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Accelerating Relative-error Bounded Lossy Compression for HPC datasets with Precomputation-Based Mechanisms

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Outline

- Background of research
- Our design
- Evaluation
- Conclusion



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Background

- Scientific simulations
 - Climate scientists need to run large ensembles of high-fidelity 1kmX1km simulations. Estimating even one ensemble member per simulated day may generate **260 TB** of data every **16s** across the ensemble.
 - Cosmological simulation may produce **40PB** of data when simulating 1 trillion of particles in hundreds of snapshots.
- Data reduction is required
 - Lossless compression
 - Simulation data often exhibit high entropy
 - Reduction ratio usually around 2:1
 - Lossy compression
 - More aggressive data reduction scheme
 - High reduction ratio



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Background - Lossy compressors

- ZFP

- follow the classic texture compression for image data
- Data transformation + embedded coding
- **Low** compression ratio , **High** compression speed

- SZ

- Prediction + quantization + Huffman encoding + Zstd
- **High** compression ratio, **Low** compression speed

- A dilemma: which compressor should I use?



- Question: Can we significantly improve compression speed for SZ, leading to an easy solution for users?



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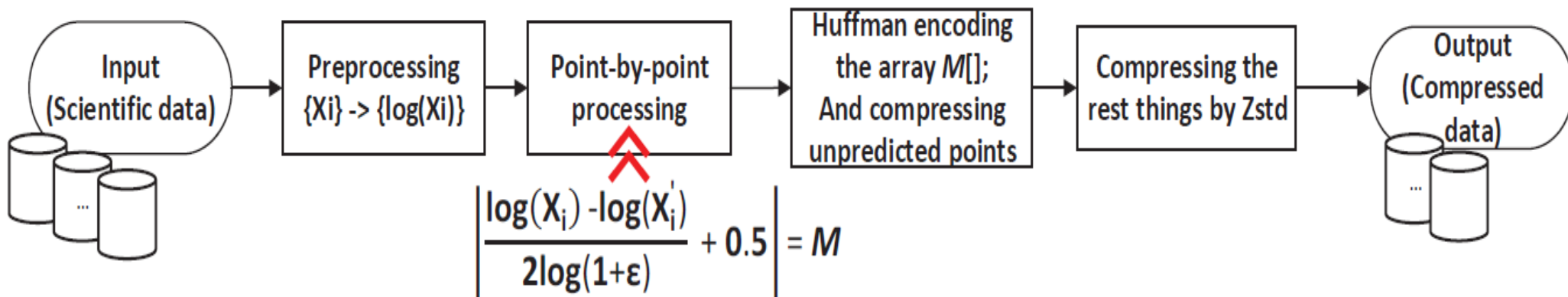
Background - Lossy compression error bound

- Absolute error bound
 - For a value f , we get $f' \in (f - \epsilon, f + \epsilon)$ is acceptable
- Pointwise relative error bound
 - For a value f , we get $f' \in (f * (1 - \epsilon), f * (1 + \epsilon))$ is acceptable
- CLUSTER18: Convert a pointwise relative error bound to an absolute error bound with a logarithmic transformation
 - $\log(f*(1 - \epsilon)) = \log(f) + \log(1 - \epsilon)$, $\log(f*(1 + \epsilon)) = \log(f) + \log(1 + \epsilon)$
 - $\log(f') \in (\log(f) + \log(1 - \epsilon), \log(f) + \log(1 + \epsilon))$



Background – design of SZ compressor for relative error control

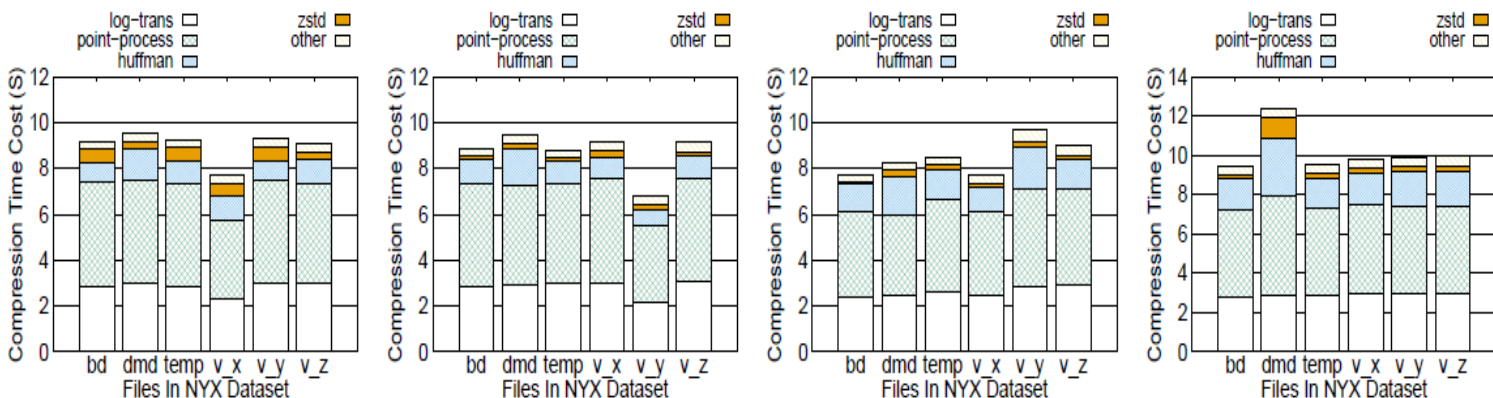
- Preprocess - Logarithmic transformation
- Point-by-point processing – prediction & quantization
- Huffman encode
- Compression with lossless compressor



