Exa-Scale FSIO
Can we get there?
Can we afford to?

05/2011 Gary Grider, LANL

LA-UR 10-04611

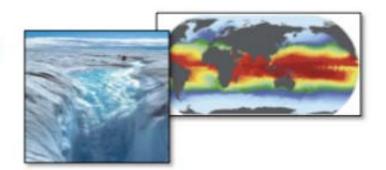
# Pop Quiz: How Old is this Guy?



## Back to Exa-FSIO

#### Mission Drivers

- Climate Change: Understanding, mitigating and adapting to the effects of global warming
  - Sea level rise
  - Severe weather
  - Regional climate change
  - Geologic carbon sequestration
- Energy: Reducing U.S. reliance on foreign energy sources and reducing the carbon footprint of energy production
  - Reducing time and cost of reactor design and deployment
  - Improving the efficiency of combustion energy sources
- National Nuclear Security: Maintaining a safe, secure and reliable nuclear stockpile
  - Stockpile certification
  - Predictive scientific challenges
  - Real-time evaluation of urban nuclear detonation





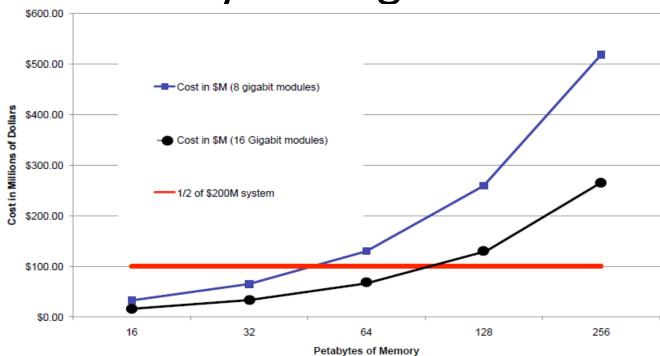


Accomplishing these missions requires exascale resources.

## Power is a Driving Issue

- Power per flop
- Power per byte
- Power per byte/sec
- Power for infrastructure
- POWER POWER POWER

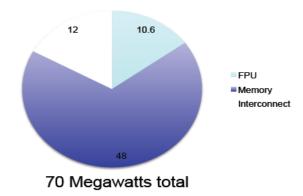
## Memory is a Big Problem

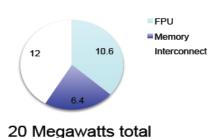


 Power Consumption with standard Technology Roadmap  Power Consumption with Investment in Advanced Memory Technology

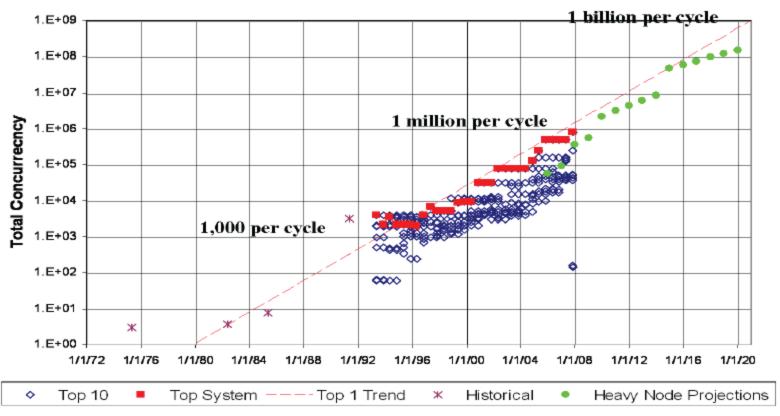
#### **Power**

Cost





### Parallelism will be Massive



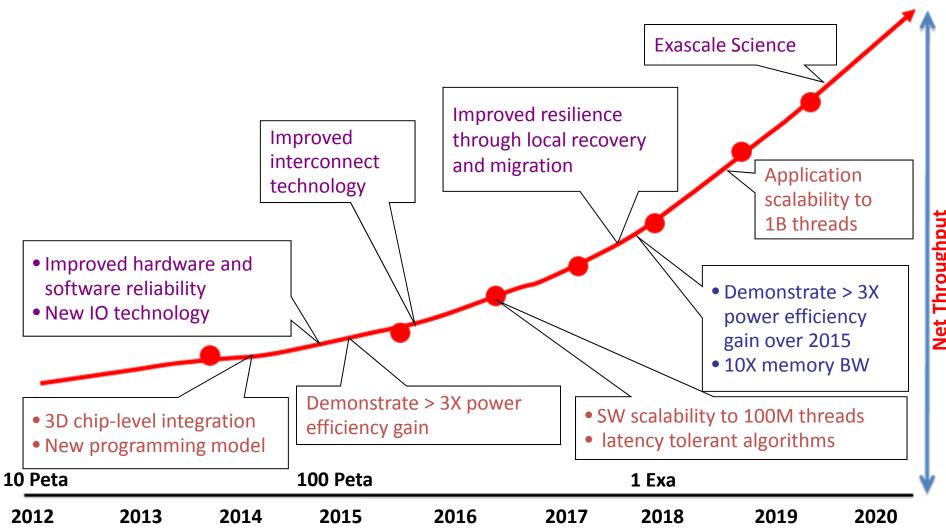
How much parallelism must be handled by the program?

