



BabuDB: Fast and Efficient File System Metadata Storage

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Motivation

- Modern parallel / distributed file systems:
 - Huge numbers of files and directories
 - Many storage servers but few metadata servers
- **Examples:**
 - Lustre, Panasas Active Scale, Google File System
- **Metadata access critical wrt. system performance**
 - ~75% of all file system calls are metadata accesses
 - Metadata servers are bottlenecks

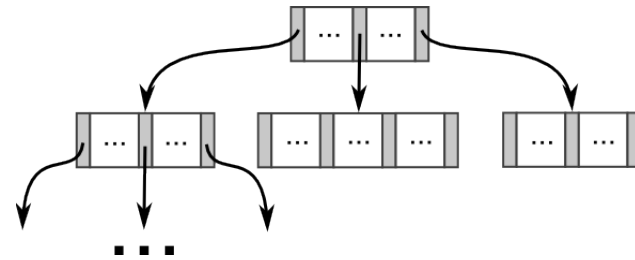
Motivation

- B-tree-like data structures used for metadata storage

- ZFS, btrfs, Lustre, PVFS2

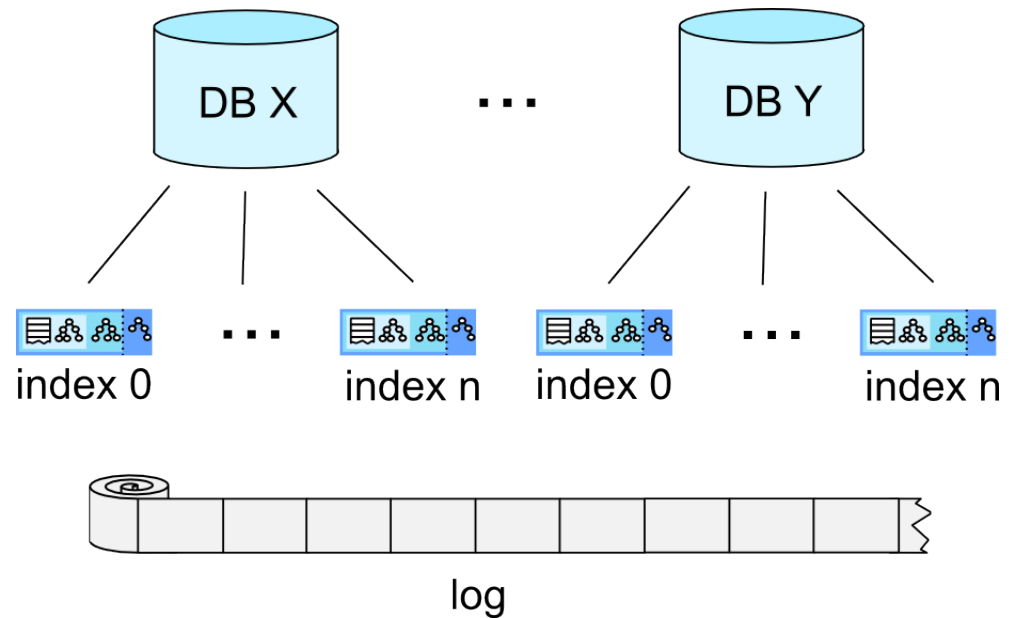
- Downsides:

- Hard to implement and test, high code complexity
 - Multi-version B-trees even more complex
 - On-disk re-balancing expensive