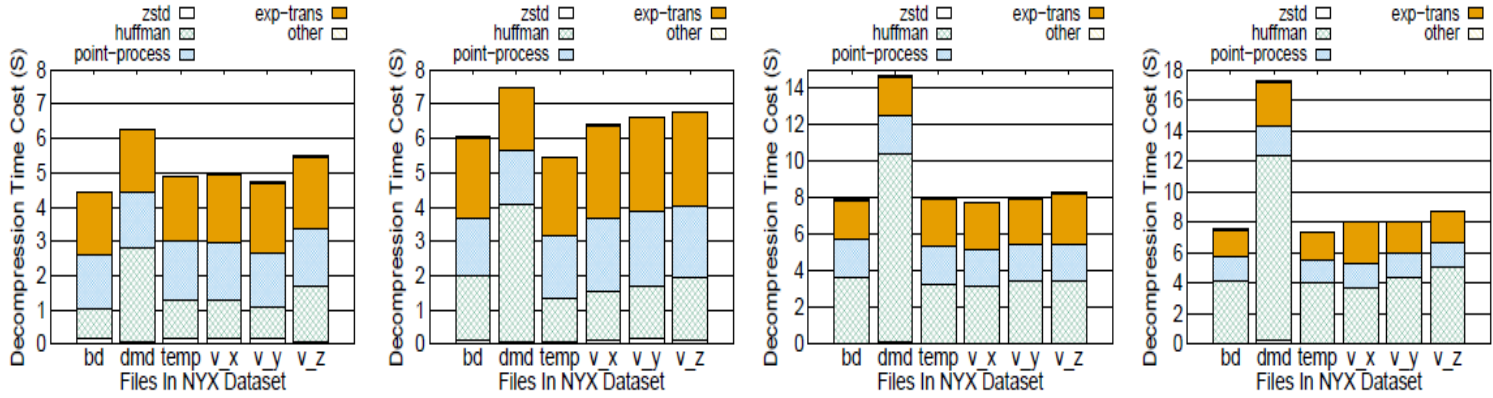
Logarithmic transformation (logX) is too expensive!



Performance breakdown of SZ Compression/Decompression



(a) 1.00E-01 (b) 1.00E-02 (c) 1.00E-03 (d) 1.00E-04



(a) 1.00E-01 (b) 1.00E-02 (c) 1.00E-03 (d) 1.00E-04

Time costs on log-trans and exp-trans stages consist about 1/3 of the total



Our design - workflow

f is a float *M* is an integer

① Quantization: $f = \frac{X_i}{X_i} = (1 + \epsilon)^{2M-1+2\delta}, \quad 0 < \delta < 1$

② Lookup the precomputed **Table T1**: $f \rightarrow M$, for point-by-point processing to obtain the $M[]$.

Table T1: $f \rightarrow M$	
$(1+\epsilon)^{2(0-K/2-1)} \sim (1+\epsilon)^{2(0-K/2+1)}$	$\rightarrow 0$
$(1+\epsilon)^{2(1-K/2-1)} \sim (1+\epsilon)^{2(1-K/2+1)}$	$\rightarrow 1$
$(1+\epsilon)^{2(2-K/2-1)} \sim (1+\epsilon)^{2(2-K/2+1)}$	$\rightarrow 2$
$(1+\epsilon)^{2(3-K/2-1)} \sim (1+\epsilon)^{2(3-K/2+1)}$	$\rightarrow 3$
...	

③ Entropy-encoding the $M[]$, and compressing the rest things by Ztsd.

④ Lookup the precomputed **Table T2**: $M \rightarrow f$, point-by-point, for decompression.

Table T2: $M \rightarrow f$	
0	$\rightarrow (1+\epsilon)^{2(0-K/2)}$
1	$\rightarrow (1+\epsilon)^{2(1-K/2)}$
2	$\rightarrow (1+\epsilon)^{2(2-K/2)}$
4	$\rightarrow (1+\epsilon)^{2(4-K/2)}$
...	

