Classifying Data to Reduce Data Movement in Shingled Write Disks

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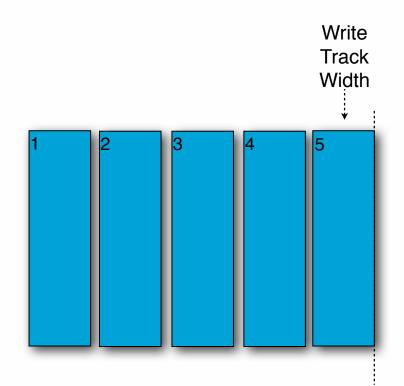


* Layers tracks like shingles on a roof

 Takes advantage of the fact that the read head is smaller than the write head

Problems

- A write can destroy data on subsequent tracks
- No more random writes and in-place updates



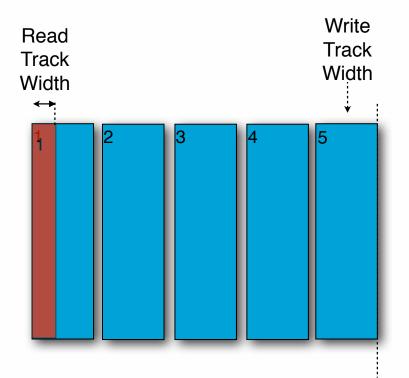


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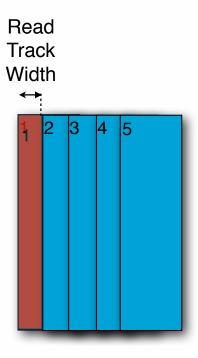
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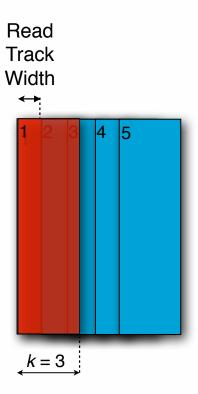
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- * Band compaction is necessary for reclaiming space in SMR disks
- But, how do you approach band compaction?
- Our work focuses on minimizing long term data movement over the life of a shingled disk
 - We use write heat as a metric to reduce this long term data movement
- Simulated LFS with a block-based API



Why Do You Need to Classify Data?

- Perform band compaction less often
- Moving fewer blocks when doing band compaction
 - Improves system responsiveness
 - Reduces overall system activity





Our experiments cover 4-band compaction

- Simulate the effect of compaction in a space-constrained system
- On-demand

Pseudo-code for multiple band compaction:

- Read all live data in the multiple bands
- Sort in ascending order by block write heat
- Write live data to one of the bands read from
- If band is full and there is still live data
 - Write to another of the bands read from



5



- Developed and tested three heuristics
 - Empty (Greedy)
 - Cold-weight
- Only cold-weight considers write heat when classifying data blocks





- Write to all segments in the log
- * When you reach the log's tail
 - Prioritize writing to any empty segment
- If there are no empty segments
 - Select the segment with the least live data









%free + (w_{cold} × %cold) + (w_{hot} × %hot)





%free + (w_{cold} × %cold) + (w_{hot} × %hot)

% free + ($w_{cold} \times %$ cold) + ($w_{hot} \times (1 - %$ free - % cold))





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% free + ($w_{cold} \times %cold$) + w_{hot} - ($w_{hot} \times %free$) - ($w_{hot} \times %cold$)





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% free + % cold ×
$$\left(\frac{W_{cold} - W_{hot}}{1 - W_{hot}}\right)$$