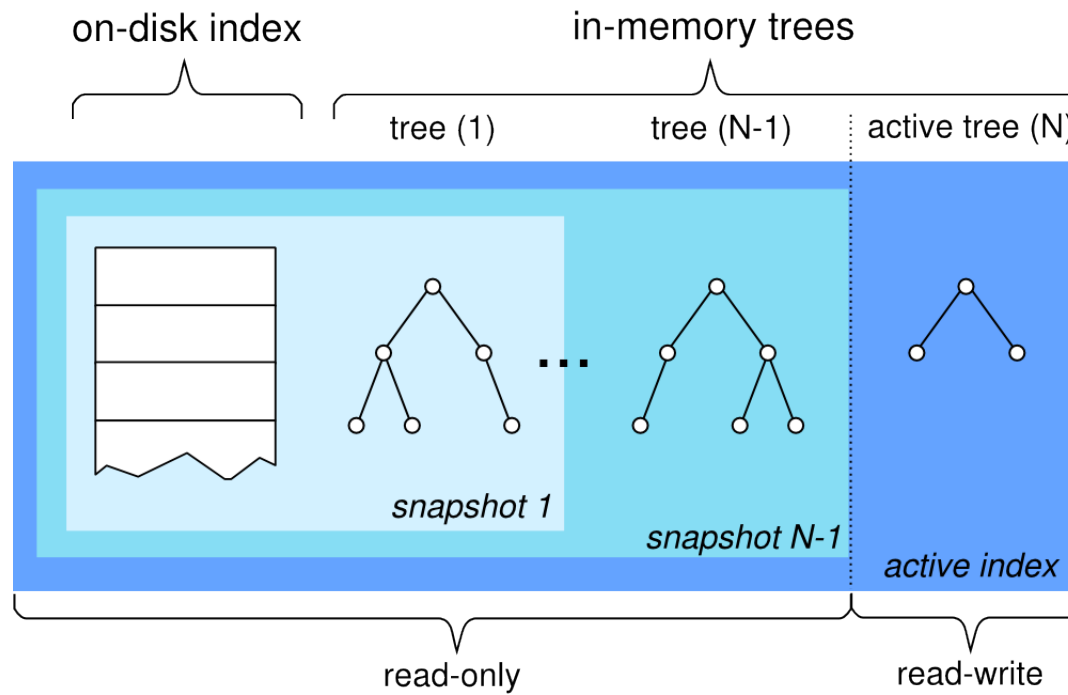
BabuDB

- Key-value store

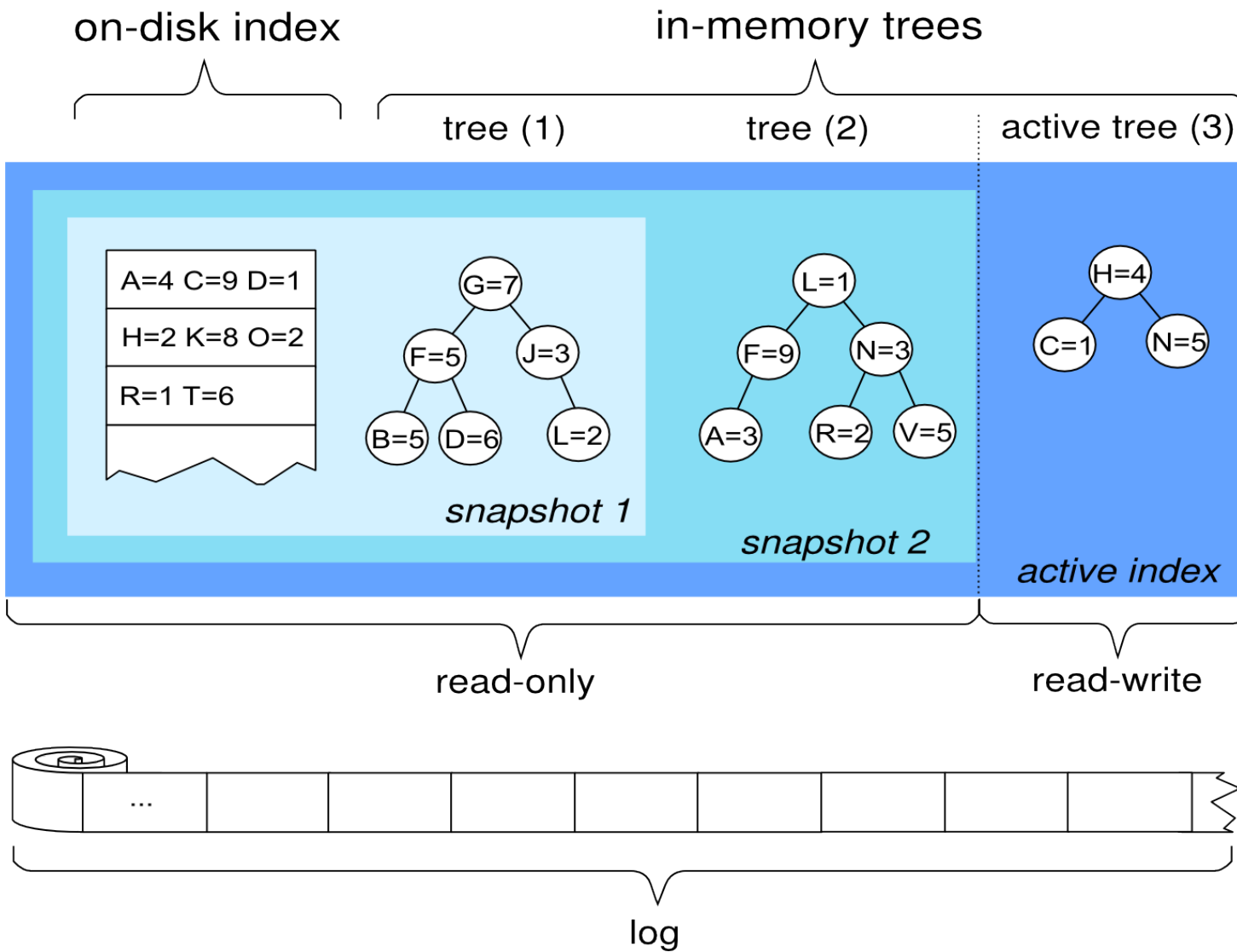


- FS metadata: key-value pairs stored in DB indices

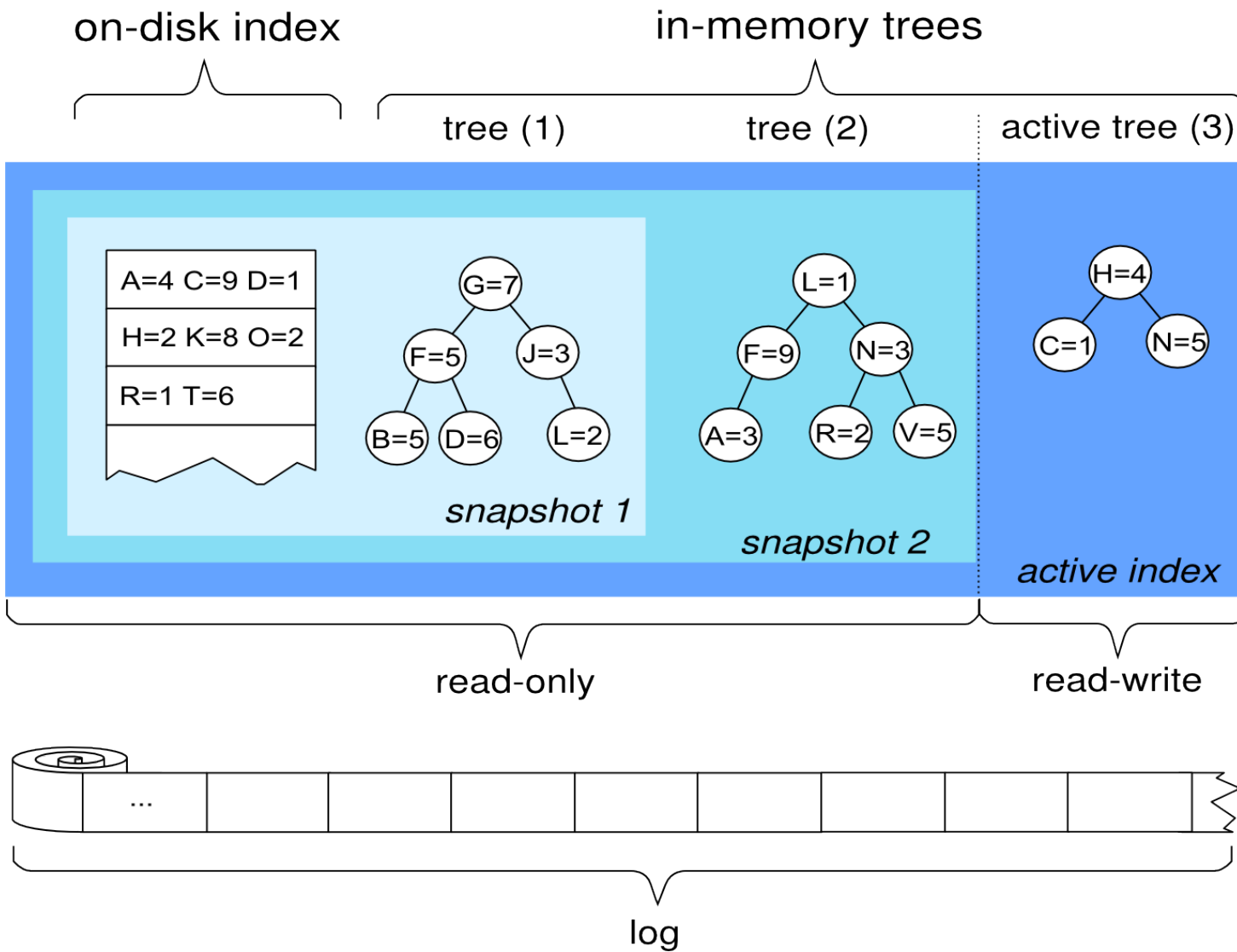
BabuDB: Index



Example

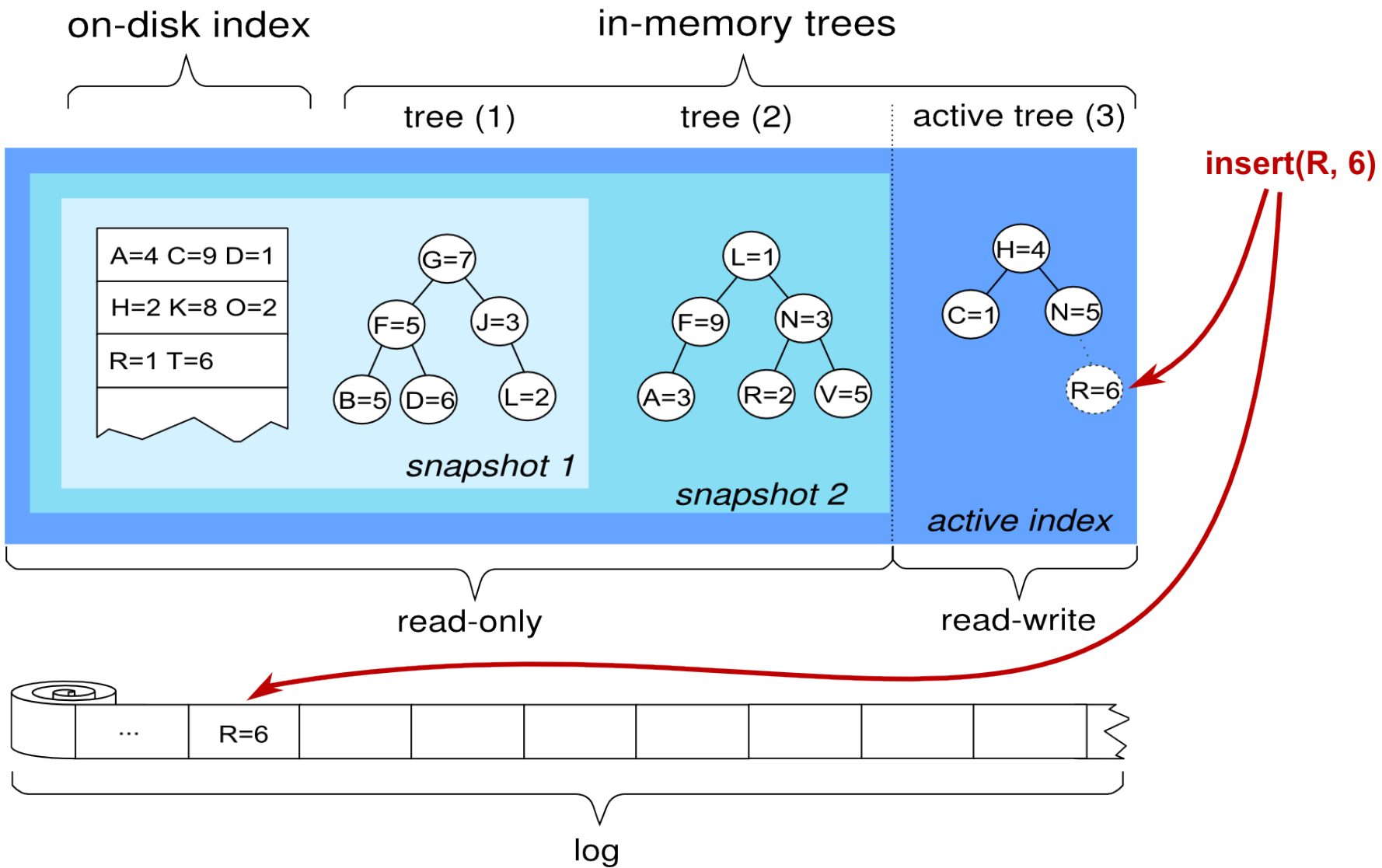


Example: Insertions

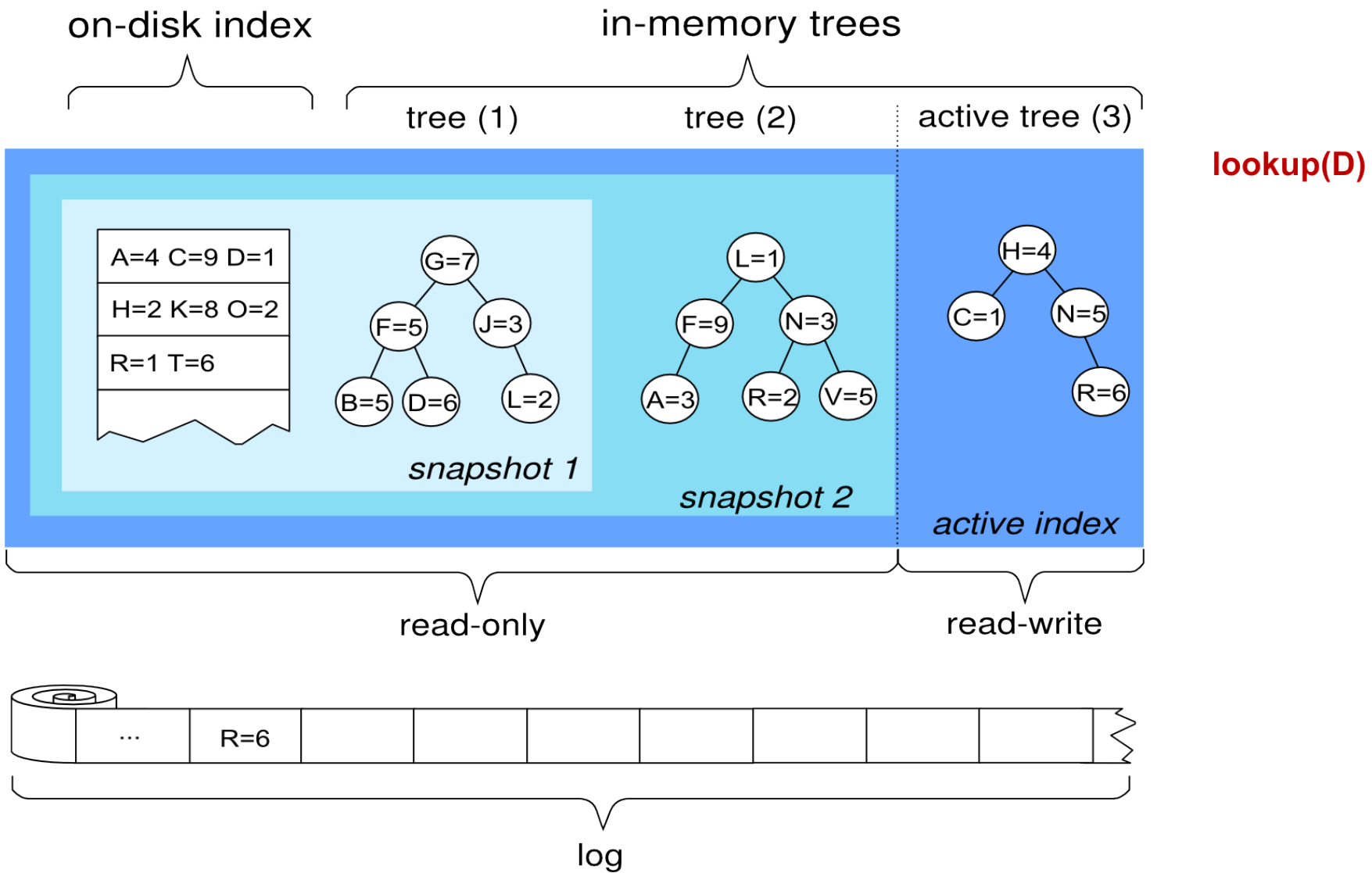


insert(R, 6)