- No longer to calculate the quantization factor, but look up tables.
- Using Table T1 to get quantization factor from f
- Using Table T2 to get a approximate value of f from quantization factor



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Our design - Model A

- A simple idea
- We call $((1 + \varepsilon)^{2M-1}, (1 + \varepsilon)^{2M+1})$ as $PI(M)$
- Separate the value field of f into many grids.
- Each grid maps to a $PI(M)$



A general description to model A

$$f = \frac{X_i}{X'_i} \approx (1 + \varepsilon)^{2M}$$

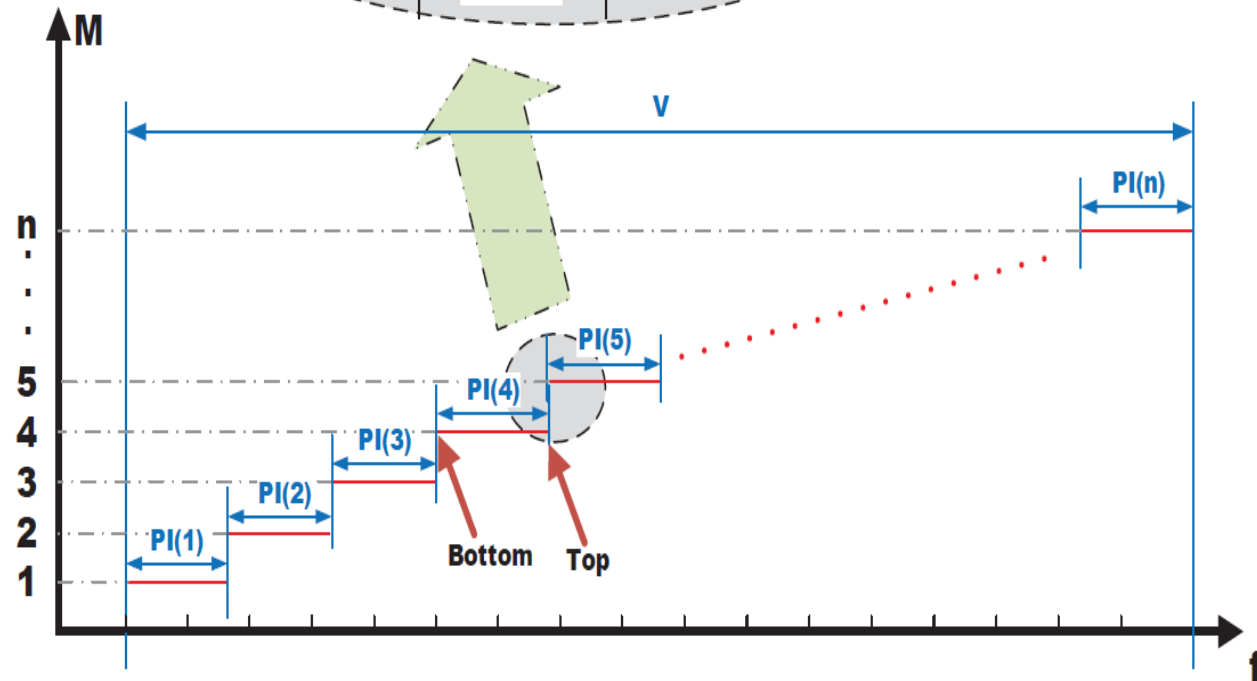
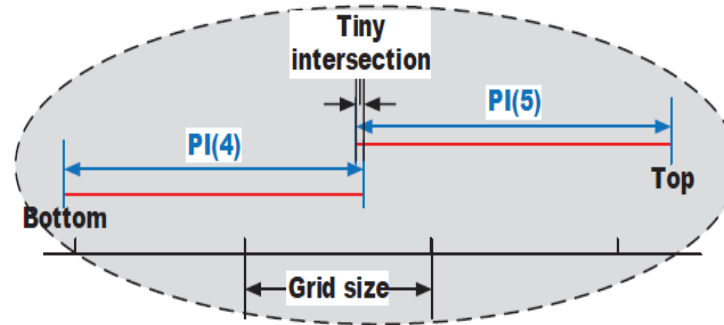


$$(1 + \varepsilon)^{2M-1} \leq f \leq \frac{(1 + \varepsilon)^{2M}}{1 - \varepsilon}$$



PI interval

$$\left[(1 + \varepsilon)^{(2M-1)}, \frac{(1 + \varepsilon)^{2M}}{1 - \varepsilon} \right]$$