From Peter Kogge (on behalf of Exascale Working Group), "Architectural Challenges at the Exascale Frontier", June 20, 2008

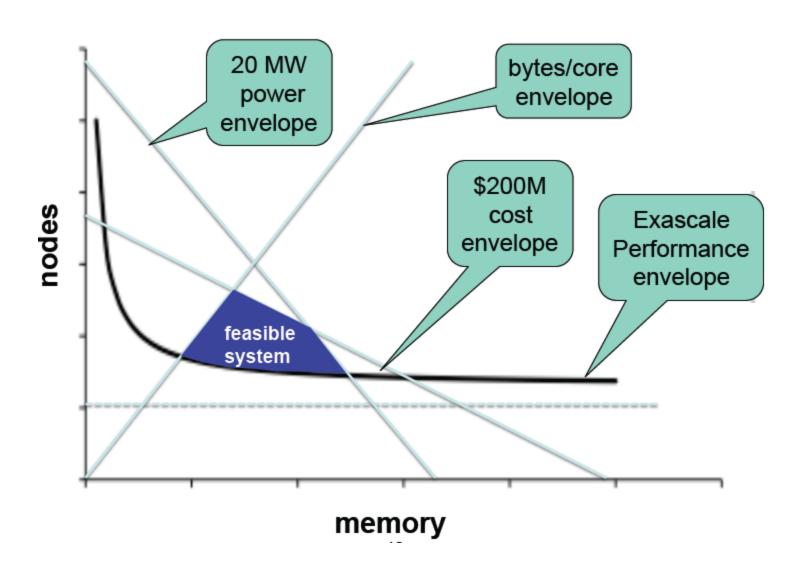
Need 1Million-way parallelism to reach an Exaflop . . .

And possibly another 100x just to hide latency

# Technology Roadmap



# It is a Complicated Trade Space



# Reliability will be Difficult

- Industry must maintain constant FIT rate per node
  - ~1000 failures in time
- Moore's law gets us 100x improvement
  - But still have to increase node count by 10x
- So we will own 10x worse FIT rate
  - MTTI 1week to 1 day
  - MTTI 1 day to 1 hour

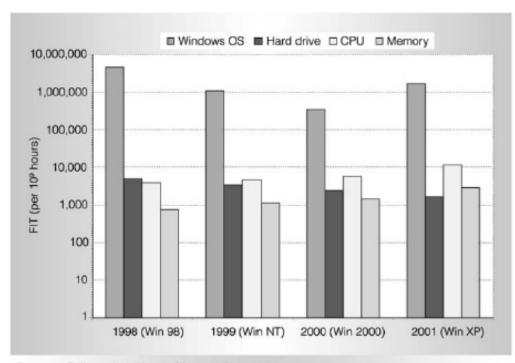


Figure 2. Failures in billions of hours of operation.2-5



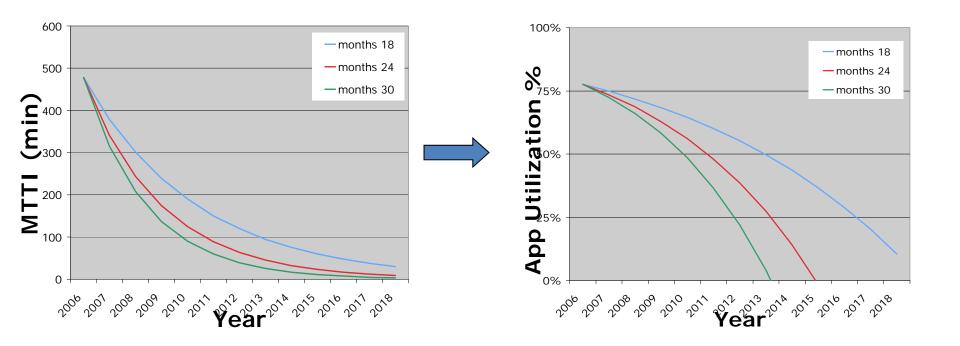
#### **Potential System Architecture Targets**

| System attributes             | 2010     | "2015"           |          | "2018"        |           |
|-------------------------------|----------|------------------|----------|---------------|-----------|
| System peak                   | 2 Peta   | 200 Petaflop/sec |          | 1 Exaflop/sec |           |
| Power                         | 6 MW     | 15 MW            |          | 20 MW         |           |
| System memory                 | 0.3 PB   | 5 PB             |          | 32-64 PB      |           |
| Node performance              | 125 GF   | 0.5 TF           | 7 TF     | 1 TF          | 10 TF     |
| Node memory BW                | 25 GB/s  | 0.1 TB/sec       | 1 TB/sec | 0.4 TB/sec    | 4 TB/sec  |
| Node concurrency              | 12       | O(100)           | O(1,000) | O(1,000)      | O(10,000) |
| System size (nodes)           | 18,700   | 50,000           | 5,000    | 1,000,000     | 100,000   |
| Total Node<br>Interconnect BW | 1.5 GB/s | 20 GB/sec        |          | 200 GB/sec    |           |
| MTTI                          | days     | O(1day)          |          | O(1 day)      |           |

