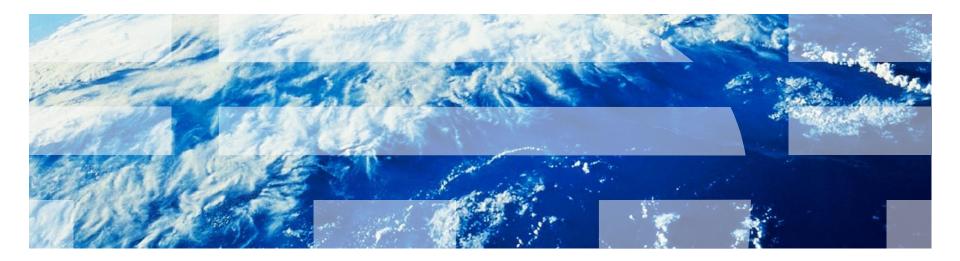


ZoneFS: Stripe Remodeling in Cloud Data Centers

Lanyue Lu - University of Wisconsin-Madison Dean Hildebrand, Renu Tewari - IBM Almaden Research Lab



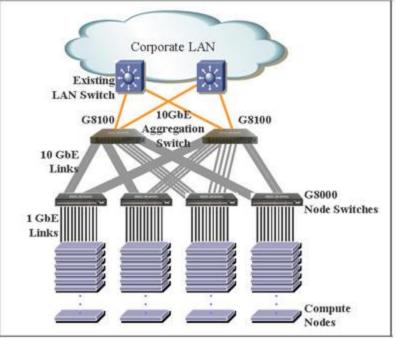


Cloud Data Centers on the Cheap

Network infrastructure

Hierarchical

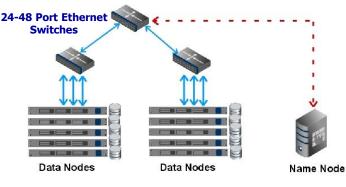
- 10GigE at root
- 1GigE at leaves (racks)
- Commodity switches
 - Limited switch port buffering
- Oversubscribed 8:1 (125Mbps) or more
- Bandwidth out of each rack is limited
 - Depends on oversubscription at connection level in hierarchy



*from Data Center Design by PTS Data Center Solutions, Inc.



Storage on the cheap (re-use existing local disks)

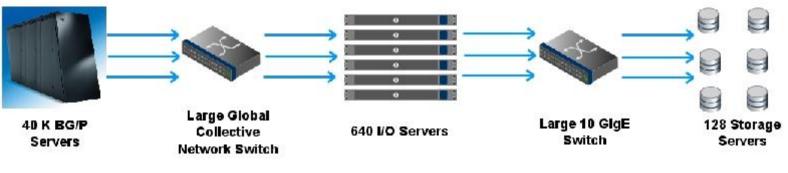


HDFS Architecture

- Storage software for cheap hardware
- o Internet-Scale File Systems (HDFS, GFS, Cloudstore)
 - Applications write all data to local node
 - Reduces parallelism, uneven load balancing
 - Create hotspots
 - Replication for fault tolerance
 - Triplication: 2 copies in 1 rack, 1 copy in another rack
 - Large stripe size (64 MB)
 - Non-standard file system API
 - Not POSIX API or semantics
 - Map/Reduce specific, cannot handle general workloads
 - Fixed 1:1 ratio of storage and compute nodes
 - Limited metadata processing capability
 - Limited data sharing
 - Low small I/O and create performance
 - Space allocation can be slow which space is limited



But Supercomputers do it (with money...)



Intrepid Supercomputer and GPFS Parallel File System

- Storage software for expensive hardware
- Parallel file systems are only as fast as the hardware
- Use large and expensive switches
 - Large port buffers allow a high-level of parallelism
 - \$700K for 128-port 10GigE switches (or >100 48-port GigE switches)
- o Compute servers realize all available bandwidth



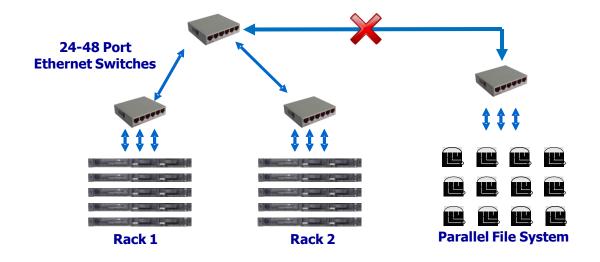
Merge Parallel File Systems and Cheap Hardware?

