# Virtualization Aware Access Control for Multitenant Filesystems

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# Storage Consolidation

#### Benefits

- Efficient usage of storage resources
- Sharing support
- Increased manageability
- Reduced cost

#### Block-level storage

• Direct virtual disk access through block interface

### File-level storage

• Direct filesystem sharing through file interface

# Storage Multitenancy

#### Goal

• Storage infrastructure shared among different tenants

### Requirements

- Scalability: Support enormous number of end users
- Isolation: Isolate the user identities and access control of different tenants
- Sharing: Flexible data sharing within or between tenants
- **Compatibility:** Compatibility with existing applications
- Manageability: Flexible resource management

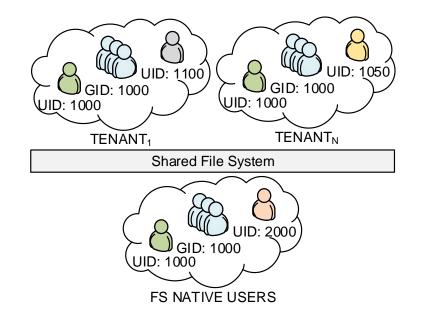
### Research focus

• Efficient and secure multitenancy in VM filesystems

# Motivation

### Problem of multitenancy

- Shared FS namespace
- Crosstalk between tenants
- Complicated security



#### Native multitenancy at the filesystem level

- Clean way to isolate multiple tenants
- Shared hardware, operating system, fileservers
- Configurable isolation, sharing, performance, manageability

# Prior approaches

#### Centralized

• The principals' identities of all tenants centrally maintained

- Poor scalability, isolation and manageability

#### Peer-to-peer

- The principals of each tenant managed locally
- Tenants communicate to publicize their principals' identities
- Overhead to periodically synchronize the tenants

### Mapping

• Local principal IDs mapped to global unique IDs

- Mapping overhead, sharing complications, security violations

# The Dike Approach

#### Hierarchical identification and authentication

- The tenants manage their principals
- The provider manages the tenants

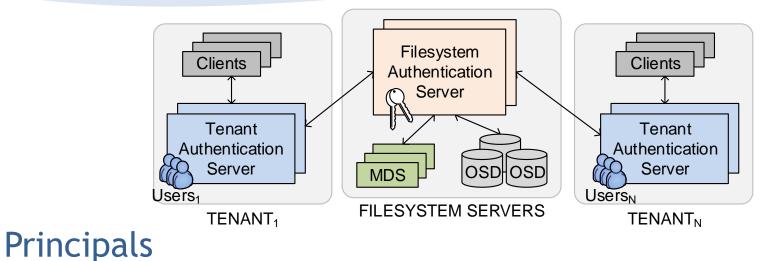
#### Native multitenant authorization

- Separate ACLs per tenant and provider
- Namespace isolation through filesystem views

#### Efficient permission management and storage

- Shared common permissions
- Inheritance of permissions

# Identification



- Tenant principals: Use/manage tenant resources
- Native FS principals: Manage the FS
- Tenant Authentication Server (TAS)
  - Certifies local clients and principals

### Filesystem Authentication Server (FAS)

• Certifies filesystem services, tenants, native principals

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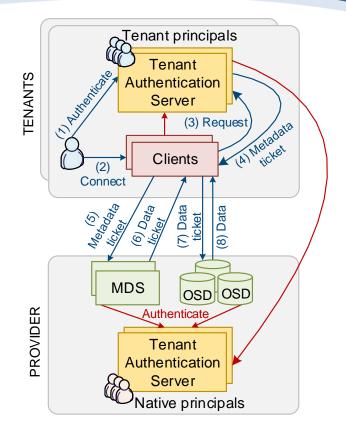
# Authentication

#### Metadata ticket

- Securely specifies tenant principal
- Provides access to MDS

### Data ticket

- Securely specifies tenant principal and permissions
- Provides access to OSDs



(1) Principal authenticated by TAS
(2) Principal requests FS access
(3) Client contacts TAS

(4) Client receives Metadata ticket

- (5) Client contacts MDS
- (6) MDS issues Data ticket
- (7) Client contacts OSD
- (8) Client accesses data

# Authorization

#### Access control isolation

- Separate ACLs per tenant, provider
- Metadata accessible through views

#### Filesystem view

 Used by native principals to manage tenants

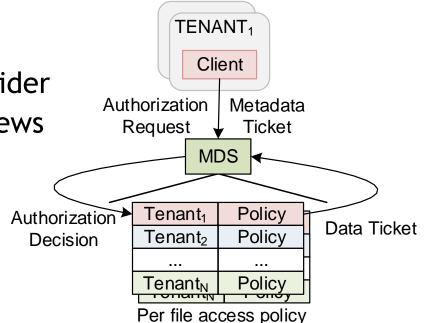
#### Tenant view

Used by tenants to access or manage tenant resources

### File sharing

- Private to a principal
- Shared across principals of one or more tenants





# **Common Permissions**

#### Goal

• Reduce filesystem load by reducing ACLs

### Per tenant permission inheritance

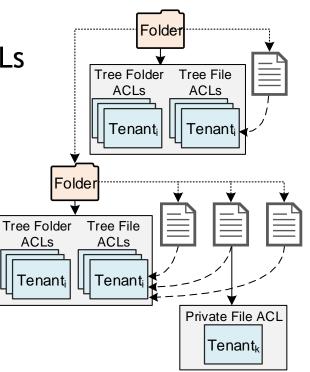
 Permissions can be inherited to child files/folders