### Gloom and Doom from 2006

- Petascale computing is coming
  - Orders of magnitude more components
  - Orders of magnitude more failures

#### Need raw data for better understanding of failures



## Past and Future Assumptions

#### Past

- All disk
- Constant ratio of total \$ to IO infra \$
- Machines wont accelerate their reliability per flop

#### Future

- Not necessarily all disk
- Not necessarily same % but close
- Machines may make accelerate progress on reliability/flop due to integration and industry desire to have constant reliability per socket

#### Can we do defensive IO at Exascale?

- If we loosen assumptions?
- If we can do it can we afford to do it?

## **New Assumptions**

| Year                 | EF  | 2010     | 2012   | 2014    | 2016    | 2018    |
|----------------------|-----|----------|--------|---------|---------|---------|
|                      |     |          |        |         |         |         |
| PF                   |     | 1.000    | 20.00  | 200.00  | 400.00  | 1000.00 |
| mem low PB           |     | 0.004    | 0.07   | 0.72    | 1.44    | 3.60    |
| mem med PB           |     | 0.020    | 0.40   | 4.00    | 8.00    | 20.00   |
| mem high PB          |     | 0.300    | 6.00   | 60.00   | 120.00  | 300.00  |
| Num Full Mem Cap     |     | 30       | 30     | 30      | 30      | 30      |
| Size Scratch PB low  |     | 0.108    | 2.16   | 21.60   | 43.20   | 108.00  |
| Size Scratch PB med  |     | 0.600    | 12.00  | 120.00  | 240.00  | 600.00  |
| Size Scratch PB high |     | 9.000    | 180.00 | 1800.00 | 3600.00 | 9000.00 |
| Time to dump Secs    |     | 1200.000 | 800.00 | 600.00  | 400.00  | 300.00  |
| Ckpt BW low TB/s     |     | 0.003    | 0.09   | 1.20    | 3.60    | 12.00   |
| Ckpt BW med TB/s     |     | 0.017    | 0.50   | 6.67    | 20.00   | 66.67   |
| Ckpt BW high TB/s    |     | 0.250    | 7.50   | 100.00  | 300.00  | 1000.00 |
|                      |     |          |        |         |         |         |
| Disk Capacity TB     |     | 2.000    | 3.92   | 7.68    | 15.06   | 29.52   |
| Disk Speed MB/s      | 100 | 100.000  | 140.00 | 196.00  | 274.40  | 384.16  |
| IO node thrput GB/s  | 100 | 1.000    | 2.000  | 4.000   | 8.000   | 16.000  |

#### Based On

- for machine sizes, mtti, etc. except 20 PB med mem machine and 30 dumps in scratch
- Seagate Disk Capacity/Size/ Pricing/Power (not shown)
- Micron Flash Capacity/Size/ Pricing/Power (not shown)
- 10% of mtti as dump time

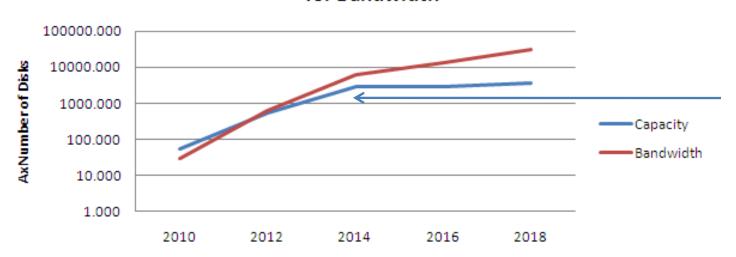
I wanted to know – what miracles will we need and to get past what problems.