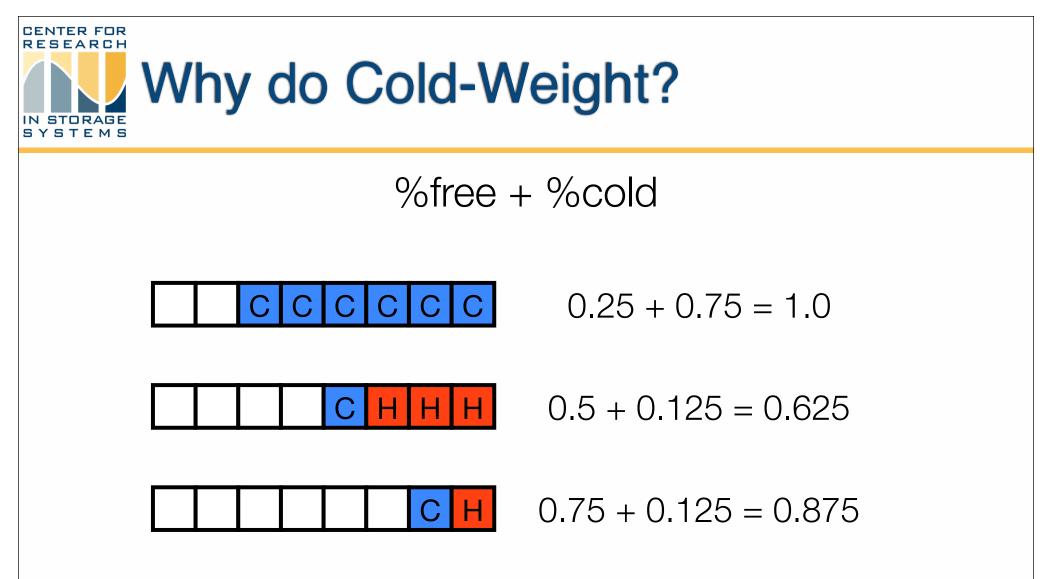


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* Everything can be expressed as a weight on the cold percentage

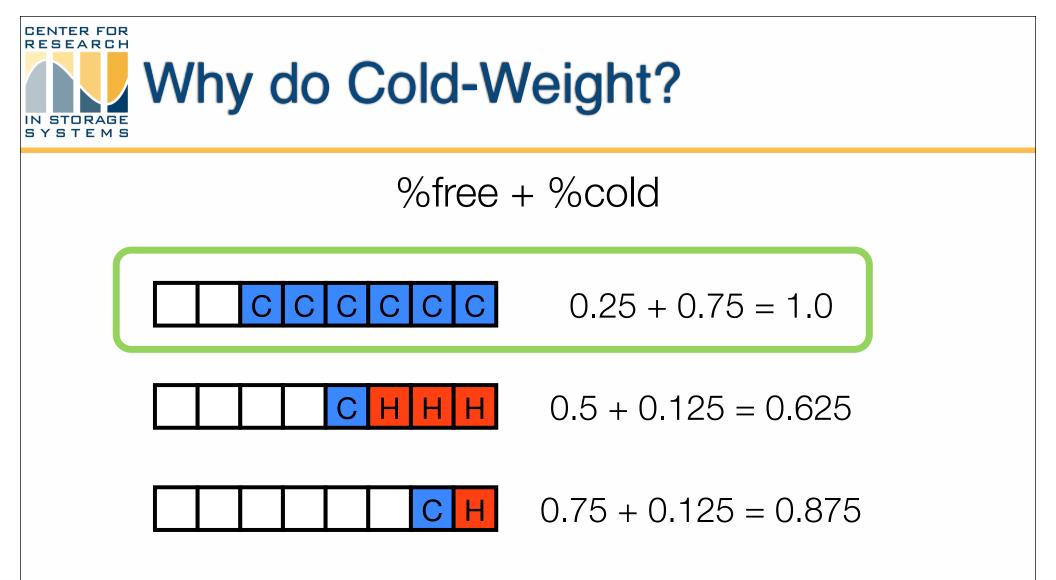
- * Write to all segments in the log
- When you reach the log's tail
 - Prioritize writing to any empty segment
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 - Select the segment with the highest value using the formula





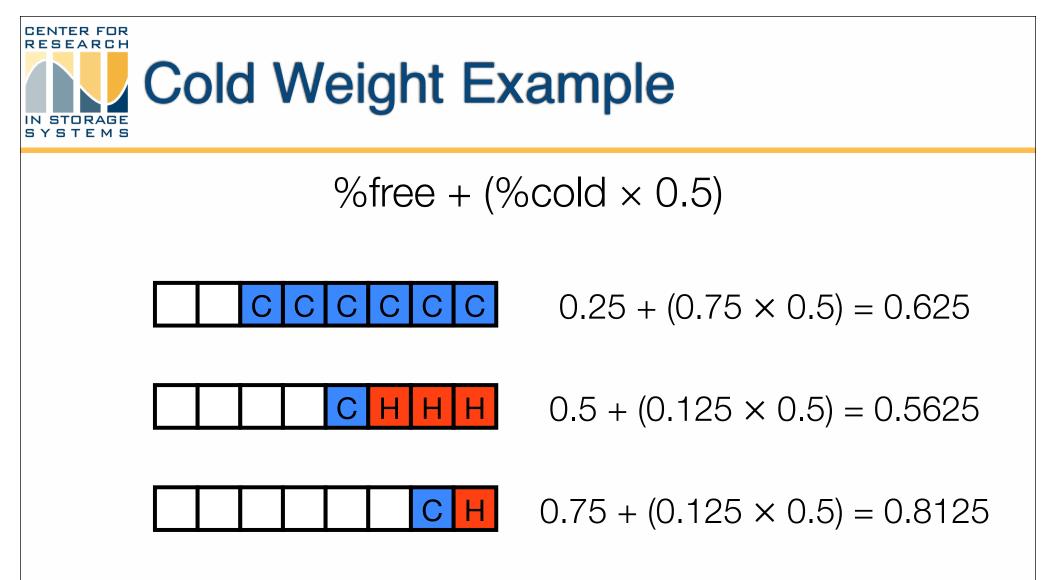
- If you don't put a weight on the cold blocks, they will have equal importance to free blocks
- Which can lead you to pick mostly full segments