#### Per tenant common permissions

• Child files can share parent's ACL

#### ACLs

- Tree folder ACLs: Folder permissions
- Tree file ACLs: Shared child files permissions
- Private file ACLs: Child file permissions explicitly set by user



# Security Analysis

#### Captured credential

• Fresh tamperproof credentials cannot be forged

#### Compromised tenant principal account

- Compromised tenant view is isolated
- Attack limited to principal's private or shared files
- Cross-tenant policy violation is prohibited

#### Attack by revoked tenant

- Restricted through deleted tenant view
- Tenant cannot access other views

#### Compromised provider administrator account

• Handle via good practices (e.g., restricted remote access)

# Prototype

#### Session

- Tenant identified by client
- Session limited to one tenant

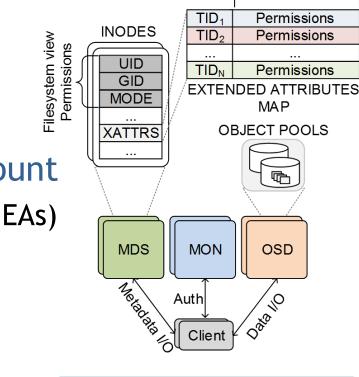
### Compromised tenant principal account

- **Tenant view:** Extended Attributes (EAs)
- Filesystem view: Regular fields
- EAs with tenant permissions not directly accessed by clients

### Capabilities

- Include principal and tenant identifiers
- Sent to clients with tenant file access





Tenant view Permissions

# **Experimentation Environment**

#### Configuration: AWS EC2 Instances

- **m1.xlarge:** x3, 4 VCPU, 15 GB RAM, Linux 3.9.3
- **t1.micro:** x32, 1 VCPU, 615 MB RAM, Linux 3.9.3

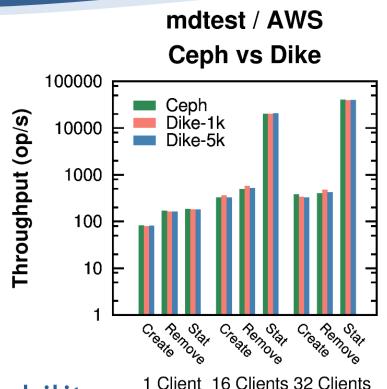
#### Filesystem configuration

- Ceph/Dike: m1.xlarge, 1xOSD+MON, 1xOSD+MDS, 1xOSD
- Gluster/Heka: m1.xlarge, 3 fileservers
- Replication factor 3

#### Microbenchmark

- mdtest
- 48000 created files and folders

Results



### Dike Client scalability

• 1  $\rightarrow$  32 clients: Similar to Ceph

### Dike Tenant scalability

•  $1k \rightarrow 5k$  tenants: 2% extra overhead

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# Results

#### Dike native multitenancy

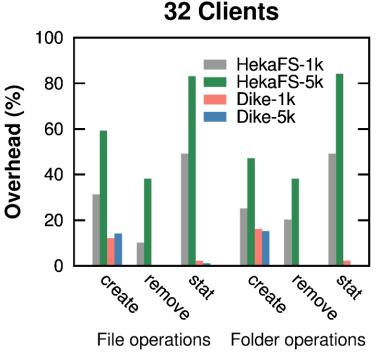
Limited overhead Scalable to thousands tenants

### Dike limited overhead

- 1k tenants overhead: up to 14%
- 5k tenants overhead: up to 16%

### ID mapping multitenancy too costly

- 1k tenants overhead: up to 49%
- 5k tenants overhead: up to 84%



mdtest / AWS

# Conclusions

#### Native filesystem multitenancy with sharing support

- Hierarchical identification scheme
- Namespace isolation: Per tenant and provider ACLs
- Per tenant common permissions and inheritance

#### Performance and security analysis

- Limited multitenancy overhead up to 16%
- Dike scalable to several thousand tenants
- Tenant principals not able to violate cross-tenant policy

#### Future work

- I/O intensive application experimentation
- Weaker trust assumptions

Backup

# Comparison of Interfaces

Benefits	Block Interface	File Interface
Compatibility	$\checkmark$	
Isolation	$\checkmark$	$\checkmark$
Sharing		$\checkmark$
Consistency, Performance		$\checkmark$
Disaster recovery, migration	$\checkmark$	$\checkmark$
Thin Provisioning	$\checkmark$	$\checkmark$
Searchability		$\checkmark$
Snapshoting, Versioning	$\checkmark$	$\checkmark$

# Dike compared to Ceph

#### Pool-level multitenancy

- Objects organized in per tenant pools
- No support for sharing files among tenants
- In Dike tenants can securely share the same pool

#### Centralized Identity management

- Keystone integration
- Poor scalability, isolation and manageability
- Dike: Hierarchical identity management scheme

#### ACLs

- Earlier versions of Ceph support Posix ACLs
- Single ACL for all tenants leads to poor isolation
- Dike: Separate ACLs per tenant and provider

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