# Status Quo: Use Disk Based Shared Global Parallel File System to Provide Dump Space

Disks Needed by Year for Low Mem Option

Notice Crossover Now We Buy for Capacity Soon We Will Buy

for Bandwidth



Notice that using these modeling parameters, we finally reach the predicted cross over point of buying disk for BW and not Capacity in 2012

2018 medium memory machine

- •4166 IO nodes, 175k disks
- •File System sees 50-100k way parallelism (assumes IOFSL)
- •\$225M pessimistic purchase (assumes no technologies pushing disk other than Flash)

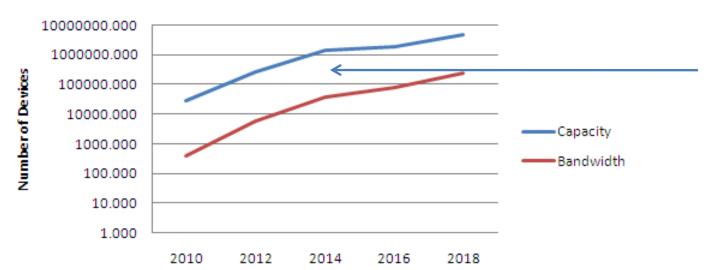
Power 1.5MWatts

Miracle Needed!

Buying disk for capacity is reasonably priced but buying disk for bandwidth gets expensive fast!

# Use MLC Based Shared Global Parallel File System to Provide Dump Space

Devices Needed by Year for Low Mem Option Notice You Always Need More Devices for Capacity



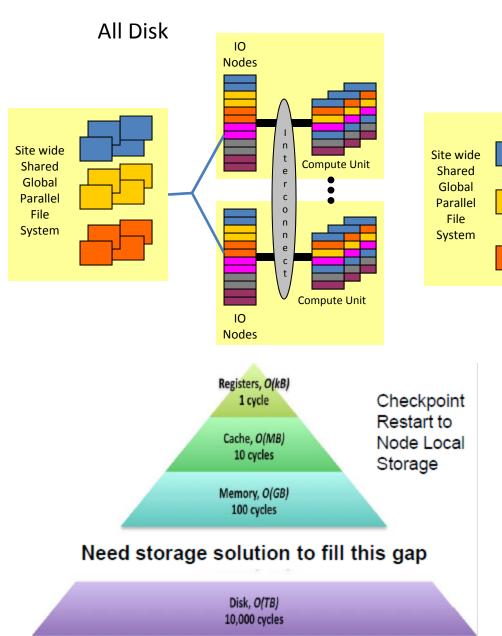
Notice that buying MLC for capacity is expensive but buying it for Bandwidth is cheaper

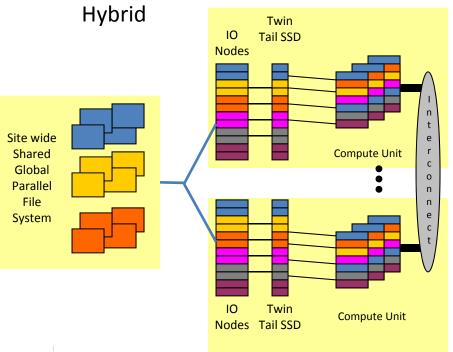
2018 medium memory machine

- •4166 IO nodes
- •File System sees 50-100k way parallelism (assumes IOFSL)
- •\$625M pessimistic purchase (assumes no technologies pushing disk other than Flash)

  Miracle Needed!
- Power 2.5MWatts (have to buy so much to get capacity)

#### **Hybrid Disk (Capacity)/SSD (Bandwidth)**

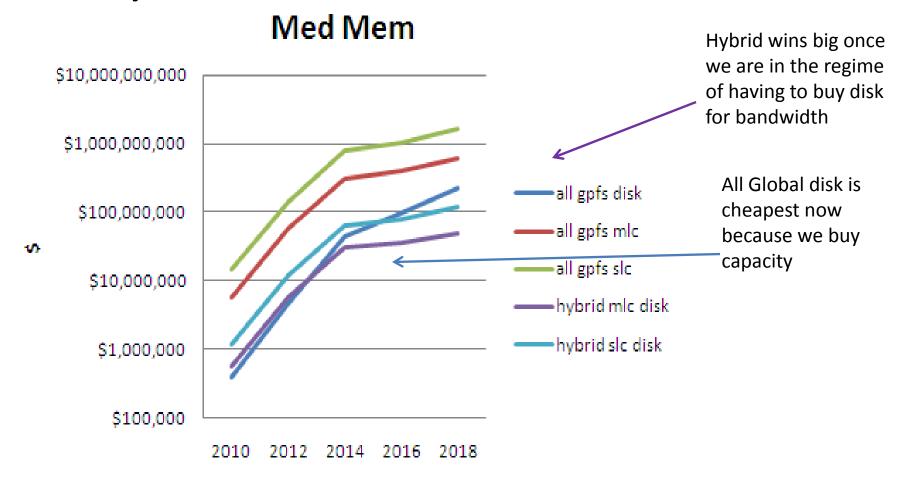




# Must move checkpoint device closer to compute memory

- on node has jitter issues
- at least near node is required
- Leads to Hybrid Storage model