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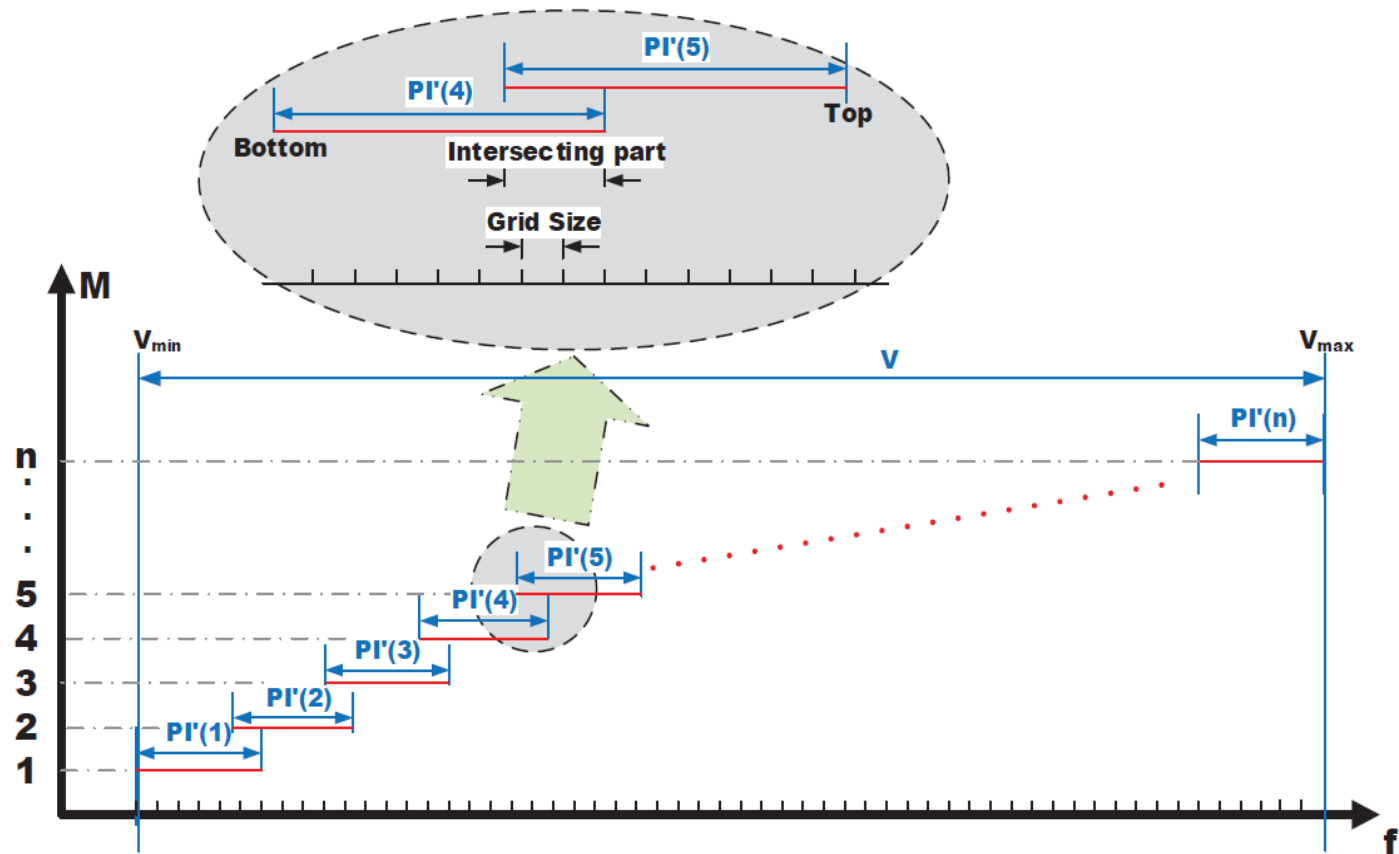
Our design - Model B

- Improved idea
- Let neighbor PIs have more intersection
- Let the size of grids is smaller than any intersection size
- We call $((1 + \varepsilon)^{M(2-\theta)-1}, (1 + \varepsilon)^{M(2-\theta)+1})$ as $PI'(M)$
- Each grid maps to a $PI'(M)$

$$(1 + \varepsilon)^{M(2-\theta)-1} \leq f' \leq \frac{(1 + \varepsilon)^{M(2-\theta)}}{1 - \varepsilon}, \quad 0 < \theta < 1$$

A general description about model

P Lemma 1. *If a grid size G is smaller than the size of any intersecting part of PI' , a PI' completely including the grid always exists.*





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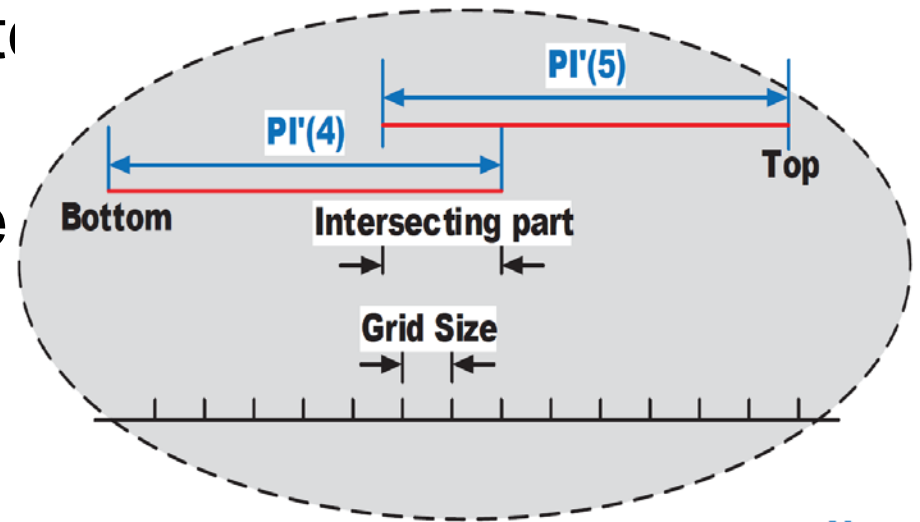


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Our design - Advantage of Model

- B** Any grid (i.e., a data point) is always included in a PI'
- Grid size is smaller than any intersection size, therefore any grid is completely included in a PI'(M)
 - Effect: Strictly respecting the bound