Goal: Leverage parallel file systems with data center infrastructures

But there are options.....



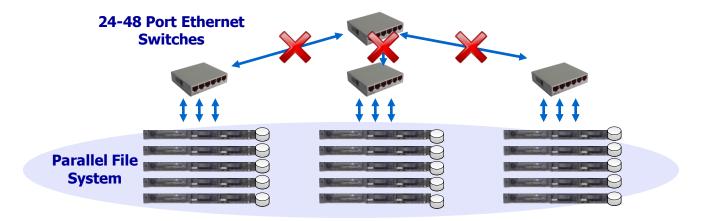
Data Centers and Parallel File Systems Option 1: Traditional (separate compute and I/O servers)



- Oversubscription creates bottlenecks
- o Compute servers realize a fraction of desired bandwidth
- Limited port buffering reduces I/O performance even further



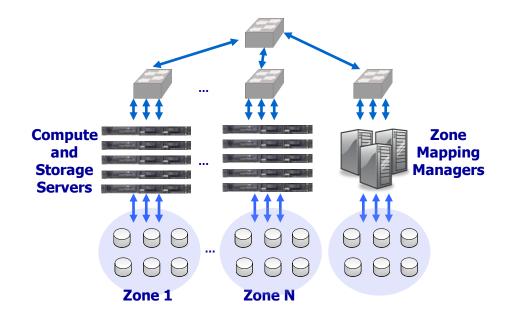
Data Centers and Parallel File Systems Option 2: Overlay compute and storage nodes



- Network oversubscription defeats purpose of data striping
- Aggregation and leaf switches scale in unison
 - 10GigE switches will emerge at leaf layer
 - 40/100 GigE switches will emerge at aggregation layer



ZoneFS Architecture



Re-organize file system into Zones

- Switch, compute nodes, storage nodes
- Flexible storage node disk architecture
 - Dedicated or non-dedicated
 - Local disk, SAN

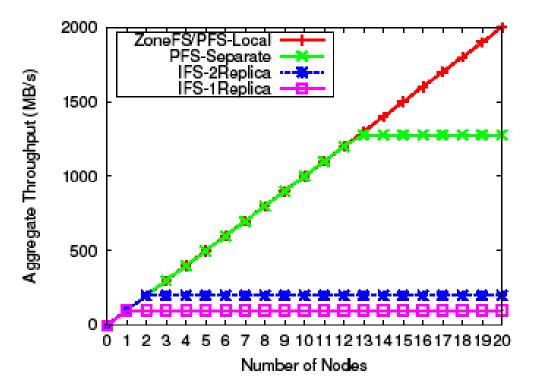
• Stripe data across storage nodes within a Zone

Level with the highest bandwidth

Compute nodes have parallel data access within their Zone



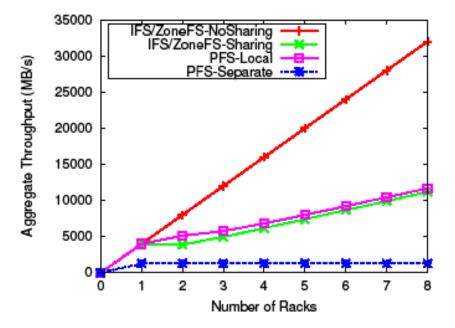
Single Rack Scalability Simulation Read single data chunk



- Internet scale file systems (IFS) limited by bandwidth to number of replicas
- Parallel File System (PFS-Separate) limited by 10GigE external switch
- ZoneFS and PFS-Local leverage full bandwidth of rack switch



Multi-Rack Scalability Simulation Access large file test



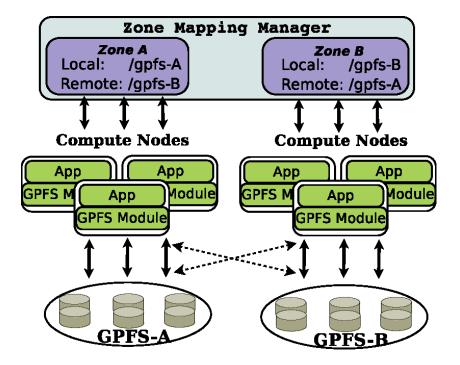
- Parallel File System (PFS-Separate) limited by 10GigE external switch
- {PFS-Local, IFS, ZoneFS}-Sharing all limited by oversubscribed switches
- {IFS, ZoneFS}-NoSharing scale linearly with additional rack switches

Notes:

- Sharing Access data across racks
- NoSharing Access data within rack
- 40 nodes per rack
- Simulated up to 48 racks