## Hybrid MLC burst / Disk Global

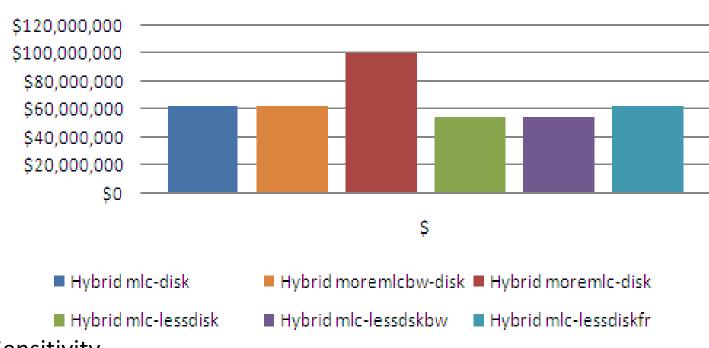


#### 2018 med mem mach

- •416 IO Nodes, 20k disks not much of a stretch
- Disk FS sees modest parallelism assumes IOFSL/burstbuffer etc.)
- •\$60M pessimistic purchase worst case (all migrated to disk and tech price)
- Power 2.2MWatts

## Hybrid MLC burst / Disk Global

#### First Order 2018 Med Memory Sensitivity Analysis



**Cost Driver Sensitivity** 

- More MLC BW (free − capacity driver)
- More MLC Cap (costly capacity driver)
- •Less Disk Cap (small savings (MLC capacity driver)
- Less Disk BW (small savings controllers/ION etc. ( MLC capacity driver)
- •Less Frequent MLC to Disk (no savings, Disk Capacity Driver)

## A Feasible Evolutionary Approach?

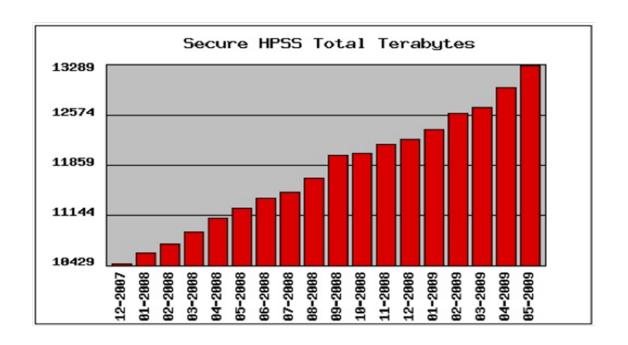
| Summary: Issue   | Action  |
|--|---|
| Probably pretty close on storage densities, bandwidths, and cost in fact it may be a bit conservative ( maybe more than a bit )                                  | ts, Continue to update model  |
| Based heavily on MTTI assumptions in the DARPA study and that study indicates a pretty large per socket improvement in MTTI without good substantiation          | t Get serious about measuring and predicting this!                              |
| Assumes that existing techniques like RAID or other redundant techniques will keep the burst buffer working often enough to n have issues without substantiation | Keep our eye on Flash<br>not reliability – prospects are<br>good given wide use |
| Assumes existing RAS techniques for file systems will be able to keep up without substantiation  | Keep our eye on this  |
| Have to have burst buffer so we will need software to manage MLC burst buffer, with bleed to global disk   | SCR LLNL / PLFS LANL / ADIOS ORNL / MPI-IO ANL. Zest PSC,                       |
| Assumes flattening to get high % of peak on disks (like log structure)   | PLFS LANL / ADIOS ORNL / MPI-IO ANL, Zest PSC,                                  |
| Need a way to deal with large numbers of files   | Giga+, etc.   |

Maybe we can get to Exascale with evolution only, but it would be pretty sad if we didn't also attempt some more fundamental revolutionary approaches!

We need both an evolutionary track and a revolutionary track!

# **Archive Analysis**

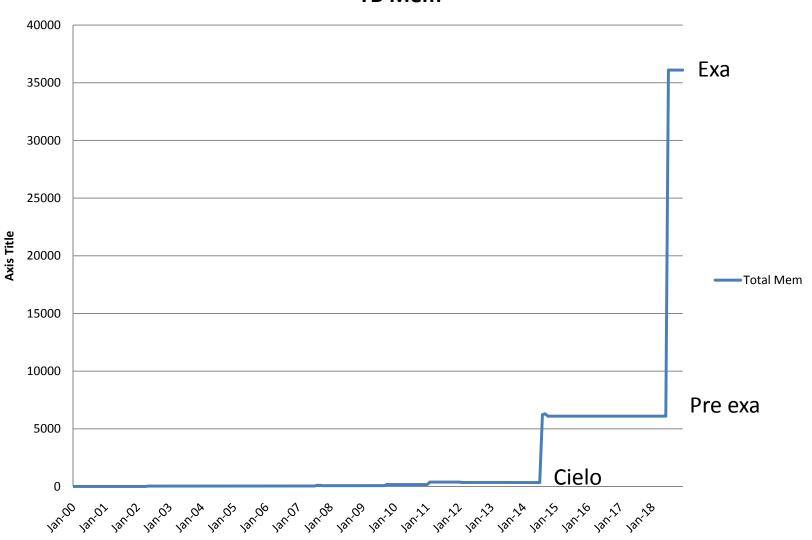
#### Can we Afford an Archive?



