PERFORMANCE ANALYSIS OF CONTAINERIZED APPLICATIONS ON LOCAL AND REMOTE STORAGE Qiumin Xu¹, Manu Awasthi², Krishna T. Malladi³, Janki Bhimani⁴, Jingpei Yang³, Murali Annavaram¹ ¹USC, ²IIT Gandhinagar ³Samsung ⁴Northeastern

Docker Becomes Very Popular

O Software container platform with many desirable features

□ Ease of deployment, developer friendliness and light virtualization

Mainstay in cloud platforms

Google Cloud Platform, Amazon EC2, Microsoft Azure

Storage Hierarchy is the key component

- **High Performance SSDs**
- **NVMe, NVMe over Fabrics**



Agenda

- **Docker, NVMe and NVMe over Fabrics (NVMf)**
- **D** How to best utilize NVMe SSDs for single container?
 - **D** Best configuration performs similar to raw performance
 - **Where do the performance anomalies come from?**
- Do Docker containers scale well on NVMe SSDs?
 - **Exemplify using Cassandra**
 - **Best strategy to divide the resources**
- **C** Scaling Docker containers on NVMe-over-Fabrics

What is Docker Container? Virtual Machine App A App B Bins/Libs **Bins/Libs** Each virtualized application Guest OS Guest OS includes an entire OS (~10s of GB) Hypervisor Host OS Server

Docker Container



Docker container comprises just application and bins/libs

C Shares the kernel with other container

Much more portable and efficient

figure from https://docs.docker.com

Non-Volatile Memory Express (NVMe)

A storage protocol standard on top of PCIe

NVMe SSDs connect through PCIe and support the standard

Since 2014 (Intel, Samsung)

Enterprise and consumer variants

NVMe SSDs leverage the interface to deliver superior perf

5X to 10X over SATA SSD^[1]





[1] Qiumin Xu et al. "Performance analysis of NVMe SSDs and their implication on real world databases." SYSTOR'15

Why NVMe over Fabrics (NVMf)?

Retains NVMe performance over network fabrics
 Eliminate unnecessary protocol translations
 Enables low latency and high IOPS remote storage



J. M. Dave Minturn, "Under the Hood with NVMe over Fabrics,", SINA Ethernet Storage Forum

Storage Architecture in Docker Storage Options: Through Docker Filesystem (Aufs, Btrfs, Overlayfs) Through Virtual Block Devices (2.a Loop-lvm, 2.b Direct-lvm) Through Docker Data Volume (-v)



Optimize Storage Configuration for Single Container Experimental Environment

Dual-socket, 12 HT cores Xeon E5-2670 V3

O enterprise-class NVMe SSD

- □ Samsung XS1715
- **C** kernel v4.6.0
- **Docker v1.11.2**
- **D** fio used for traffic generation
 - **Asynchronous IO engine**, libaio
 - **32** concurrent jobs and iodepth is 32
 - Measure steady state performance



Performance Comparison — Host Backing Filesystems



EXT4 performs 25% worse for RR

XFS performs closely resembles RAW for all but RW



XFS allows multiple processes to read a file at once

Uses allocation groups which can be accessed independently

EXT4 requires mutex locks even for read operations

Tuning the Performance Gap —Random Writes



XFS performs poorly with high thread count

Contention in exclusive locking kills the write performance

- **Used by extent look up and write checks**
- **Patch available but not for Linux 4.6 [1]**

[1] https://www.percona.com/blog/2012/03/15/ext4-vs-xfs-on-ssd/

Storage Architecture in Docker Storage Options: Through Docker Filesystem (Aufs, Btrfs, Overlayfs) Through Virtual Block Devices (2.a Loop-lvm, 2.b Direct-lvm) Through Docker Data Volume (-v)



Docker Storage Options

Option 1: Through Docker File System

Q Aufs (Advanced multi-layered Unification FileSystem):

A fast reliable unification file system

Btrfs (B-tree file system):

A modern CoW file system which implements many advanced features for fault tolerance, repair and easy administration

Overlayfs:

Another modern unification file system which has simpler design and potentially faster than Aufs



Aufs and Overlayfs performs close to raw block device for most cases

Btrfs has the worst performance for random workloads

Tuning the Performance Gap of Btrfs —Random Reads



Btrfs doesn't work well for small block size yet

D Btrfs must read the file extent before reading the file data.

D Large block size reduces the frequency of reading metadata

Tuning the Performance Gap of Btrfs —Random Reads



Btrfs doesn't work well for random writes due to CoW overhead

Storage Architecture in Docker Storage Options: Through Docker Filesystem (Aufs, Btrfs, Overlayfs) 2 Through Virtual Block Devices (2.a Loop-lvm, 2.b Direct-lvm) **3 Through Docker Data Volume (-v) Container Read / Write Operations 1** -g option **S**-v option 0 **Storage Driver** Aufs, Btrfs,



Docker Storage Configurations

Option 2: Through Virtual Block Device

Devicemapper storage driver leverages the thin provisioning and snapshotting capabilities of the kernel based Device Mapper Framework

D Loop-lvm uses sparse files to build the thin-provisioned pools

Direct-lvm uses block device to directly create the thin pools (Recommended by Docker)

Docker Storage Configurations

Option 3: Through Docker Data Volume (-v)

Data persists beyond the lifetime of the container and can be shared and accessed from other containers



* figure from https://github.com/libopenstorage/openstorage

Performance Comparison Option 2 & Option 3



Direct-lvm has worse performance for RR/RW

LVM, device mapper, and the dm-thinp kernel module introduced additional code paths and overhead may not suit IO intensive workloads

Application Performance

Cassandra Database



NoSQL database

C Scale linearly to the number of nodes in the cluster (theoretically) ^[1]

Requires data persistence

U uses docker volume to store data

[1] Rabl, Tilmann et al. "Solving Big Data Challenges for Enterprise Application Performance Management", VLDB'13

Scaling Docker Containers on NVMe

multiple containerized Cassandra Databases

Experiment Setup

Dual socket, Xeon E5 server, 10Gb ethernet
N = 1, 2, 3, ... 8 containers
Each container is driven by a YCSB client
Record Count: 100M records, 100GB in each DB
Client thread count: 16

Workloads

Workload A, 50% read, 50% update, Zipfian distribution
 Workload D, 95% read, 5% insert, normal distribution

Results-Throughput

Workload D, directly attached SSD



Aggregated throughput peaks at 4 containers

 Cgroups: 6 CPU cores, 6GB memory, 400MB/s bandwidth

Strategies for Dividing Resources



MEM has the most significant impact on throughput
 Best strategy for dividing resources using cgroups
 Assign 6 CPU cores for each container, leave other resource uncontrolled



Results-Throughput



The throughput of NVMf is within 6% - 12% compared to directly attached SSDs

Results-Latency



NVMF incurs only 2% - 15% longer latency than direct attached SSD.

Results-CPU Utilization

CPU Utilization on Target Machine 2.0% 1.5% 1.0% 0.5% 0.0% 1 2 3 8 4 5 6 7 **# of Cassandra Instances**

NVMF incurs less than **1.8%** CPU Utilization on Target Machine

SUMMARY

Best Option in Docker for NVMe Drive Performance Overlay FS + XFS + Data Volume

Best Strategy for Dividing Resources using Cgroups Control only the CPU resources

Scaling Docker Containers on NVMf Throughput: within 6% - 12% vs. DAS Latency: 2% - 15% longer than DAS

THANK YOU! QIUMIN@USC.EDU