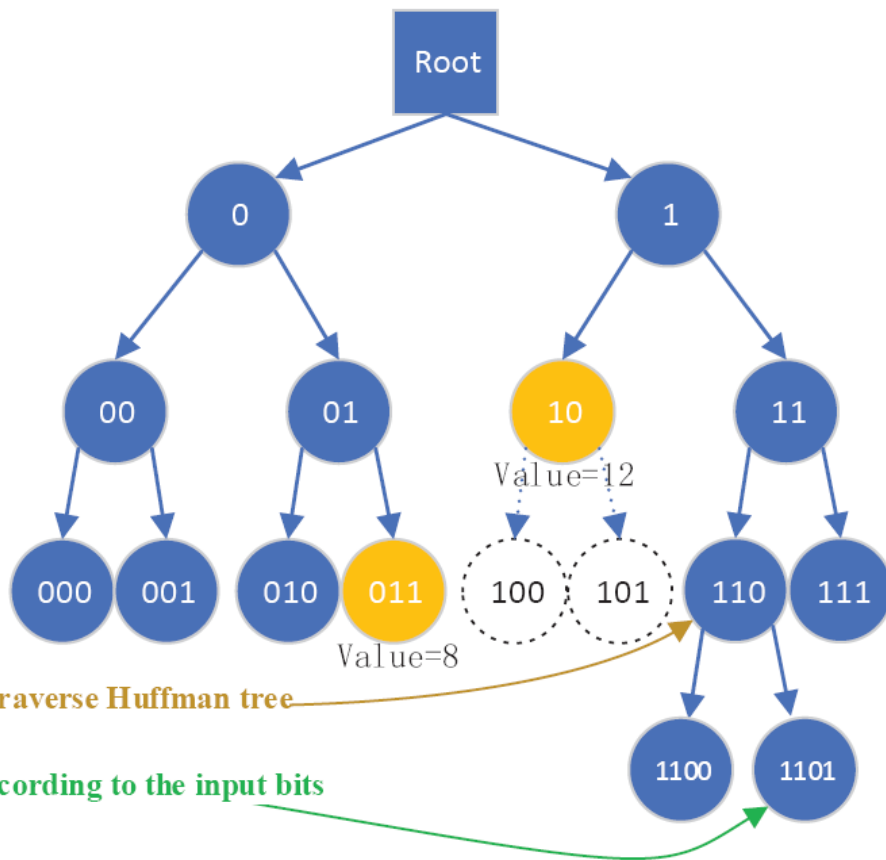




Accelerating Huffman decoding

Idea: building precomputed table to accelerate Huffman decoding

Value Table	Length Table	Node Table
000 -> -1	000 -> 3	000 -> &Node_000
001 -> -1	001 -> 3	001 -> &Node_001
010 -> -1	010 -> 3	010 -> &Node_010
011 -> 8	011 -> 3	011 -> NULL
100 -> 12	100 -> 2	100 -> NULL
101 -> 12	101 -> 2	101 -> NULL
110 -> -1	110 -> 3	110 -> &Node_110
111 -> -1	111 -> 3	111 -> &Node_111



① Get value
 ② Get Huffman code length
 ③ Pop 3 bits
 ④ Get Value
 ⑤
 ⑥ Pop 2 bits
 ⑦
 ⑧ Locate Node
 ⑨ Continue traverse Huffman tree

Input: 011 101 10101...

Output: 8 12

⑩ Choose 1101 according to the input bits



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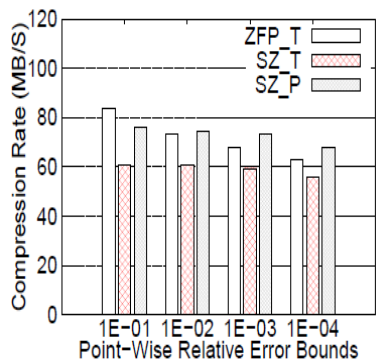
Performance Evaluation

- Environment
 - 2.4GHz Intel Xeon E5-2640 v4 (3D, 3.1GB) processors
 - 256GB memory
- Datasets
 - CESM (2D, 2.0GB)
 - Hurrican (3D, 1.9GB)
 - HACC (1D, 6.3GB)