- Unlimited archive will become cost prohibitive
- Past method of using bandwidth to archive as rate limiter may not be adequate going forward

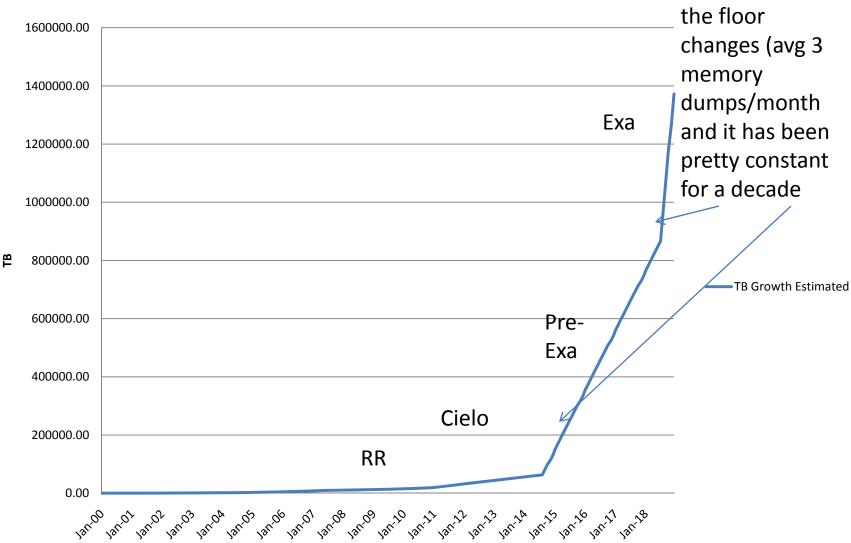
# Archive Growth Depends on TB of Memory on the Compute Floor





#### **Archive Growth TB**

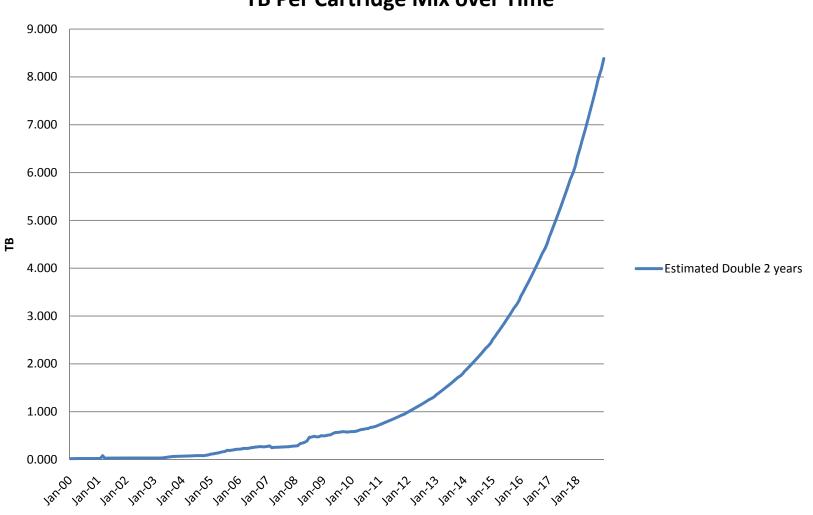




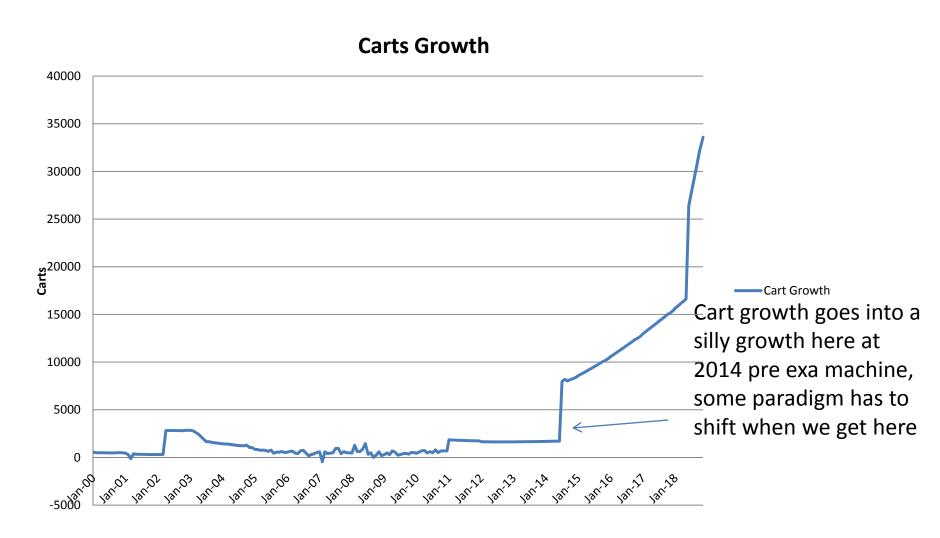
Notice the slope changes when the memory on