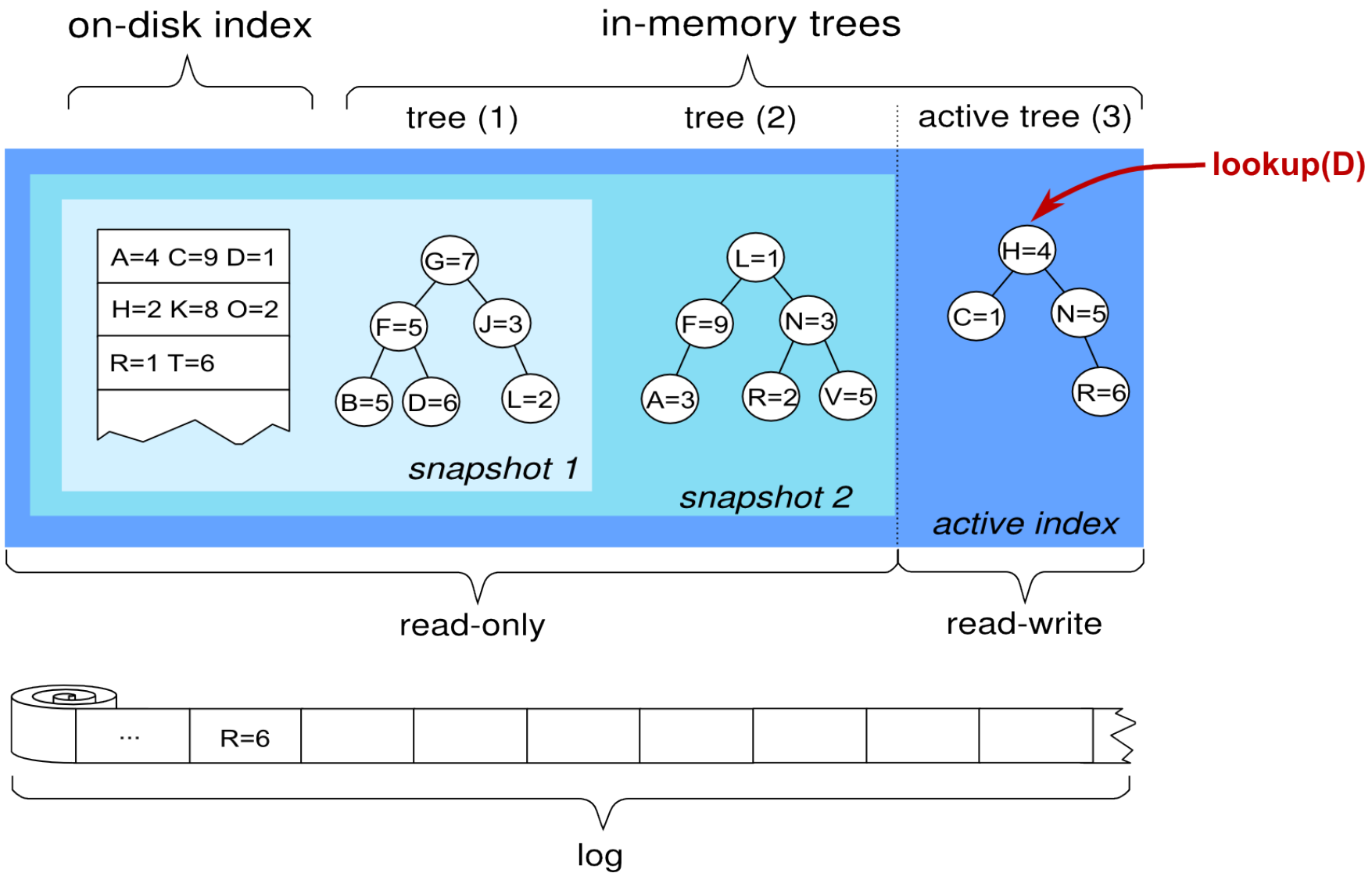
Example: Insertions



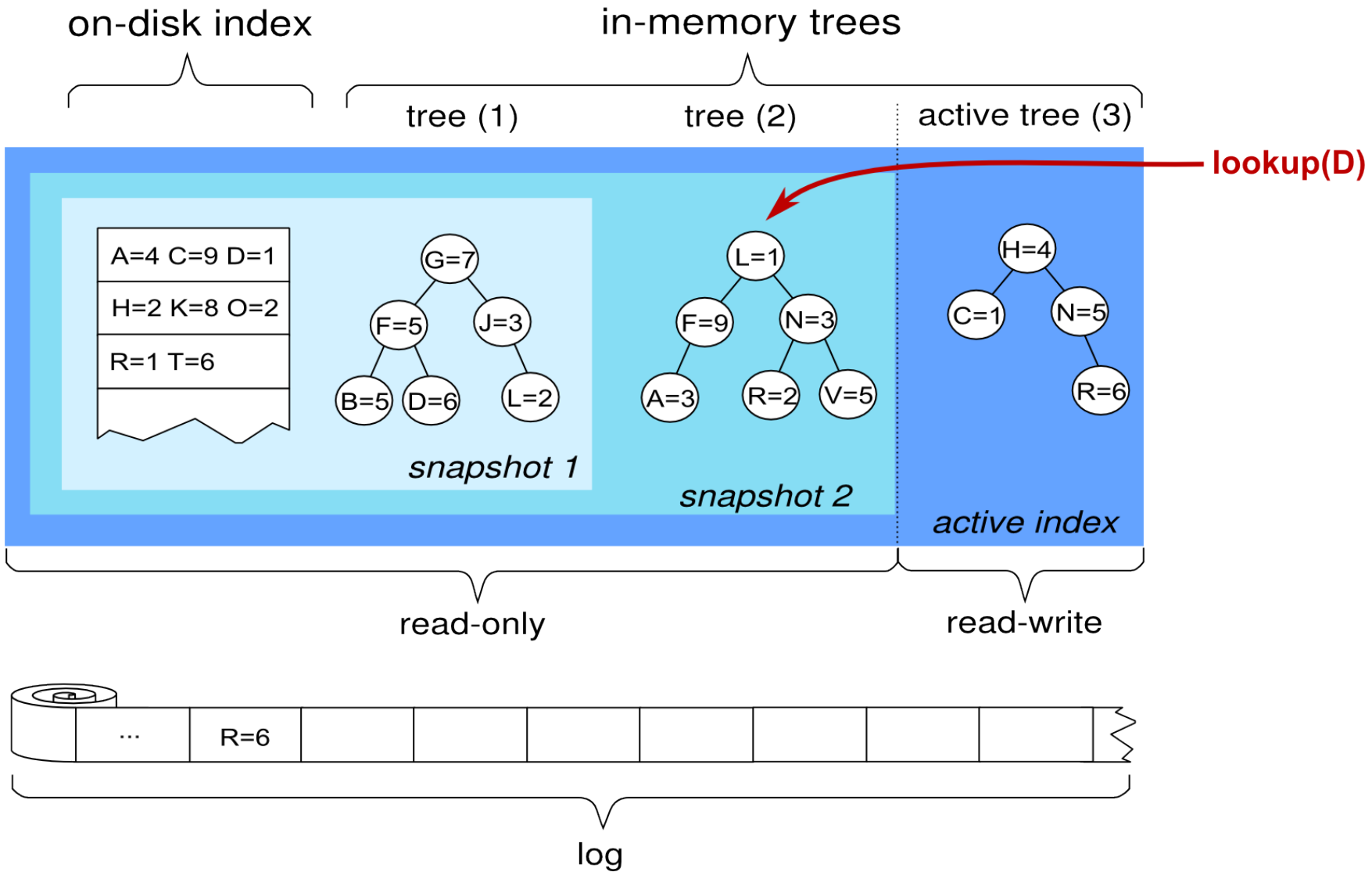
Example: Lookups



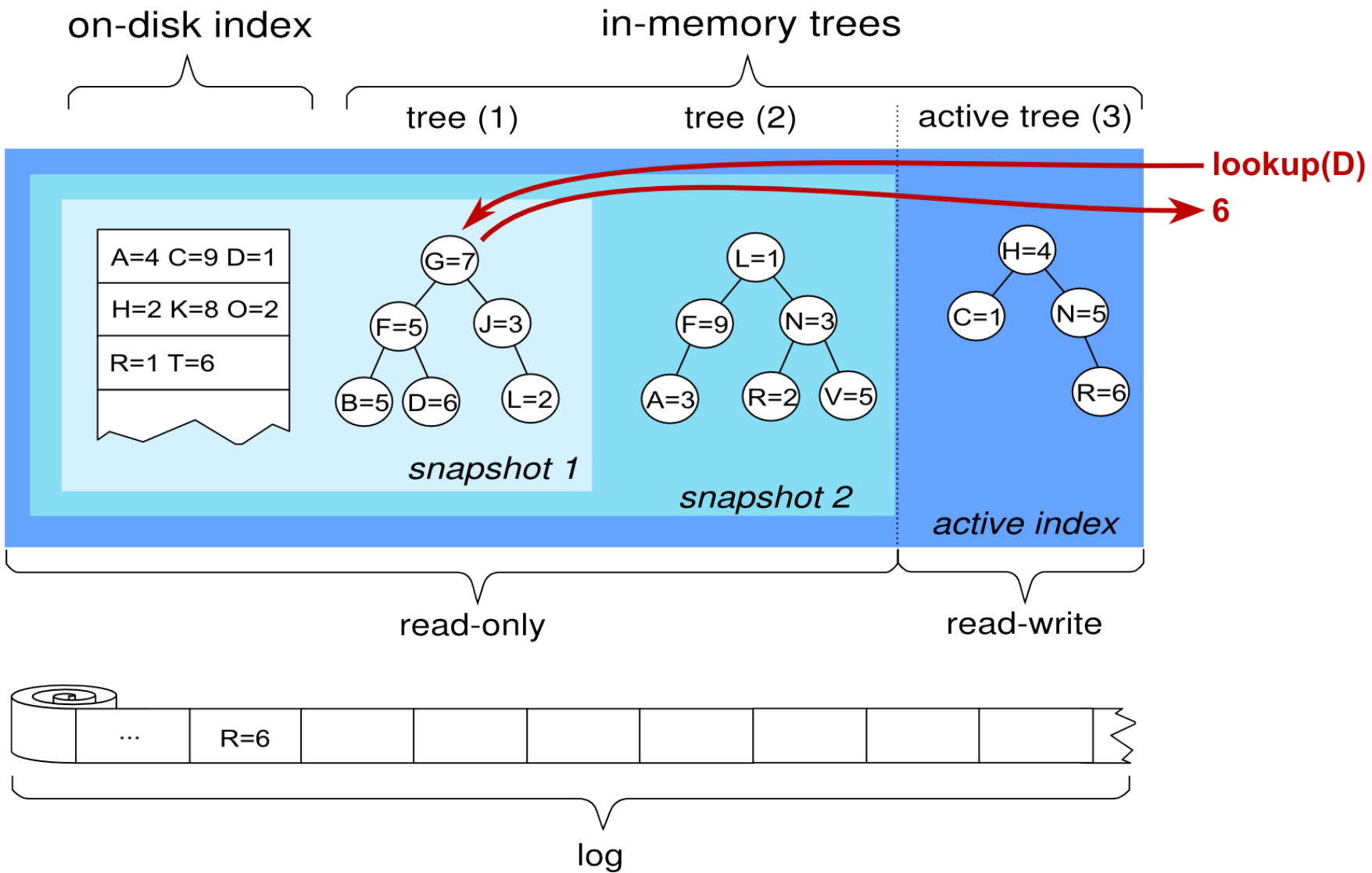