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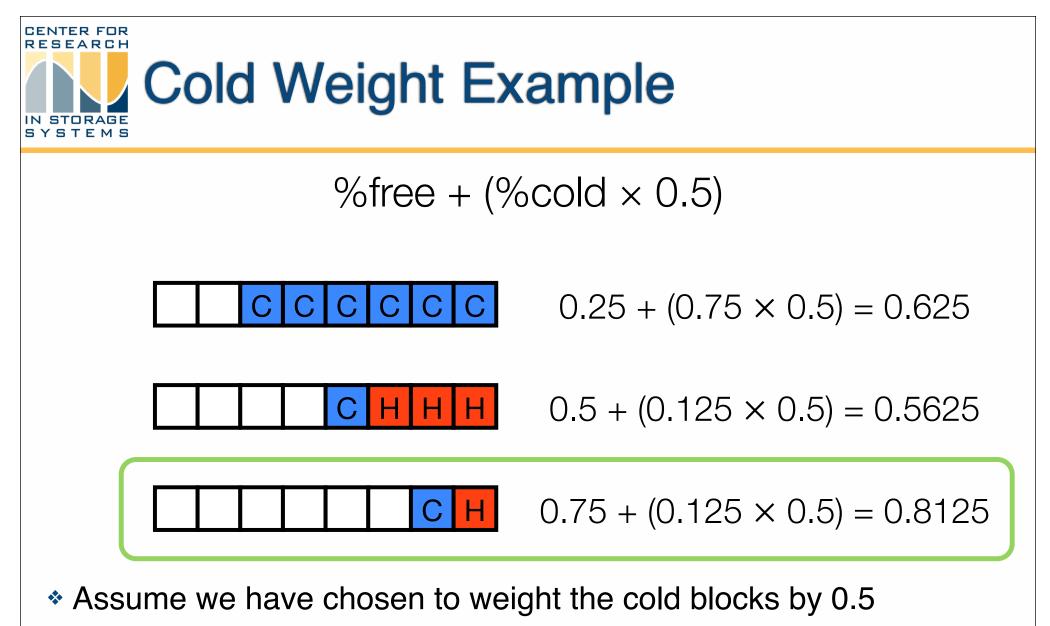




Assume we have chosen to weight the cold blocks by 0.5

* We end up picking the mostly free segment using the weighting



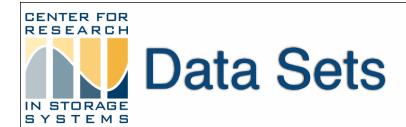


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Cooling + Write Buffer

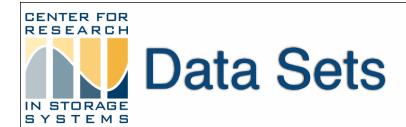
- Data blocks are cooled during segment cleaning
 - All live blocks in the segments selected for cleaning have their heat counts reset to the lowest value
- In order to more accurately separate hot and cold data before it is written to disk, we use a 2-segment sized write buffer
- When the write buffer is full, we determine if it has more hot or cold data
- If it has more hot data than cold we write out the hottest data to the current segment





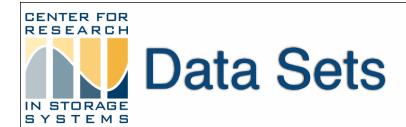
	Project	Source
Number of Write Requests	2,496,935	2,170,271
Total Data Written	26 GB	31 GB
Total Unique Data (live at the end of the trace)	9.5 GB	4.4 GB
Total Data Written Only Once	7.5 GB	3.6 GB





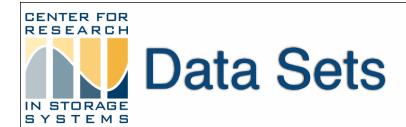
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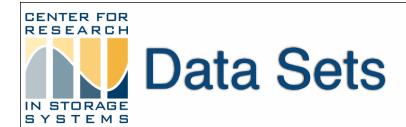
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Experiment	Blocks Moved	% Difference
Empty/Greedy	2,378,357	_
Cold Weight 10%	2,193,595	7.77%
Cold Weight 20%	2,206,924	7.21%
Cold Weight 30%	2,082,264	12.45%
Cold Weight 40%	2,142,427	9.92%
Cold Weight 50%	2,345,437	1.38%