# Effective Cartridge Density Considering 3 Generations of Technology

#### **TB Per Cartridge Mix over Time**



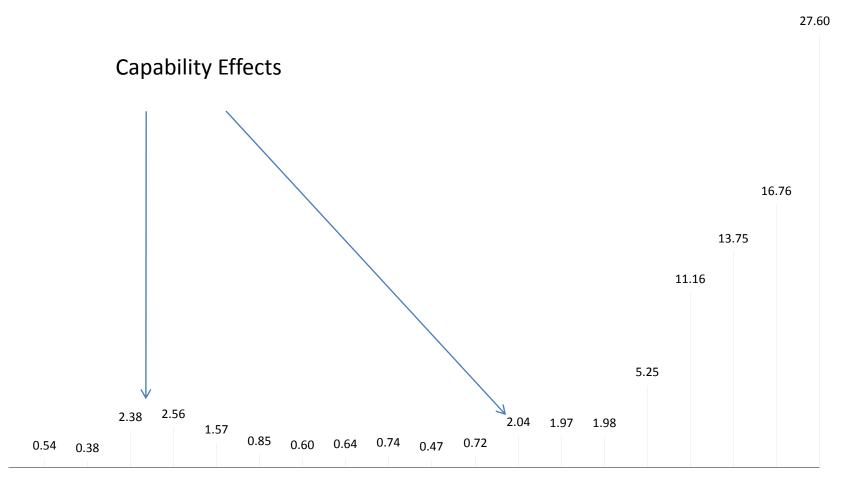
# Cartridge Growth (new data and shrink data on latest cart tech)



# Yearly \$ on Carts

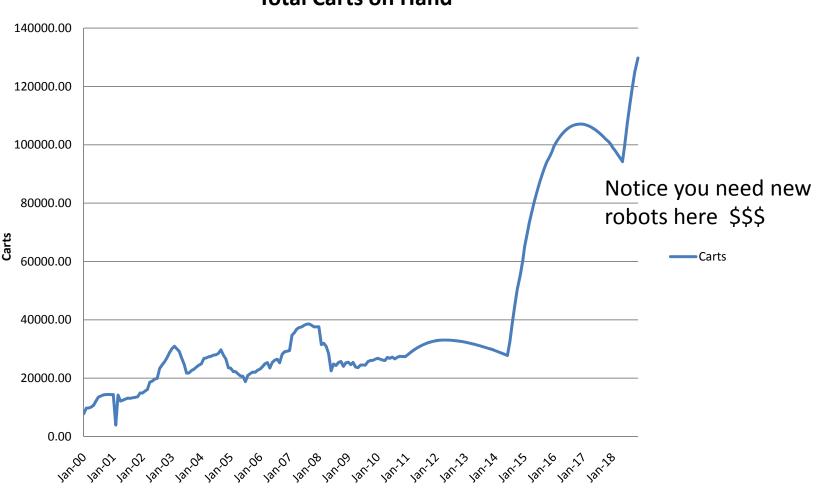
#### **Yearly \$M on Carts**

■\$ for new Carts



#### **Total Carts on Hand**

#### **Total Carts on Hand**

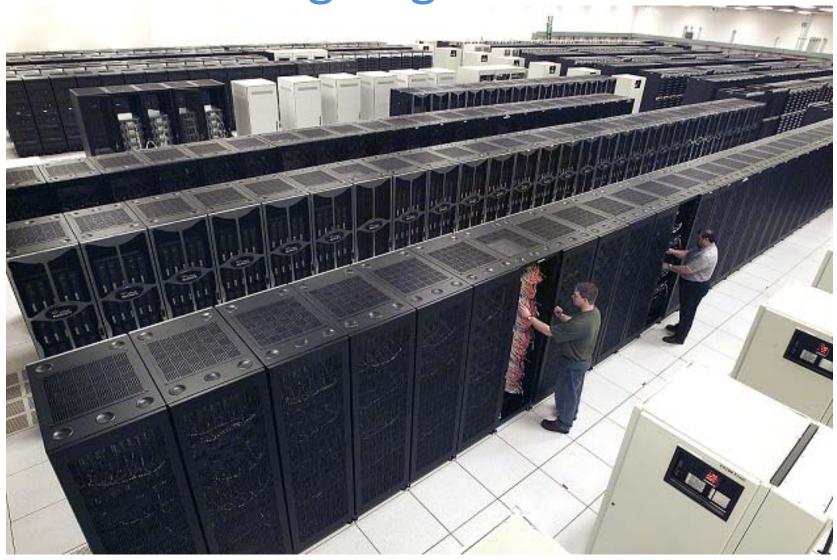


# PRObE

An NSF large scale systems research center in Los Alamos, New Mexico

http://www.newmexicoconsortium.org/probe

# LANL was going to trash this!



#### PRObE to the rescue!

- NSF Funds the New Mexico Consortium (NMC) to bring LANL supercomputers back to life
- PRObE –