Example: Lookups



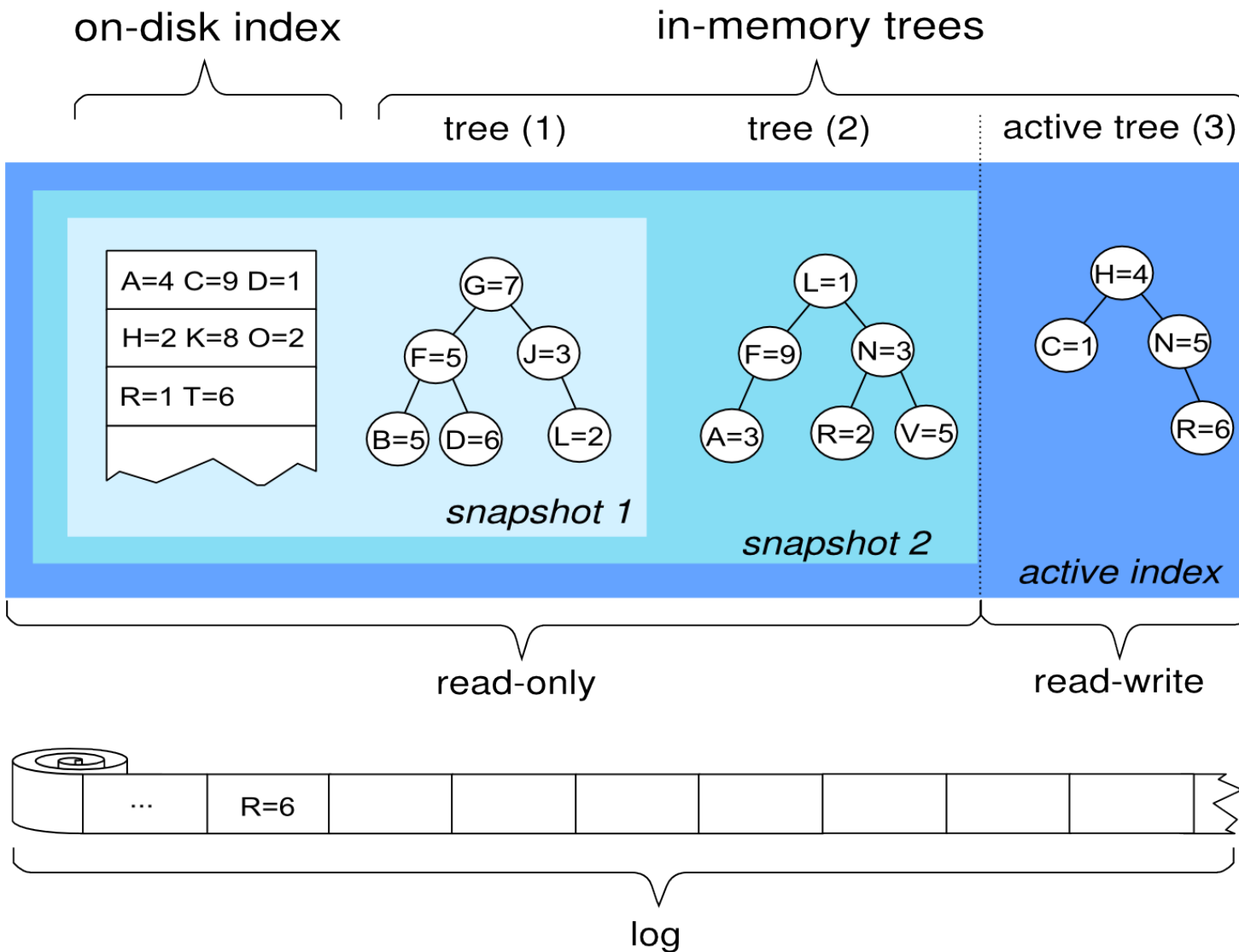
Example: Lookups



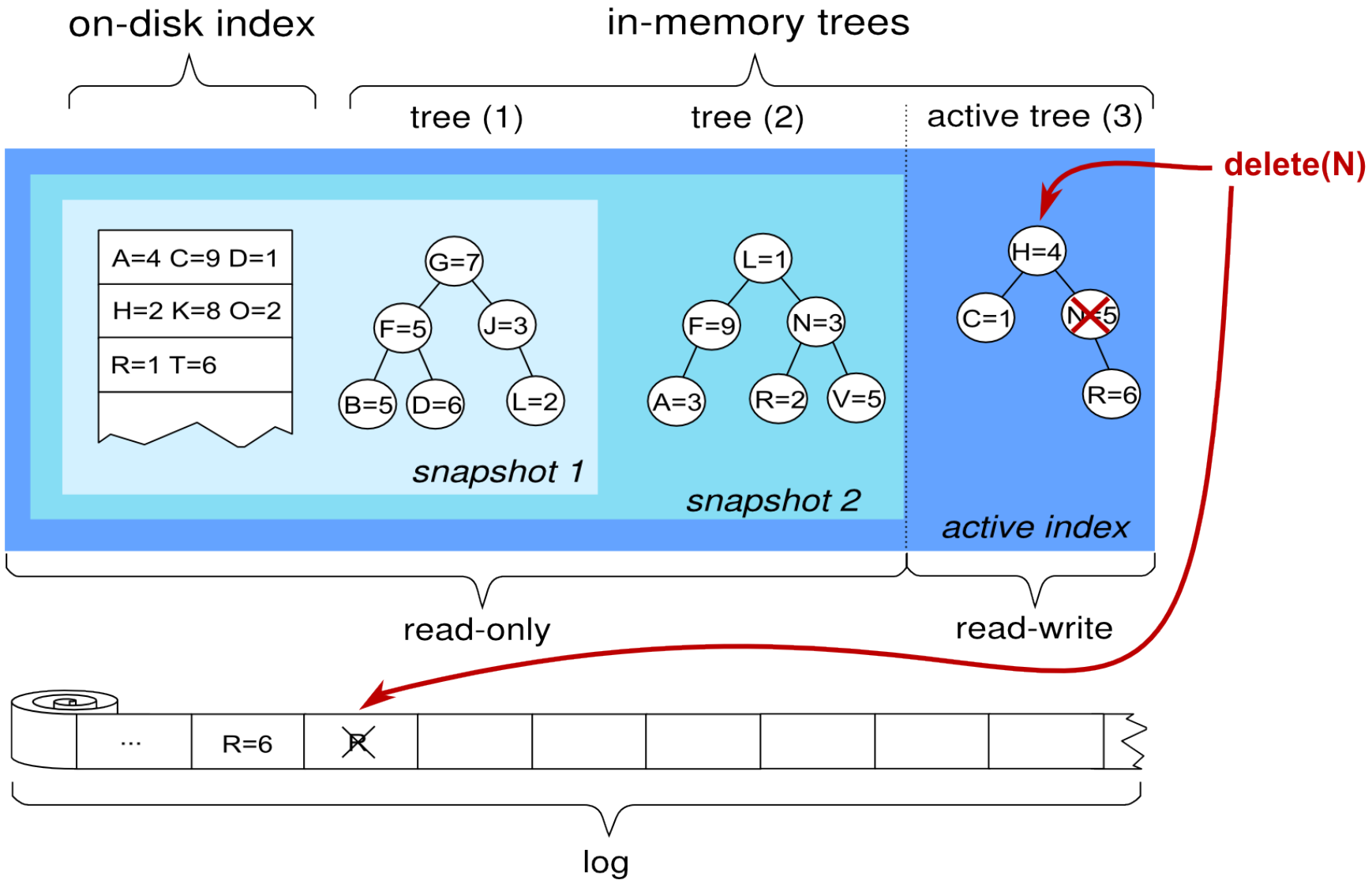
Example: Lookups



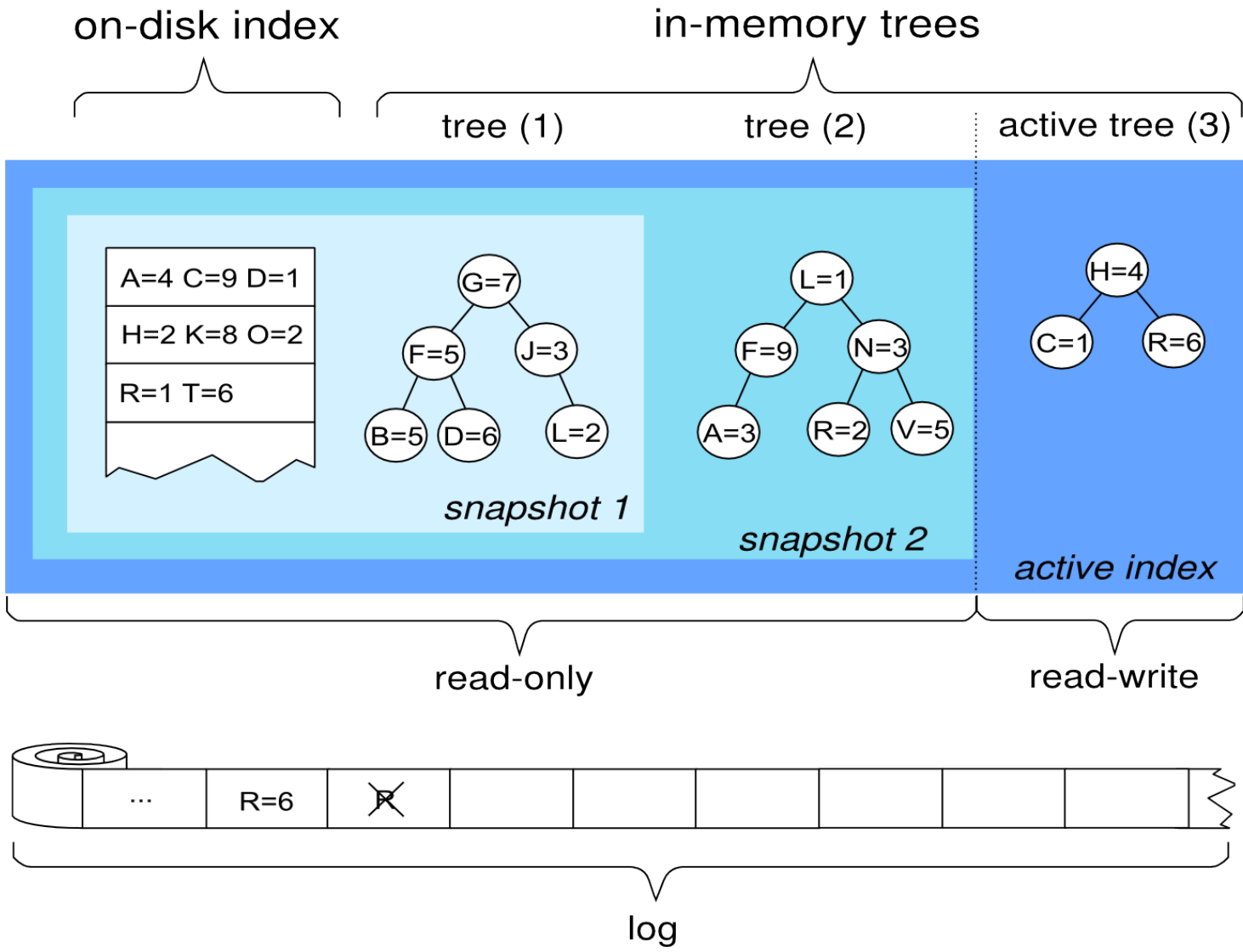