ZoneFS Prototype

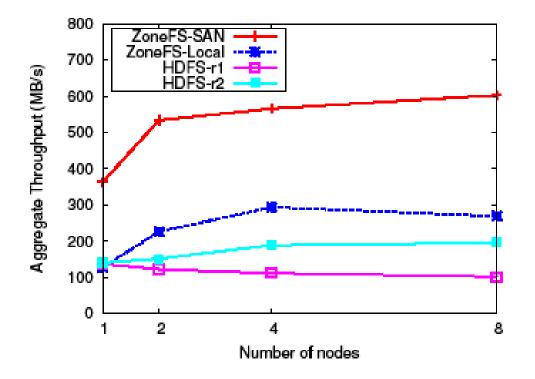


- Layered above GPFS in Linux 0
 - SAN and DAS
- File system per rack Ο
 - Inter-rack access available through GPFS client to remote file systems
- Job Scheduler
 - Hadoop or MPI

 - GPFS pick any random node
 ZoneFS pick random node in correct zone



Hotspot read performance



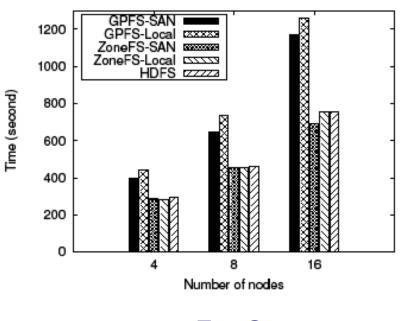
o Setup

Single rack

- All data on a single node
- HDFS constrained by number of replicas
- o ZoneFS-Local performance constrained by limited port-buffering in Netgear switch
- ZoneFS-SAN limited by SAN bandwidth

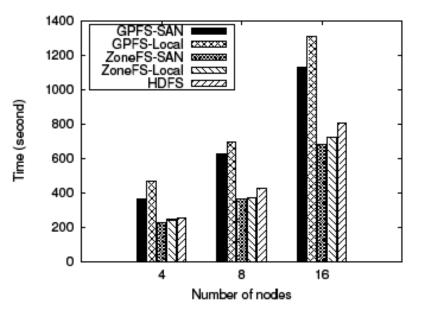


Hadoop Teragen and Terasort



TeraGen (write intensive)

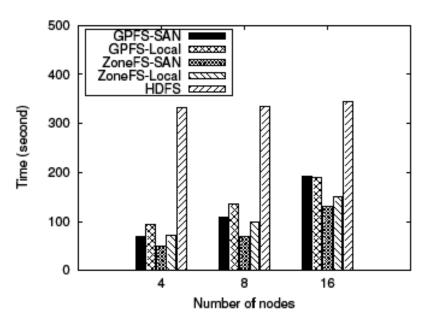
- o Setup
 - 2 racks
 - Increase nodes in-tandem between racks
 - IGB file per node
- GPFS limited by inter-rack switch



Terasort (read and write intensive)

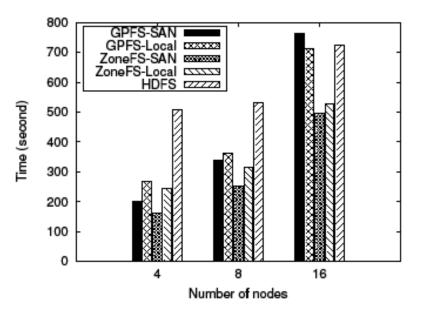


General workload - Image conversion



1MB JPG to 9MB BMP

- o Setup
 - Relatively small files
 - 200 images per node
 - 2 racks
- GPFS constrained by inter-rack switch
- HDFS constrained by single NameNode
- ZoneFS-SAN limited by disk bandwidth
- ZoneFS-local limited by switch port-buffering



4MB JPG to 35MB BMP



Summary

- Enterprise solution
 - Support general applications (POSIX API and semantics)
 - Support Information Lifecycle Management (ILM)
- Scalability and Performance
 - Independent scaling of switch, compute, and storage nodes
 - Add zones as needed
 - Intra-switch parallelism improves load balancing
 - Avoid the inter-rack network bottleneck experience by parallel file systems
 - Good performance for both analytics and general workloads
- Availability and flexibility
 - Avoid ingesting input data and offloading results
 - Single storage system for primary and analytic data
 - Flexible options for compute and storage nodes
 - Local disks or SAN per rack
 - Vary the ratio of compute to storage nodes
 - Independent management of compute and storage node



Thank You !

Questions ?

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