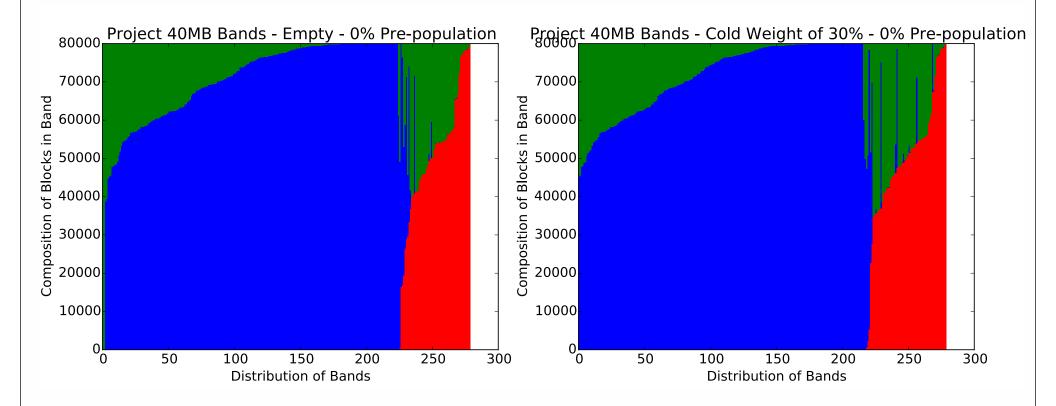




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- * Green is free, blue is cold, red is hot
- We keep cold segments fuller
- * We have less cold data because we've needed to compact fewer segments

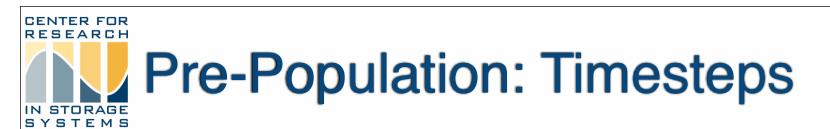




- * We extended the MSR traces by randomly reordering each trace
- We have initially tested our implementation using 2 levels of prepopulation: 50% and 100% pre-population
- We cut the write requests into groups of 10 timesteps
 - A timestep is one second
- * 10 timesteps is at least 10 seconds
 - It could be longer if there is inactivity in the trace

Each level of pre-population is a different random ordering



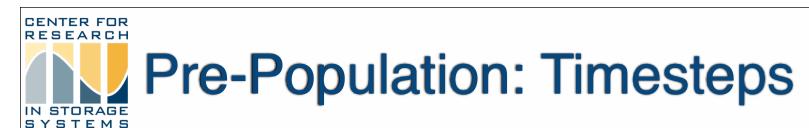


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mestamp		LBA
20.3	W	200
21.1	W	201
24.8	W	202
25.2	W	203
25.7	W	205
25.9	W	206
26.4	W	207
50.0	W	300
54.5	W	400
58.6	W	250
60.7	W	111

- Let's break these up into a group of 3 timesteps
- The first column is timestamp information in seconds
- Each distinct second is a timestep





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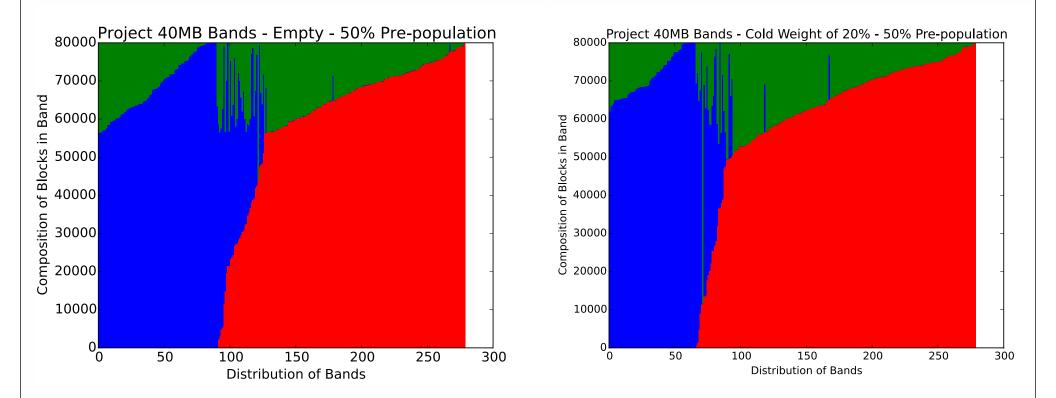
Experiment	Blocks Moved	% Difference
Empty/Greedy	71,568,329	-
Cold Weight 10%	70,985,425	0.81%
Cold Weight 20%	68,478,899	4.32%
Cold Weight 30%	71,966,567	-0.56%
Cold Weight 40%	89,102,638	-24.50%
Cold Weight 50%	3,729,860,479	-5,111%