Example: Deletions



Example: Deletions

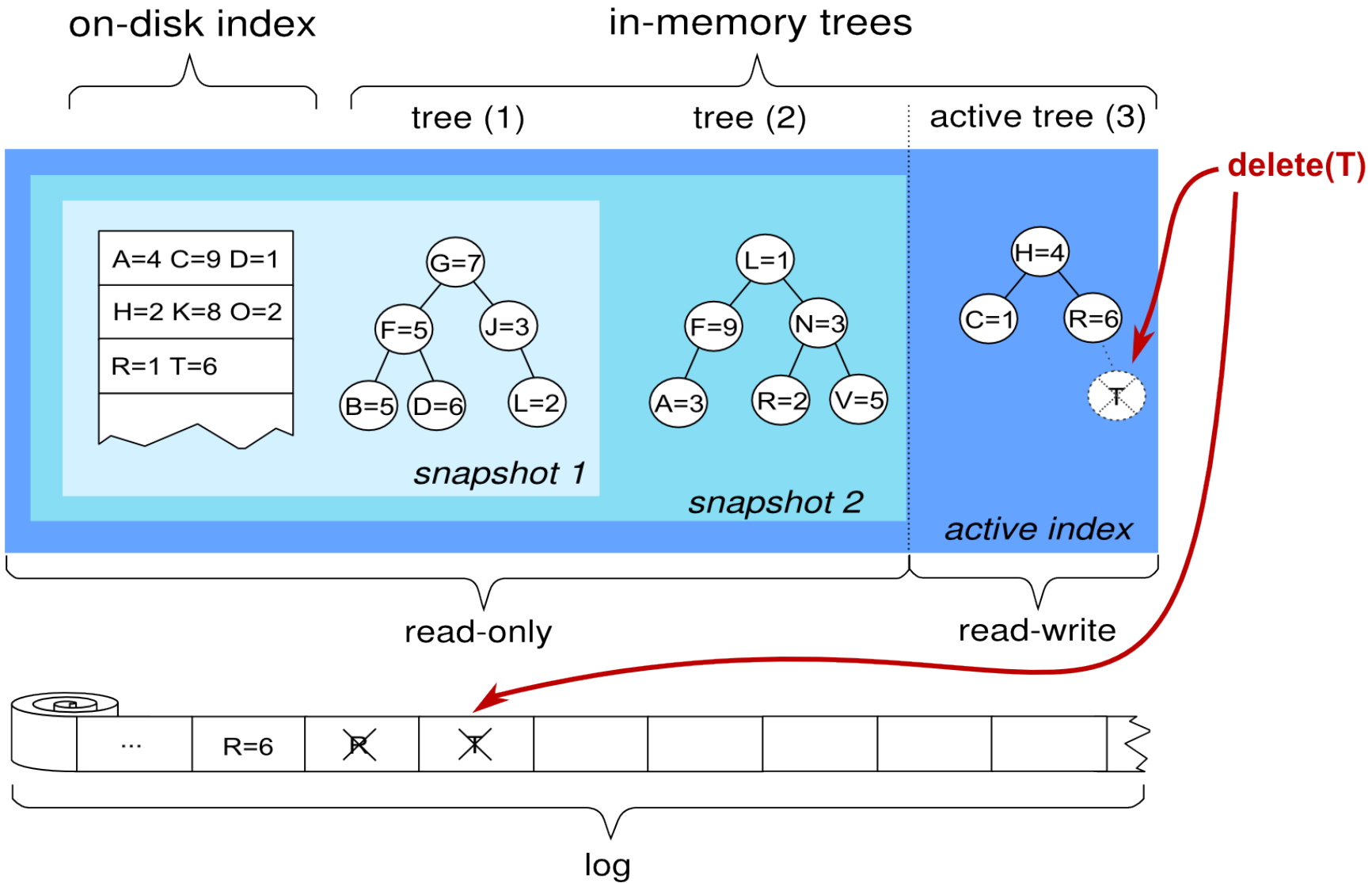


Example: Deletions

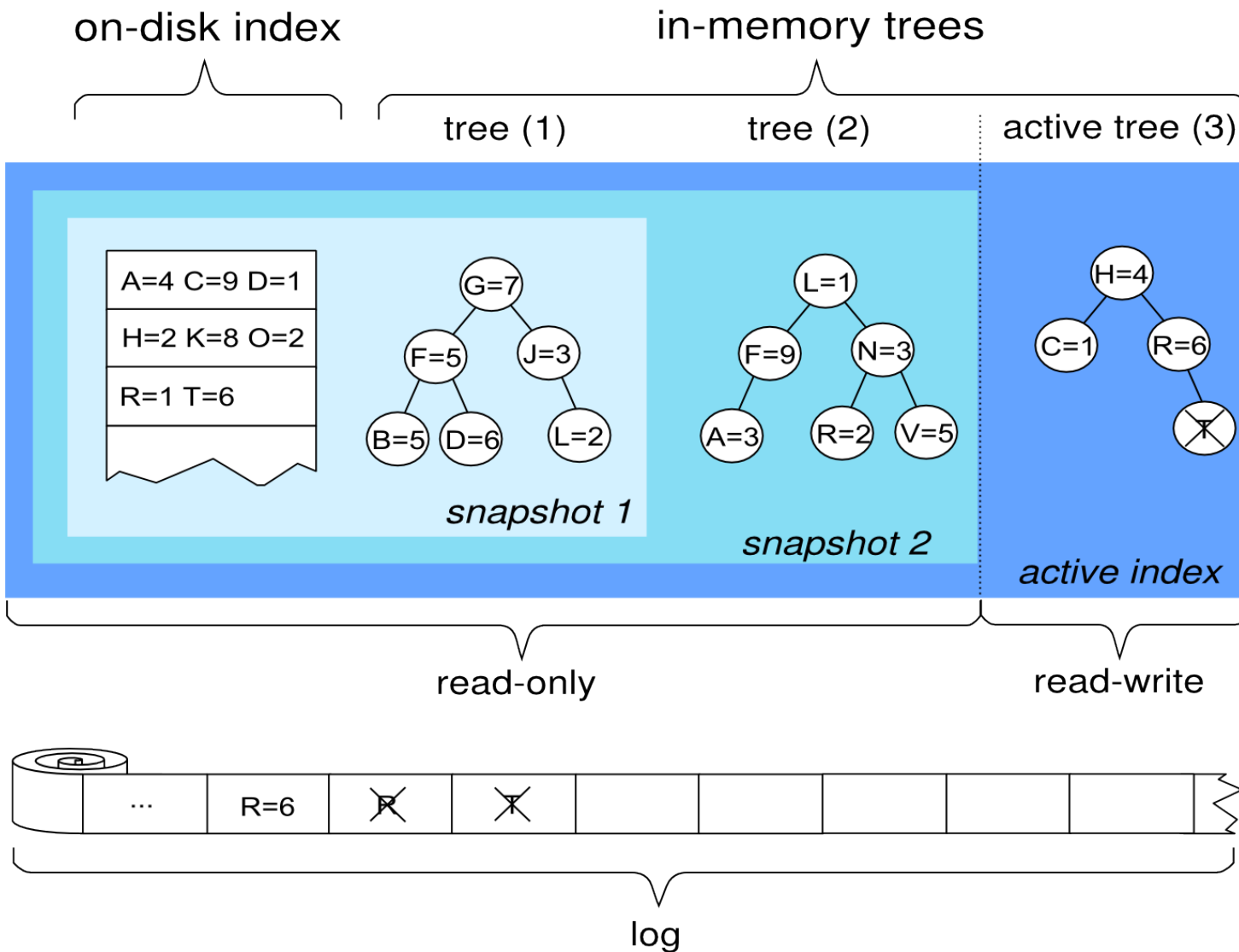


delete(T)

