 - Number of transfers
 - Spatial access patterns
 - Temporal access patterns
 - Read/write ratios
- Indirect operational parameters
 - Memory allocation
 - Processor allocation
 - Process priority
 - Synchronization
- Synchronized / Unsynchronized workload generators
 - High resolution common reference clock

Benchmark Control Hierarchy



: Performance Data Collection

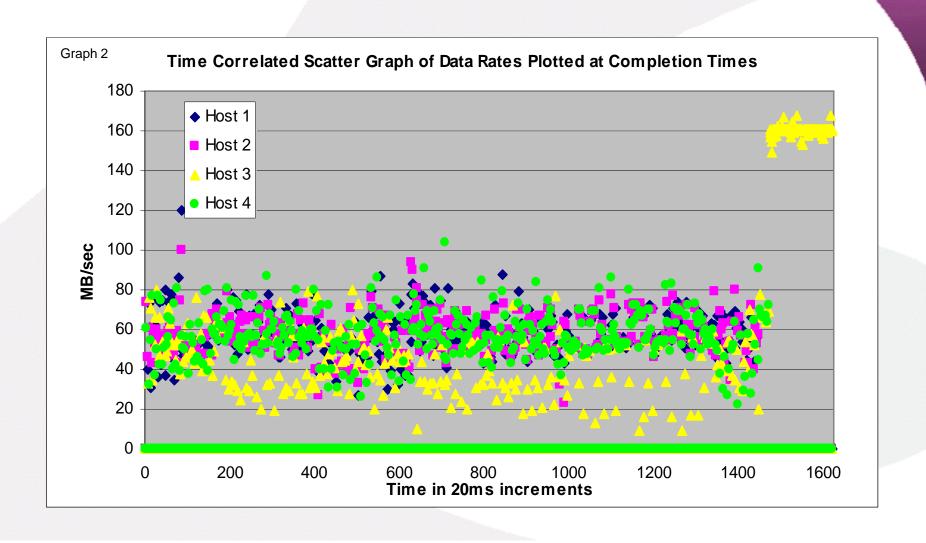
- A fully deployed I/O benchmark would need to run nearly 10,000 I/O threads, each generating results that need to be collected, condensed, and displayed
- The network I/O traffic for collecting the results in real time and/or post mortem is significant
- The performance results data collection process cannot interfere with the data transfer for the benchmark



Run-time Monitoring and Report Generation.....

- Detailed reports from
 - Each thread performance of an individual thread
 - Each node aggregate performance of all threads on a node
 - entire system aggregate performance of all threads on all nodes across the system
- Real-time (run-time) time-correlated reports interactive displays and visualization of traffic, performance, and bottlenecks
- Trace data analysis tools Post mortem analysis and visualization
- Bottleneck Isolation tools real-time and post-mortem
- Summary reports for "benchmark" purposes

Example of Time-Correlated Performance Data.....



: Summary

- The *process* of designing and running an I/O benchmark program that is attempting to
 - mimic the behavior of an application or a class of applications,
 - interpreting the results
 - Provide information that can be used to fine tune the I/O subsystem and/or file system(s)
- Provide detailed, real-time system-, application-, and nodewide perspective I/O monitoring capability for identifying performance bottlenecks of the benchmark
- Tight management of the variables that influence the I/O performance during a benchmark run



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