

A Simple Mass Storage System for the SRB Data Grid

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- Motivations for implementing a Mass Storage System in a data grid
- Simple MSS design
- An Overview of SRB architecture and features
- Simple MSS implementation

Motivations for a MSS

- Cost of licensing commercial MSS systems.
- Efficiency and performance
 - Tightly integrated with SRB data grid functions
 - Use SRB feature such as file replication, server directed parallel I/O, latency management.
- Elimination of duplication of features
 - Better disk cache utilization
 - Single authentication scheme
 - SRB data grid abstractions logical name space, storage repository abstraction, logical storage resource naming and metadata management
- Provide a storage system that spans remote caches and distributed archival devices

Simple MSS - Design Goal

- A distributed farm of disk cache resources backed by a tape library
 - Use a pool of distributed cache resources to form one large cache resource
 - A tape library system to control the mounting and dismounting of tapes.
 - Support Storage Tech silo running ACSLS software
- Physical location (on cache or tape) transparency
 - User uses same access mechanism to access data independent of data location

Simple MSS - Design Goal

- Files should always be staged to cache before I/O operations
 - Tool to manage disk cache
- Large files should be stored in segments.
 - handle files of very large size
 - parallel data transfer between tapes and cache system can be implemented (not yet implemented)

Simple MSS – Components Needed

- A client-server architecture
 - authentication scheme
 - a framework for client-server and server-server communication
 - A federated server system cache and tape resources located on different hosts.
- A metadata server that maintains
 - logical POSIX-like name space
 - mapping of each logical file name to its physical location.

Simple MSS – Components Needed

- Server infrastructure and meta data that
 - allow files stored in the MSS to appear the same as other files stored on disk caches
 - translate user requests to physical actions using metadata
 - driver functions for basic tape I/O operations.
 - driver functions for basic cache I/O operations.
 - Functions to stage files from tape to cache and dump files from cache to tapes.
 - Data transfer between the cache system and clients.

SRB Architecture

• Federated middleware system

- Client/server model
 - Federated servers with uniform interfaces
 - Access to all resources in the federation
- MCAT metadata catalog
 - Information repository abstraction
 - Client access through SRB server

Simplified SRB Model



SRB Server Design

- Three layer architecture
 - Top layer (client abstraction layer)
 - Interacts with clients and other servers through tcp/ip sockets
 - User authentication
 - Handle function requests parses requests and calls handlers in middle and bottom layers.
 - Middle layer (logical layer)
 - Input parameters are in their logical representations (logical name space, logical resource name)
 - Queries MCAT, translates from logical to physical representations (e.g., host address, type of resources, actual path where the file is located)
 - Calls functions in the bottom (physical) layer to access resources.

SRB Server Design

- Bottom layer (physical layer), handles three types of resources
 - File system
 - Drivers for the 17 POSIX functions (creat, open, read, write..)
 - FS supported : UNIX, HPSS, DPSS, UniTree, gridFTP (to be released), SRB's own tape library system (to be released)
 - Include/exfSw.h defines driver function mapping
 - DB large objects
 - Similar to files except stored in DB
 - Drivers for Oracle, DB-2 and Illustra
 - DB tables
 - Access external DB tables (query, insert, ...)

Federated SRB Operation



- Abstraction of User Space
 - Single sign-on
 - No need for UNIX account on every systems
 - Multiple authentication schemes supported
 - Certificates, (secure) passwords, tickets, group permissions
 - Robust access control
 - User level, grant access to multiple users
 - Group level
 - Tickets

- Abstraction of Data and Collections
 - Logical Name space -
 - UNIX like directories (collections) and files (data)
 - Mapping of logical name to physical attributes host address, physical path.
 - Single logical name space mapped to data on multiple resources
 - POSIX like API for making collections (mkdir) and data creation (creat)
 - Data replication
 - replica on different resource
 - same logical name but different replica number

- Virtualisation of Resources
 - Mapping of a logical resource name to physical attributes: Resource Location, Type & Access transparency
 - Client use a single logical name to specify a resource
 - Resource group (logical resource) bundling of resources, automatic replication, transparent caching/archival storage
- Uniform Access Methods
 - Use same APIs, Command Line, GUI Browsers, Web-Access (Portal,WSDL, CGI) to access various resources

- Performance enhancement
 - Client and server-driven Parallel I/O strategies
 - Interface with HPSS's mover protocol for parallel I/O
 - Container physical grouping of small files for tape I/O

Containers

- Physical Grouping of Objects
- Similar to tar but many more features
 - Each inContainer object appears in logical name space
 - Can be accessed as normal object
 - Data objects in a container moved together
 - To aid access patterns
 - To take advantage of resource characteristics
 - Automatic staging(caching) and Archiving sync to backend
- Family of Containers
 - New container automatically created when filled
- Containers for Collections
 - Associate a collection with a container

Access Control

- Access controls on logical names
 - Datasets
 - Collections
 - Resources
- Multi-level access
 - Read, Annotate, Write, Curate, Own
- Access control for users and groups
- Ticket-based access control
- Audit access

Simple MSS – Components Needed

• A tape library server to schedule, mount and dismount tapes.

• A tape database that tracks the usage of all tapes controlled by the MSS.

• A set of tape management utilities to initialize and migrate data

- Use many existing features in SRB
 - Federated client/server architecture
 - MCAT logical POSIX-like name space
 - Parallel I/O for data movement
- Mechanisms added to the SRB
 - A tape library server for mounting and dismounting tapes
 - Drivers for Tape I/O in SRB server
 - metadata needed to manage files on the MSS.
 - functions to stage files from tape to cache.
 - Tape management functions and utilities.

- Compound resource
 - Appear as a single resource to user
 - Internally:
 - A pool of distributed cache resource for front end
 - A tape library system resource for back end

- Compound object
 - Object stored in compound resource
 - Behave as a single normal SRB object
 - I/O always always done on cache resource first
 - When object opened for read/write:
 - Check whether already on cache.
 - Stage to one of the cache resources.
 - Dirty bit set if object modified
 - API for synchronizing dirty copies to tape and purging cache copies by sys admin
 - File locking on object during I/O

- Tape library server
 - An independent server
 - For STK silo running ACSLS software
 - Use the same SRB server framework
 - Authentication scheme
 - Client-server function call
 - Scheduling, mounting and dismounting tapes
 - Handle multiple drive types
 - Handles only 3 function calls
 - mount tape
 - dismount tape
 - report priorities by drive types (based on number of drives idling and in use)

- Other supporting features
 - Drivers in SRB servers for tape I/O operations
 - A database schema and metadata that tracks tape usage
 - Tape labels controlled by MSS
 - Current tape positions
 - Bytes written for each tape
 - Full flag
 - Functions to insert, modify and query the tape meta data
 - Tape management utilities tape initialization, data migration, tape metadata manipulation and query, etc.

Comparison to IEEE MSSRME

- Provides similar functionality
- Implementation similarities and differences.
 - Differences attributed to SRB abstraction mechanisms
 - SRB MSS uses the underlying File System for managing data storage
 - Reference Model maps bitfiles to the logical and physical volume abstractions
 - Name Server Similar to SRB's POSIX-like name space
 - Reference model's Bitfile Server, Storage Server and Mover are combined into a single SRB resource server
 - Similar Migration-Purge facility

Simple MSS - Status

• First version released with SRB 2.0 (2/19/03)

• Validated through significant testing. About to enter production at SDSC

- Performance on SDSC resources
 - Cache to client up to 40 Mbytes/sec
 - Tape to cache 5-7 Mbytes/sec (3590 tape)

More Information

http://www.npaci.edu/DICE