Example: Deletions

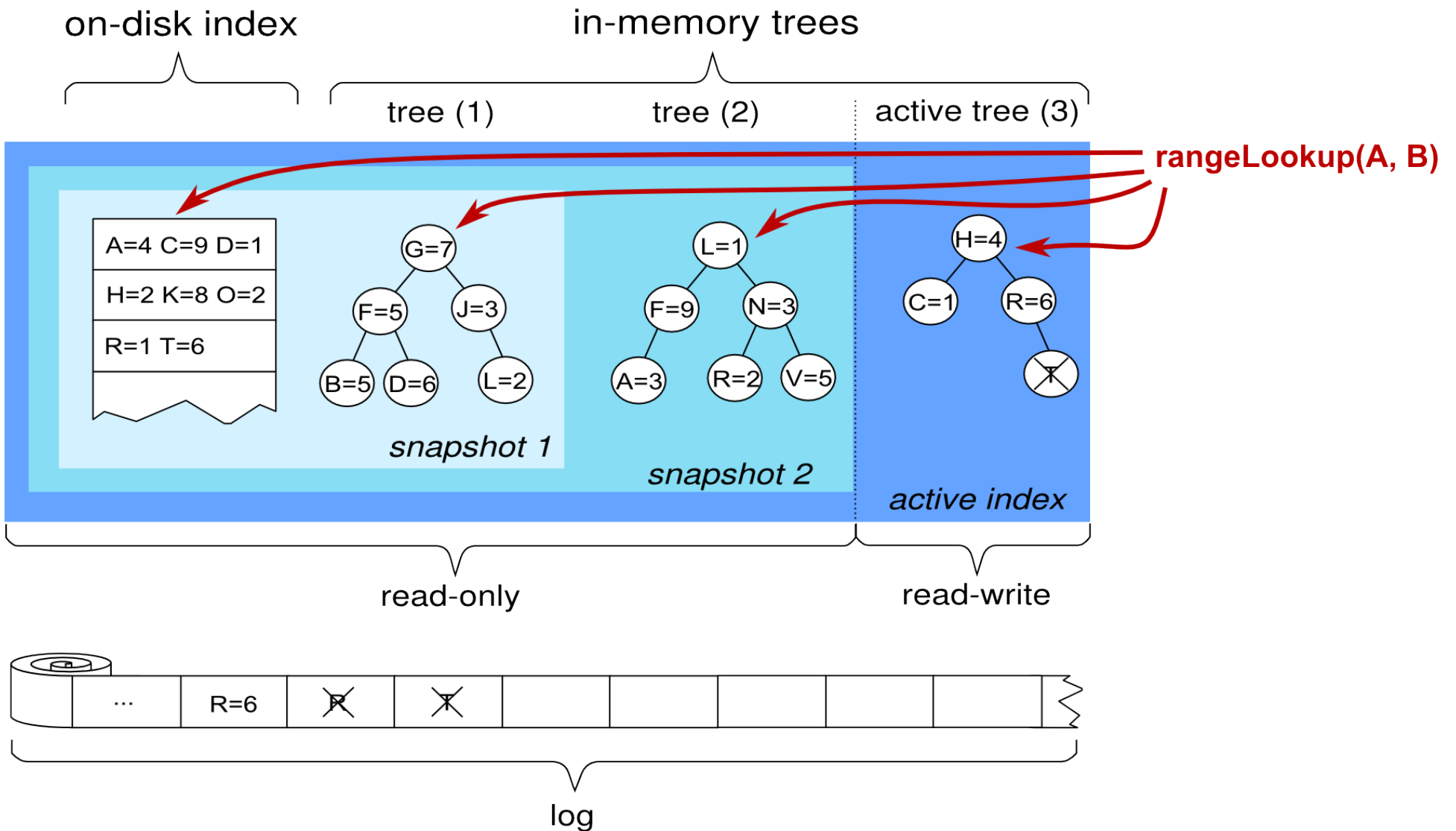


Example: Range Lookups

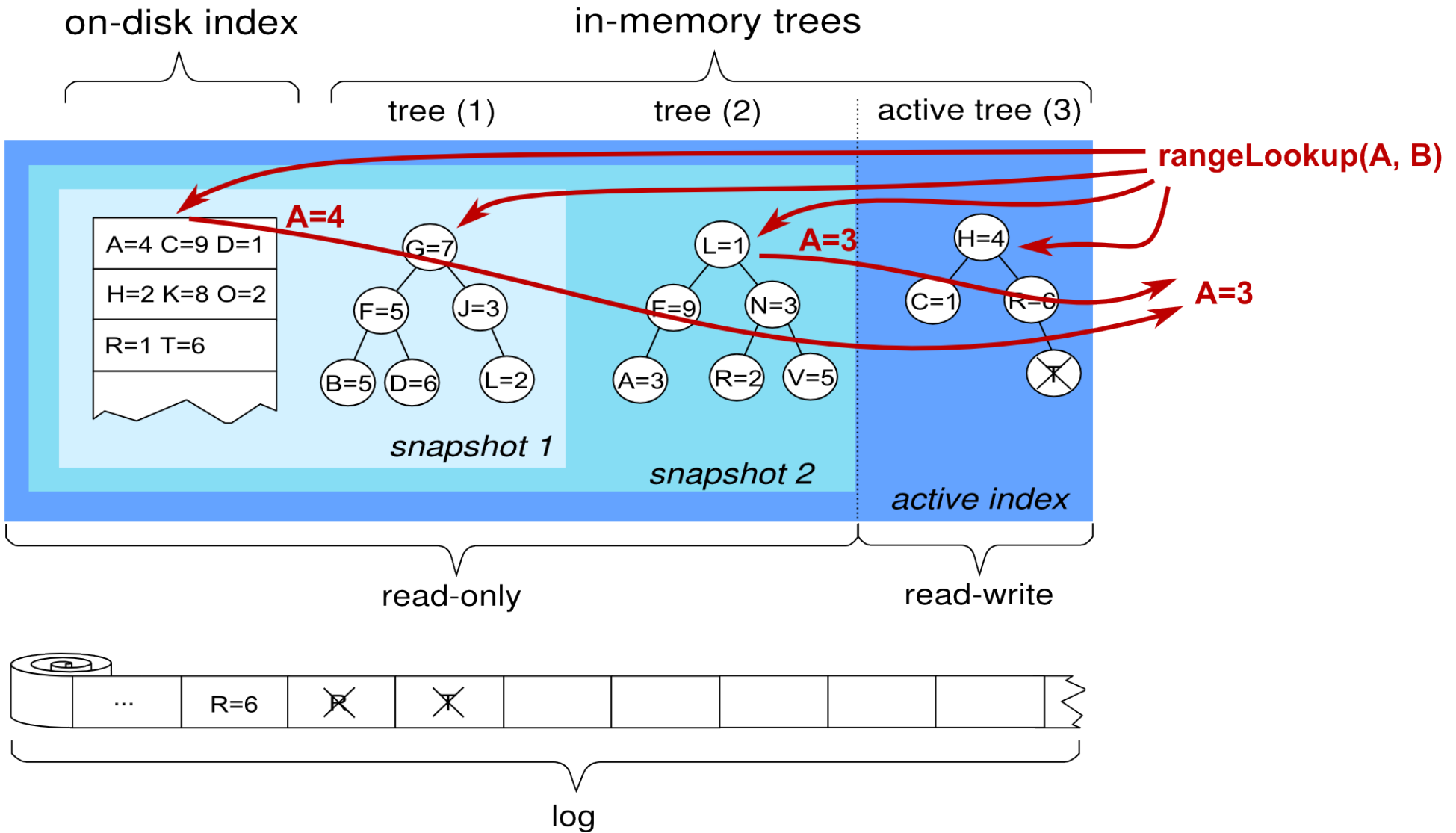


rangeLookup(A, C)

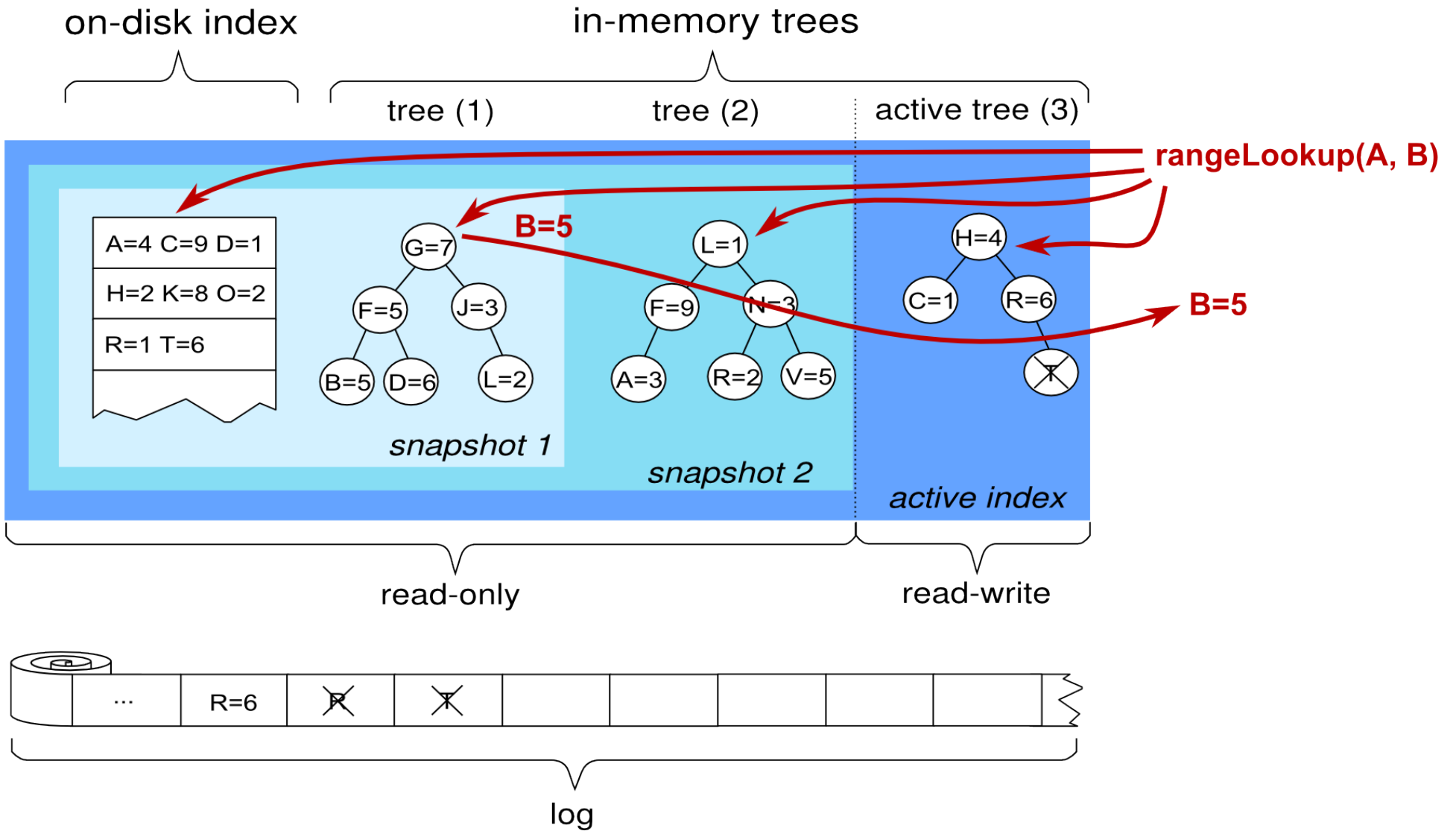
Example: Range Lookups



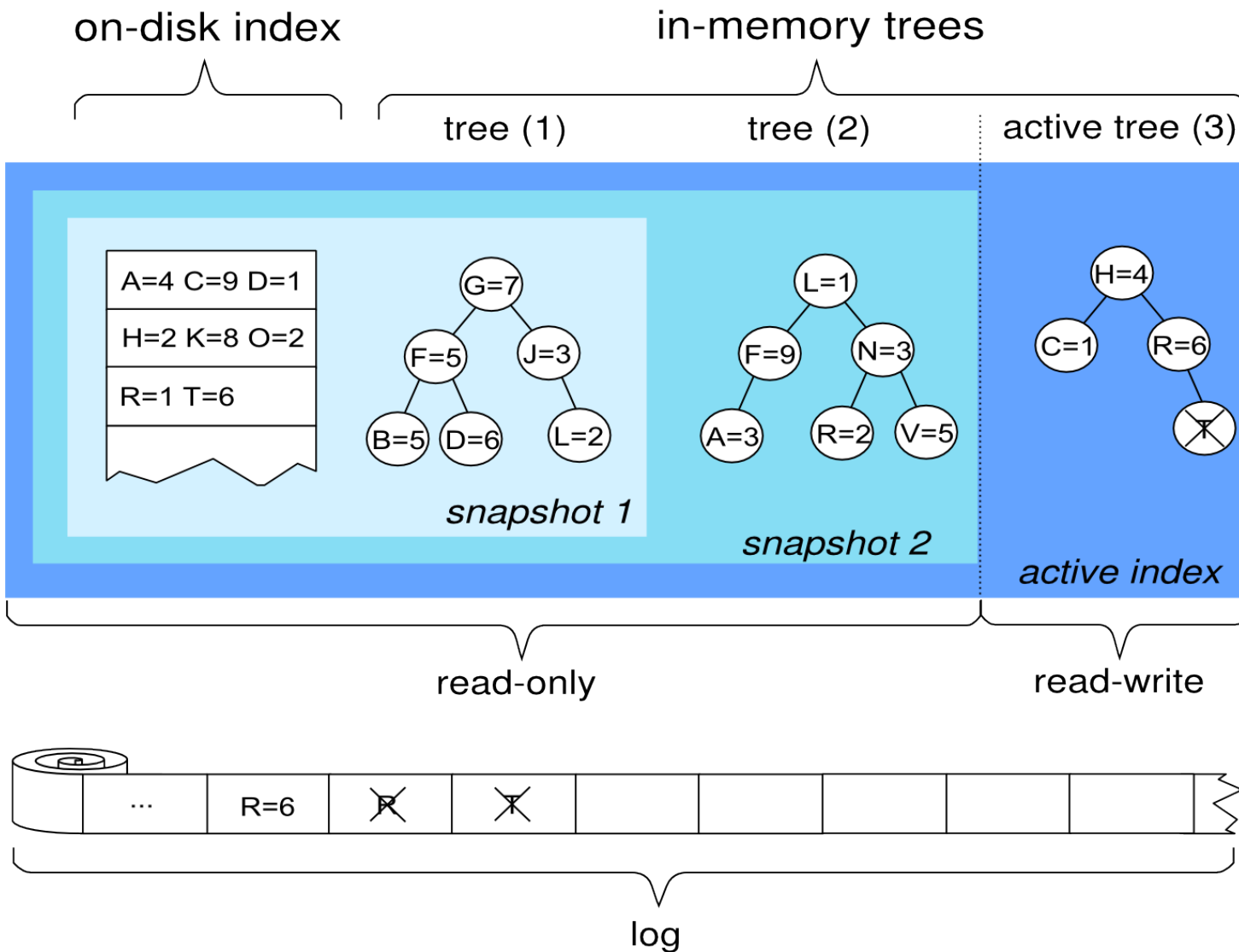
Example: Range Lookups



Example: Range Lookups



Example: Checkpoints



Example: Checkpoints

on-disk index

in-memory trees

tree (1)

tree (2)

tree (3)