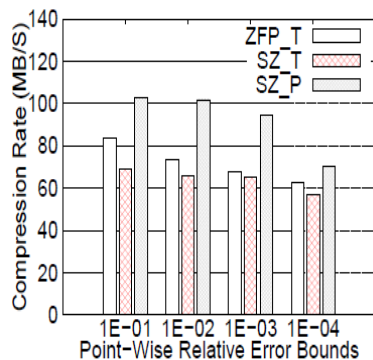


Compression/Decompression

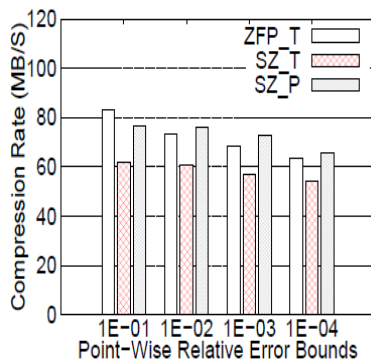
Rate



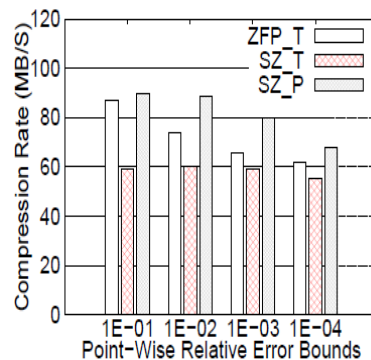
(a) CESM Dataset



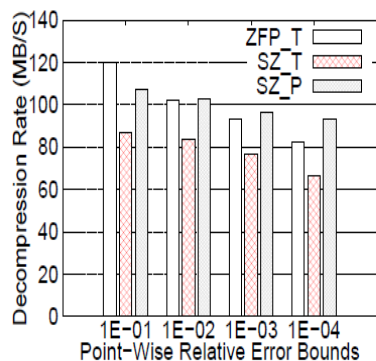
(b) HACC Dataset



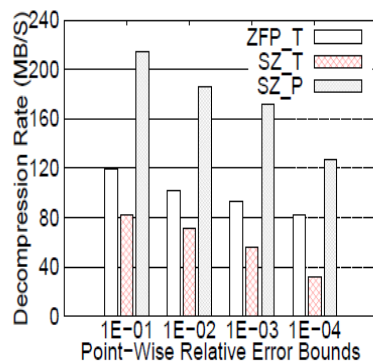
(c) Hurricane Dataset



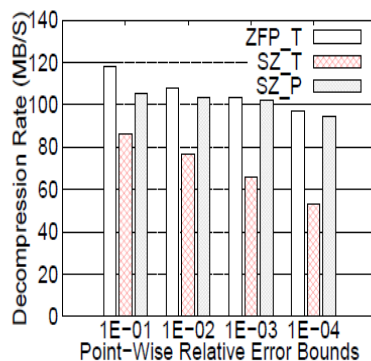
(d) NYX Dataset



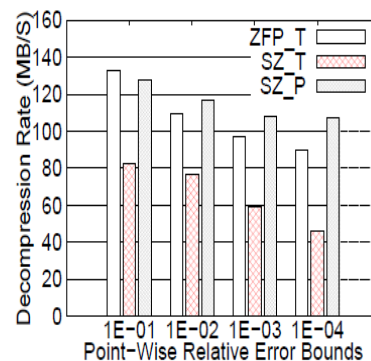
(a) CESM Dataset



(b) HACC Dataset



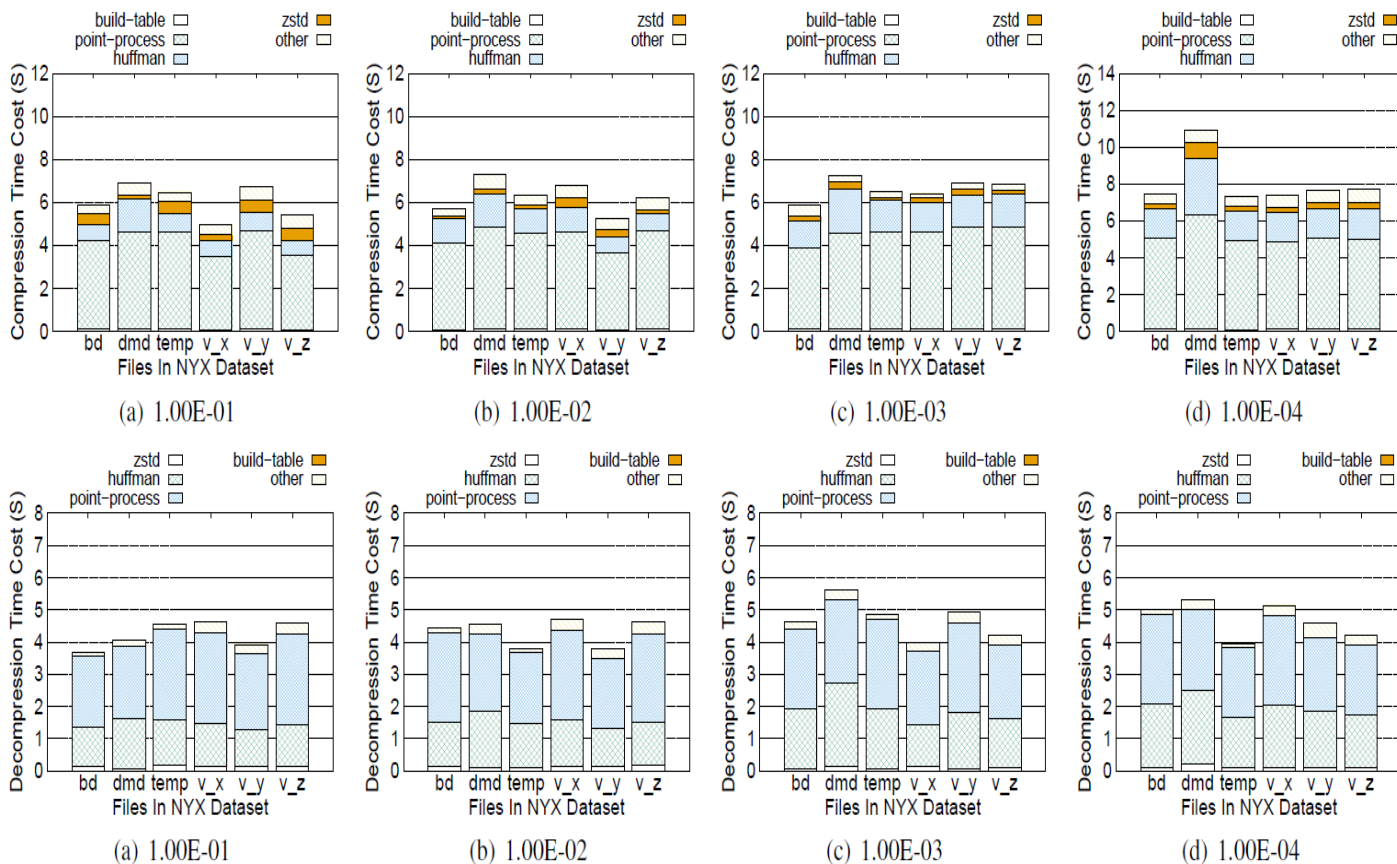
(c) Hurricane Dataset



(d) NYX Dataset

Our Approach is about 1.2x ~ 1.5x than original SZ on compression rate and 1.3x ~ 3.0x on decompression rate.