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Distribution Graphs: 50% Pre-population



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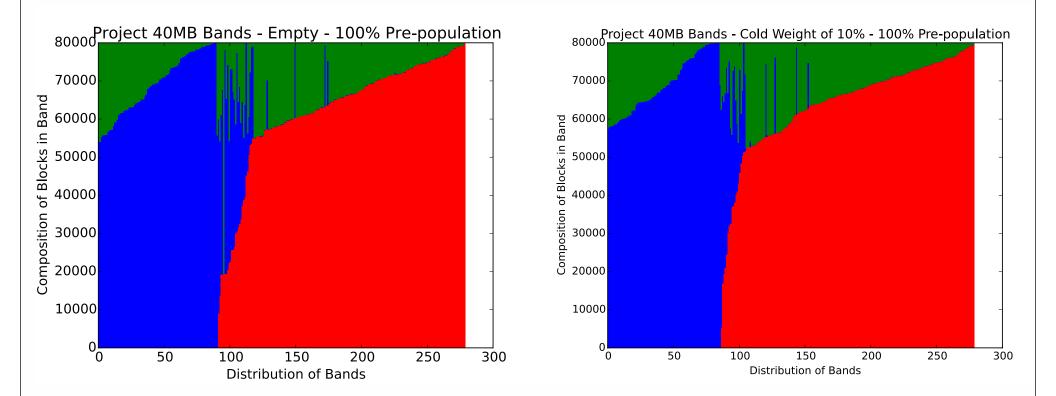
Experiment	Blocks Moved	% Difference
Empty/Greedy	72,702,645	-
Cold Weight 10%	66,031,282	9.18%
Cold Weight 20%	68,707,085	5.50%
Cold Weight 30%	70,874,379	2.51%
Cold Weight 40%	84,283,621	-15.93%
Cold Weight 50%	12,547,962,187	-17,159%



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Distribution Graphs: 100% Pre-population



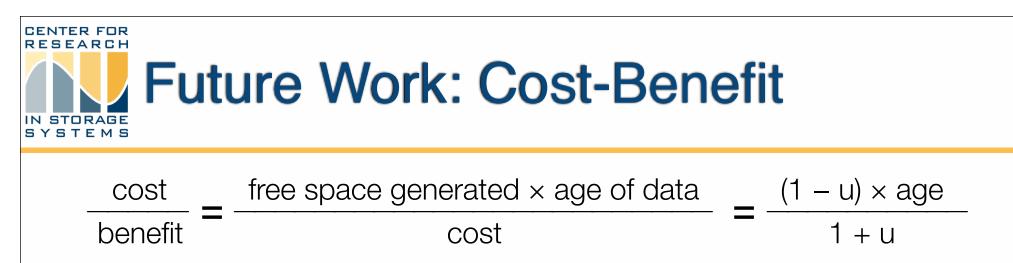
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Separating During Band SYSTEMS Compaction

- * We still have about 5-10% of segments that contain a mix of hot and cold
- Why?
 - This happens because we can write hot data to a band that has been returned by band compaction that is full of cold data
- * This is immediate future work and we will explore two possibilities
 - Hot and cold bands
 - Hot, cold, and "was hot" bands
- Incoming hot data will go to the hot band, cold data and compacted data will go to the cold band
 - "Was hot" will be specifically for data that was hot and is now cold due to compaction





- Using the formula and definitions from Rosenblum's dissertation
 - *u* is the utilization of a segment (how full it is)
 - age is the most recent modify time of any block in a segment
- Write to all segments in the log
- When you reach the log's tail
 - Prioritize writing to any empty segment
- If there are no empty segments
 - Select the segment with the highest value using the cost-benefit formula



Future Work: Dynamic Weighting

- We have promising results with setting static weights
 - They are set at the start of the experiment and are unchanging
- We can improve on these results by manipulating the weight on the cold data
- Our current design will change the weight on code by looking at the overall heat of the data on disk
 - If it's more hot than cold than the weight on cold is more important





- Don't use "how hot is this?", use "how cold is this?"
- Weighting is very important, don't assign equal weights to cold and free





Thank you! Questions? snjones@cs.ucsc.edu