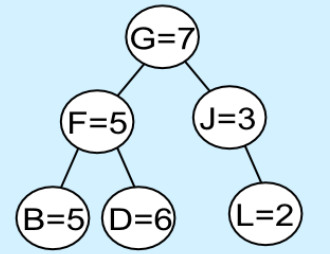
active tree (4)

checkpoint()

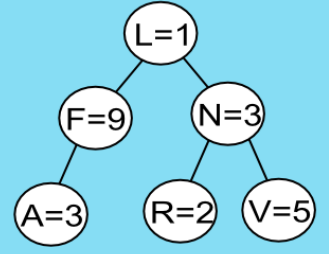
1. create snapshot



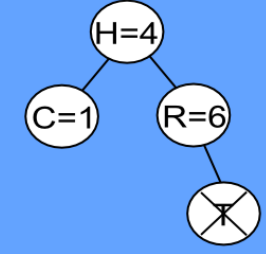
A=4 C=9 D=1
H=2 K=8 O=2
R=1 T=6
~



snapshot 1



snapshot 2

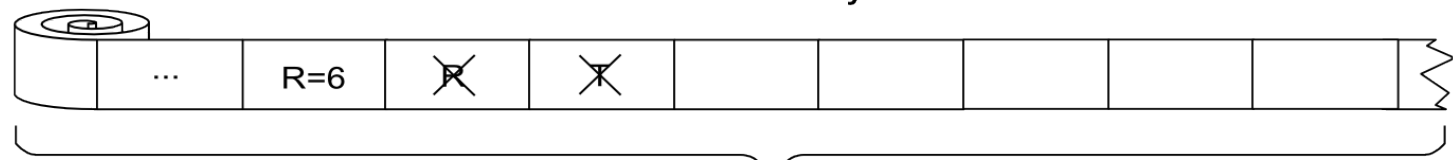


snapshot 3

active index

read-only

read-write



log

Example: Checkpoints

on-disk index

in-memory trees

tree (1)

tree (2)

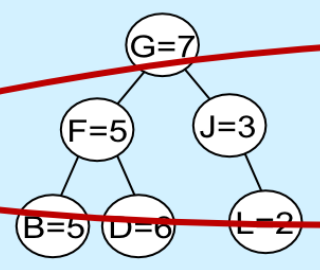
tree (3)

active tree (4)

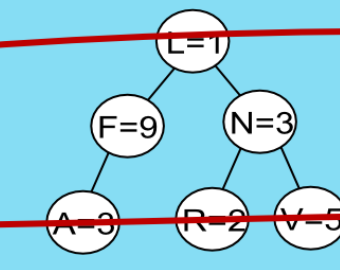
checkpoint()

1. create snapshot
2. rangeLookup(*) at snapshot 3, write new on-disk index

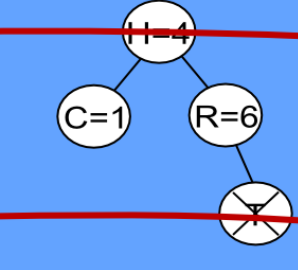
A=4	C=9	D=1
H=2	K=8	O=2
R=1	T=6	



snapshot 1



snapshot 2

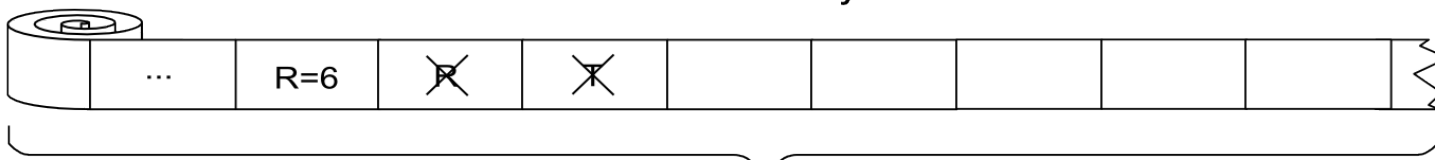


snapshot 3

active index

read-only

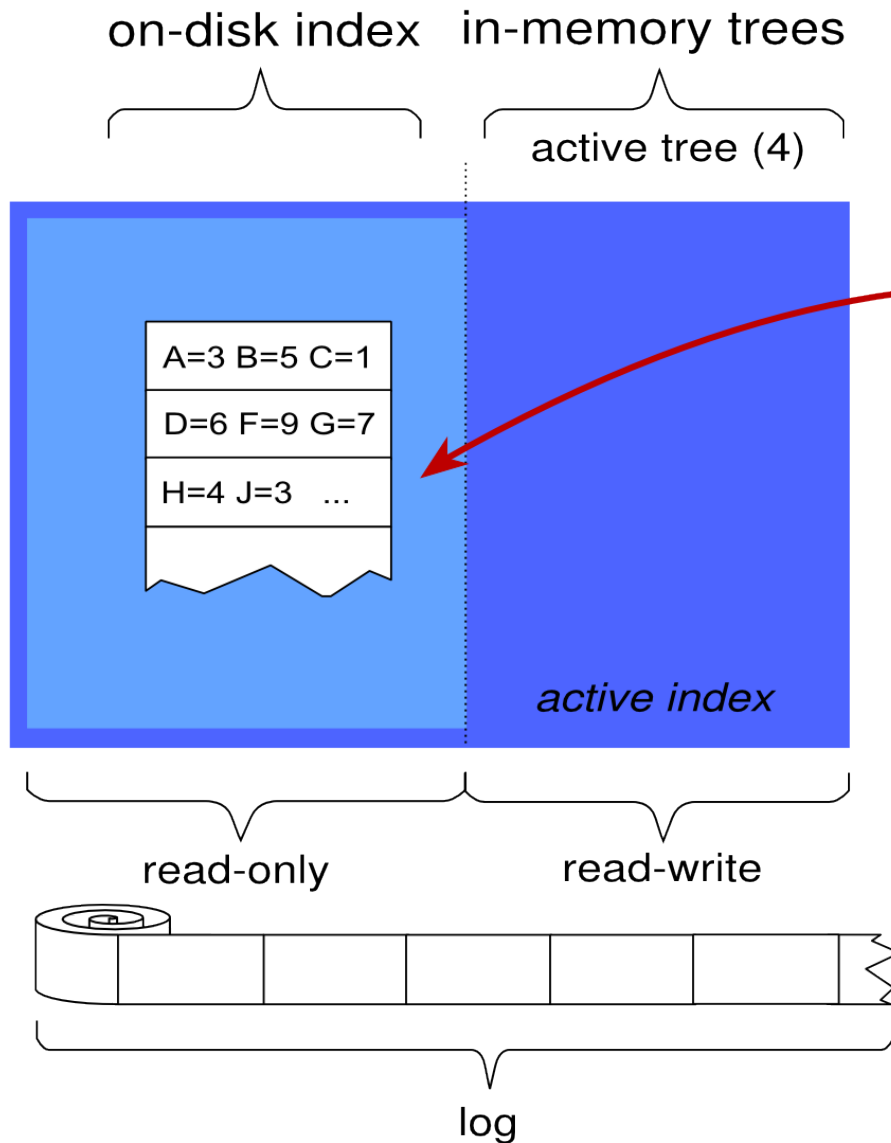
read-write



log

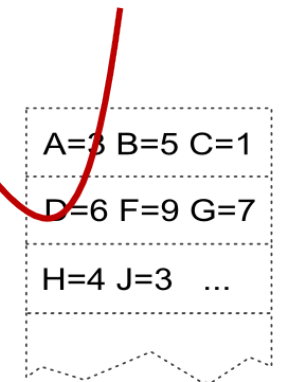
A=3	B=5	C=1
D=6	F=9	G=7
H=4	J=3	...

Example: Checkpoints

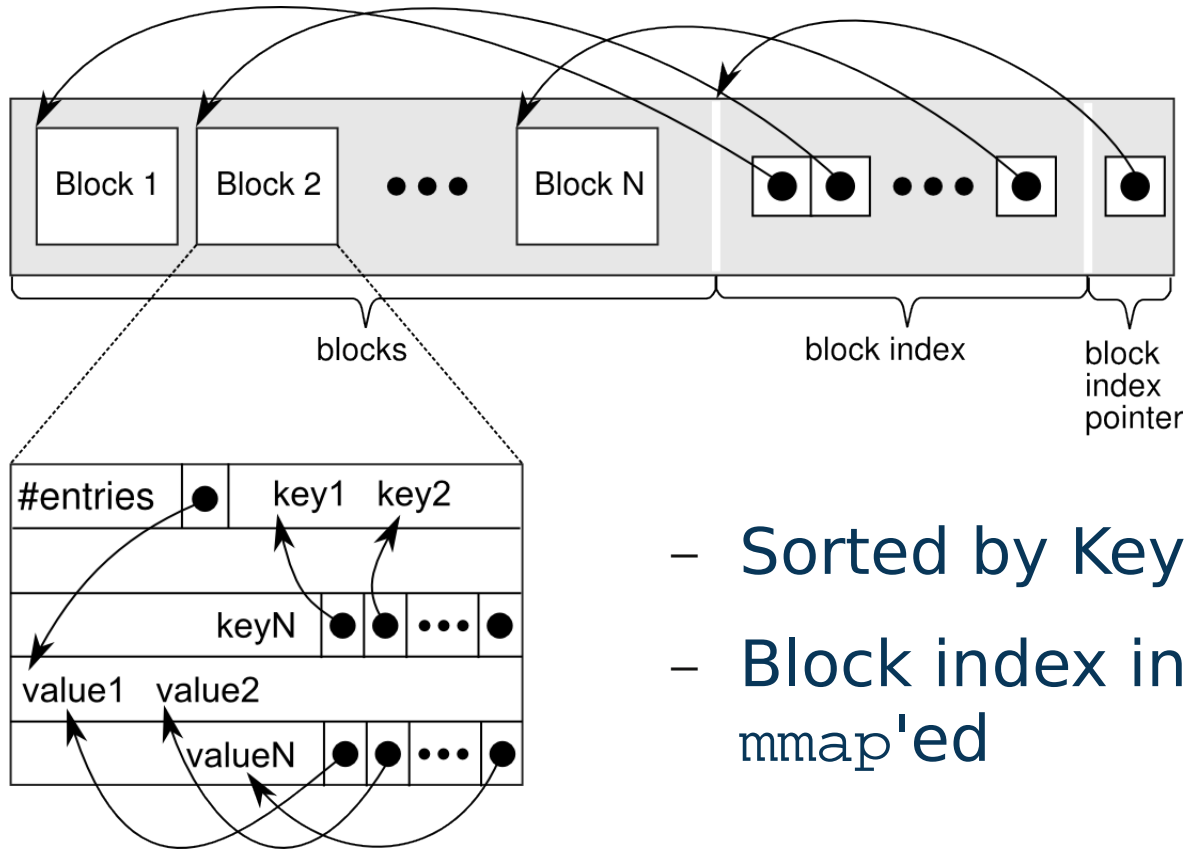


checkpoint()

1. create snapshot
2. rangeLookup(*) at snapshot 3, write new on-disk index
3. replace on-disk index, discard snapshots, purge log



On-disk Index



- Sorted by Keys
- Block index in RAM, blocks mmap'ed

BabuDB: Related Work

- Inspired by log-structured merge trees (LSM-trees)
 - Only one on-disk index
 - No „rolling merge“
- Made popular by Google Bigtable
 - Insert/lookup/merge similar as in Bigtable's Tablets