Compression/Decompression breakdown

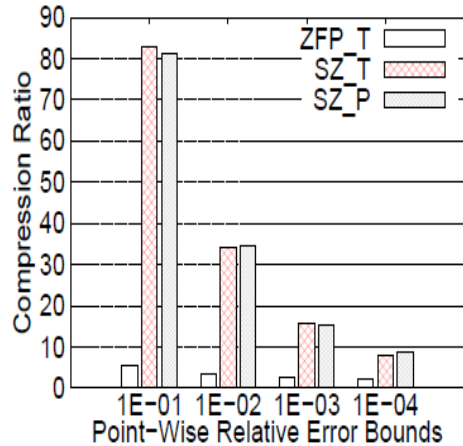


No time cost on log-trans and exp-trans. Time cost on build-table stage is very small.

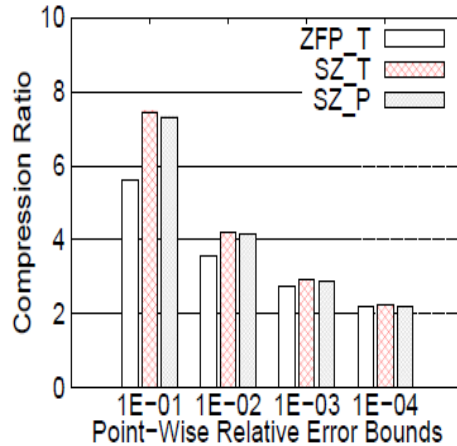


Compression Ratio

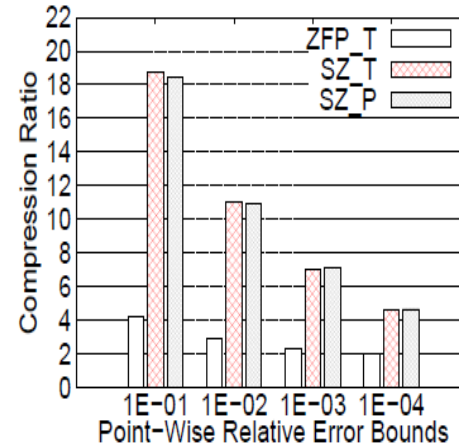
We can observe that our solution (SZ_T) has very similar compression ratios with SZ_T.



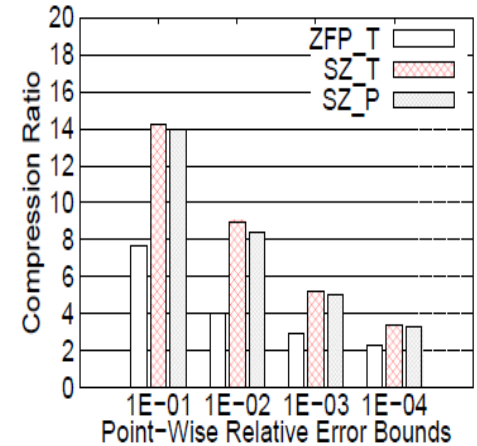
(a) CESM Dataset



(b) HACC Dataset



(c) Hurricane Dataset



(d) NYX Dataset



Data quality

POINT-WISE RELATIVE ERROR BOUND ON 3 REPRESENTATIVE FIELDS IN NYX.