Parallel Reconfigurable

**Observational Environment** 

#### Motivation

- Systems research community lacks very *large* dedicated resource for experiments, fault injection, and hardware control.
- Research on large compute resources often constrained by imposed software stack
- Large systems are hurried through testing phase into production. Inhibits systems research at scale.
- And...

### What is PRObE?

- Low level systems research center
- Days to weeks of dedicated usage of a large computer resource for projects
  - Physical and remote access
- Complete control of hardware and software
- Enables fault injection and failure statistics collection
- End-of-life destructive testing
- Supports parallel and data intensive workloads

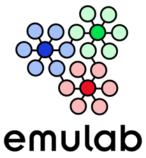
# Brought to you by:







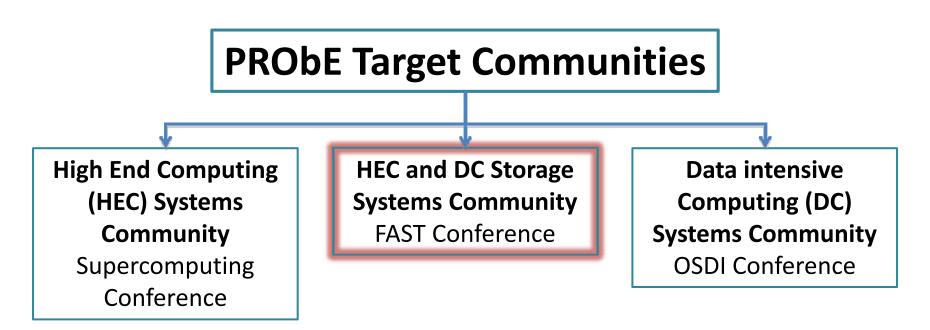
EST.1943 -





# For Systems Research Users

- NFS's "who can apply" rules
  - Includes international and corporate research projects (partnership with US university preferred)



# **PRObE Decision Making**

3 Committees, members selected from community



Government (1), Industry (1), and University Top
Systems Researchers (6)

# PROJECT SELECTION COMMITTEE

University Top Systems Researchers

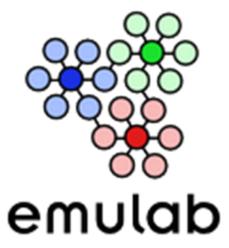
MANAGEMENT GROUP

**University Top Systems Researchers** 

USER ENVIRONMENT
COMMITTEE

### Software

- First, "none" is allowed
  - Researchers can put any software they want onto the clusters
  - Full OpenCirrus stack possible
- Second, a well known tool managing clusters of hardware for research
  - Emulab (www.emulab.org), Flux Group, U. Utah
  - Widely used in academic systems community
  - Enhanced for PRObE hardware, scale, networks, resource partitioning policies, remote power and console, failure injection, deep instrumentation



### Cluster Installation Timeline

| When      | Nodes   | Core<br>s | Purpose / Type                    | Where | Name      |
|-----------|---|-----------|-----------------------------------|-------|-----------|
| Q1 CY2011 | 128   | 256       | Front end test cluster (IB)       | CMU   | Marmot    |
| Q3 CY2011 | 128   | 256       | Front end (Myri)                  | NMC   | Denali    |
| Q3 CY2011 | 36  | 1728      | High core count cluster (IB)      | CMU   | Susitna   |
| Q4 CY2011 | 1024  | 2048      | High node count cluster (Myri)    | NMC   | Sitka     |
| Q1 CY2012 | 1024  | 2048      | High node count cluster (IB)      | NMC   | Kodiak    |
| Q3 CY2013 | 16  | 128       | Front end (IB)                    | NMC   | Yakutat   |
| Q3 CY2013 | 200   | 1600      | High node count cluster (IB)      | NMC   | Nome      |
| Q4 CY2013 | 36  | 3456      | High core count cluster (100GigE) | CMU   | Matanuska |
| Q2 CY2014 | Q2 CY2014 Next high node count cluster identified and |           |                                   |       |           |

..first 1024 node cluster decommissioned to make room for next large cluster. Research contest to see how best to torture the machine on its way out will be conducted.

#### **Contacts**

- Website
  - http://www.newmexicoconsortium.org/probe
    - Will soon house: Wiki's, Published data
       Committee Nomination & Election pages
- Email
  - probe@newmexicoconsortium.org