BabuDB: Metadata Mapping

- Mapping a hierarchical directory tree to a flat database index:

key:

parentID	filename	type
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value (type 1):

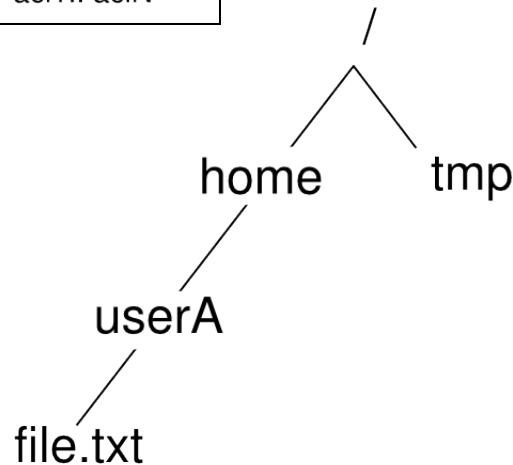
atime	ctime	mtime	file size
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value (type 2):

ID	content locations	perm.	owner	link count	...
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value (type 3):

xattr1.. xattrN	acl1.. aclN
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key	value
0,/,1	atime=2009-01-01 12:00 CET...
0,/,2	ID=1,perm=rwxr-x---...
0,/,3	empty
1,home,1	atime=2009-01-01 12:00 CET...
1,home,2	ID=2,perm=rwxr-x---...
1,home,3	empty
1,tmp,1	atime=2008-10-21 05:21 CET...
1,tmp,2	ID=3,perm=rwxrwx---...
1,tmp,3	empty
2,userA,1	atime=2009-01-01 12:00 CET...
2,userA,2	ID=4,perm=rwx-----...
2,userA,3	empty
4,file.txt,1	atime=2008-10-05 23:49 CET...
4,file.txt,2	ID=5,perm=rwx-----...
4,file.txt,3	empty

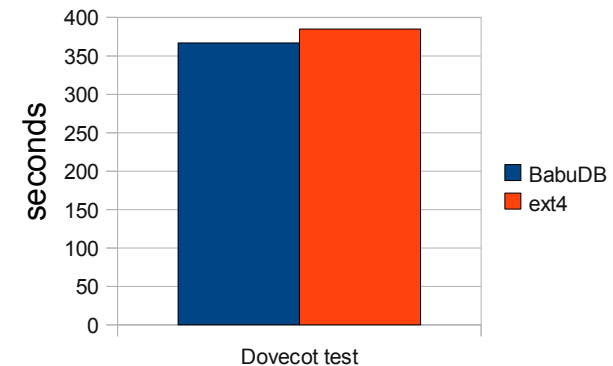
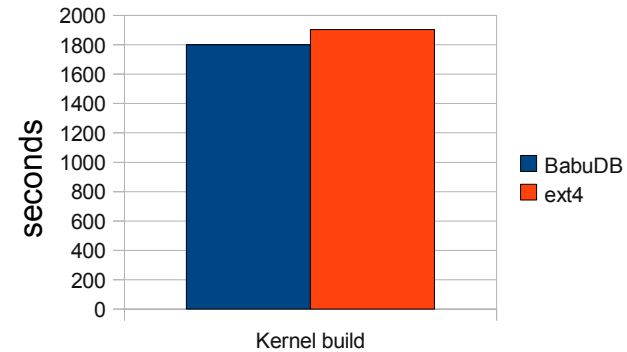
BabuDB: Advantages

- Why BabuDB for File System Metadata?
 - Short-lived files
 - 50% of all files deleted within 5 minutes
 - Atomic file system operations w/o locking or transactions
 - e.g. `rename`
 - Directory content in contiguous disk regions
 - Efficient `readdir + stat`
 - Snapshots
 - No need for multi-version data structures

BabuDB: Evaluation

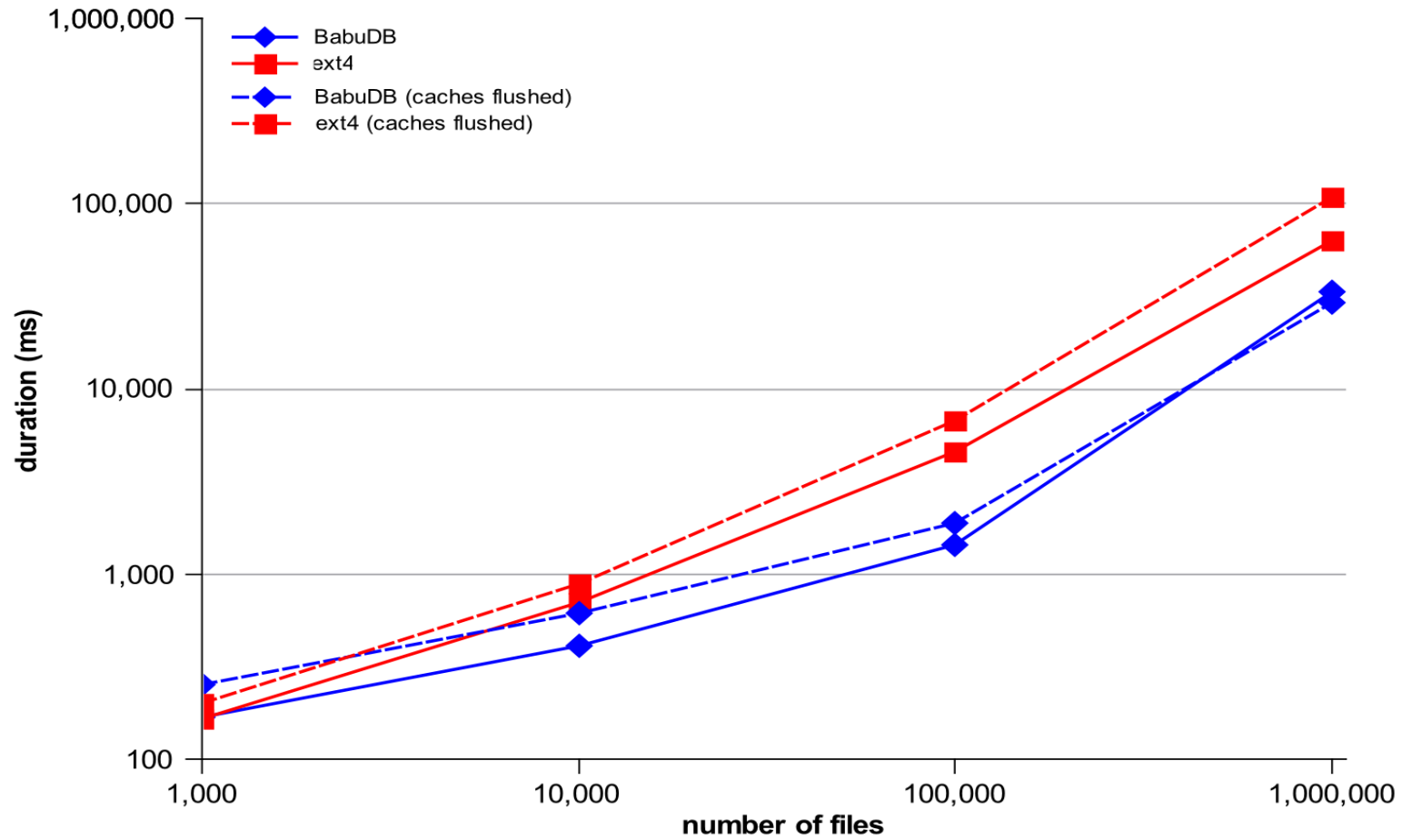
- Linux kernel build
 - ~10M calls: 44% stat, 40% open, 15% readlink, 1% others

- Dovecot mail server + `imaptest`
 - ~2M calls: 51% stat, 48% open, 1% others



BabuDB: Evaluation

– Listing directory content



Summary

- BabuDB is ...
 - an efficient key-value store
 - optimized for file system metadata but also suitable for other purposes
 - suitable for large-scale databases
 - available for Java and C++ under BSD license
 - used in the XtremFS metadata server



<http://babudb.googlecode.com>



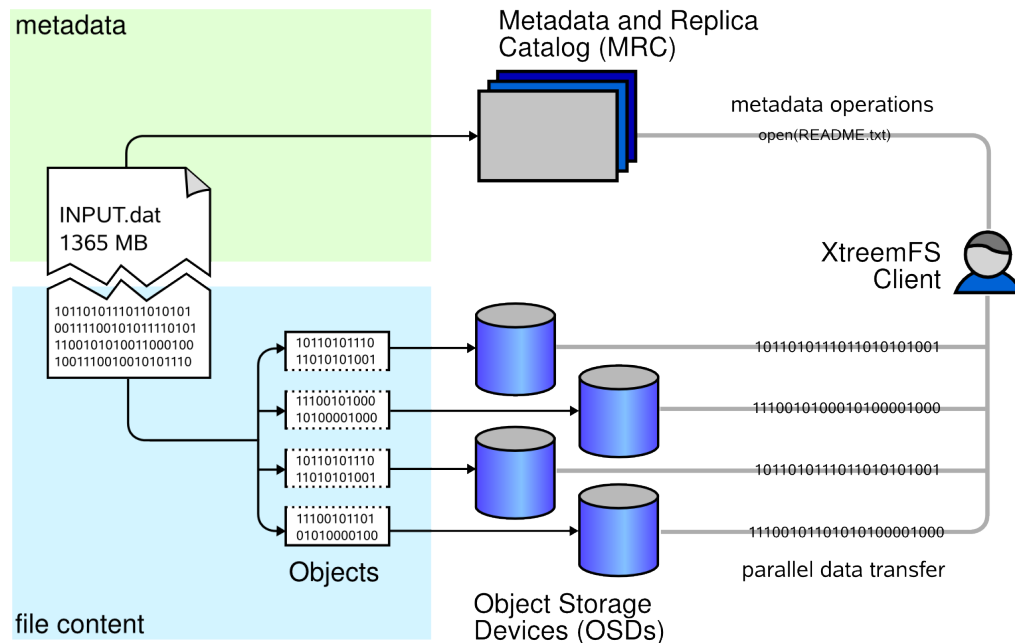
<http://www.xtreemfs.org>

Thank you for your attention!

Background: XtreamFS

– XtreamFS: a distributed replicated Internet file system

- part of the XtreamOS research project
- developed since 2006 by partners from Germany, Spain and Italy



– Object-based architecture:

- **MRC** stores metadata
- **OSDs** store pure file content as objects
- **Clients** provide POSIX file system interface

www.xtreemfs.org

The XtremOS Project

- Research project funded by the European Commission
- 19 partners from Europe and China
- XtremFS is the data management component
 - developed by ZIB, NEC HPC Europe, Barcelona Supercomputing Center and ICAR-CNR Italy
 - ~ 3 years of development
 - first public release in August 2008



XtreemFS: Overview

- What is XtreemFS?
 - a **distributed** and **replicated** **POSIX** compliant file system
 - **off-the-shelf** Servers - no expensive hardware
 - servers in **Java**, runs on Linux / OS X / Solaris
 - client in **C**, runs on Linux / OS X / Windows
 - **secure** (X.509 and SSL)
 - **easy** to install and maintain
 - open source (GPL)



File System Landscape