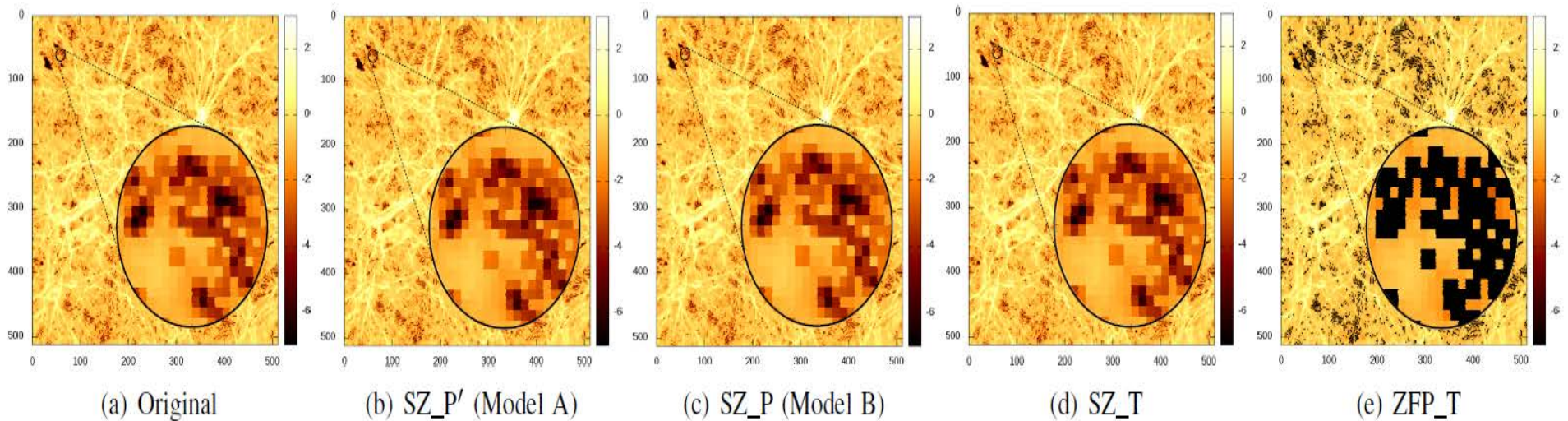
pwrEb	type	dark_matter_density				velocity_x				temperature			
		MAX E	NRMSE	PSNR	CR	MAX E	NRMSE	PSNR	CR	MAX E	NRMSE	PSNR	CR
1E-01	SZ_P'	1.52E-01	3.97E-05	88.02	6.25	1.50E-01	2.27E-04	72.88	23.48	1.52E-01	4.70E-03	46.56	18.83
	SZ_T	1.00E-01	3.49E-05	89.13	6.19	1.00E-01	2.05E-04	73.75	24.97	1.00E-01	4.32E-03	47.29	19.85
	SZ_P	9.99E-02	3.26E-05	89.73	6.03	9.97E-02	1.96E-04	74.15	25.99	9.97E-02	4.10E-03	47.75	20.79
	ZFP_T	5.07E-02	4.64E-06	106.66	3.32	4.80E-02	2.89E-04	70.79	18.40	5.17E-02	2.74E-05	91.25	14.00
1E-02	SZ_P'	1.75E-02	4.05E-06	107.84	3.85	1.70E-02	2.36E-05	92.55	13.46	1.70E-02	5.39E-04	65.37	14.37
	SZ_T	1.00E-02	3.55E-06	108.99	3.85	1.00E-02	2.00E-05	93.97	14.06	1.00E-02	4.50E-04	66.93	13.55
	SZ_P	1.00E-02	3.42E-06	109.31	3.80	9.96E-03	1.82E-05	94.78	12.98	9.96E-03	4.23E-04	67.47	11.93
	ZFP_T	3.02E-03	2.75E-07	131.22	2.35	3.33E-03	3.45E-05	89.24	6.59	3.16E-03	1.75E-06	115.15	5.21
1E-03	SZ_P'	1.96E-03	4.30E-07	127.34	2.75	1.95E-03	2.58E-06	111.76	6.75	1.95E-03	6.70E-05	83.48	8.02
	SZ_T	9.97E-04	3.51E-07	129.10	2.74	9.98E-04	1.98E-06	114.08	6.61	9.98E-04	4.51E-05	86.91	7.63
	SZ_P	1.00E-03	3.44E-07	129.27	2.72	9.99E-04	1.79E-06	114.93	6.49	9.99E-04	4.25E-05	87.44	7.12
	ZFP_T	3.90E-04	3.56E-08	148.98	1.92	3.95E-04	4.63E-06	106.69	4.08	3.97E-04	2.23E-07	133.04	3.50
1E-04	SZ_P'	1.60E-04	3.96E-08	148.04	2.12	1.60E-04	2.10E-07	133.55	3.93	1.60E-04	4.98E-06	106.05	4.39
	SZ_T	9.80E-05	3.48E-08	149.16	2.09	9.90E-05	1.98E-07	134.05	3.92	9.90E-05	4.43E-06	107.08	4.38
	SZ_P	1.00E-04	3.38E-08	149.43	2.01	1.00E-04	1.72E-07	135.29	3.88	1.00E-04	4.25E-06	107.43	4.27
	ZFP_T	5.08E-05	4.49E-09	166.96	1.63	4.99E-05	5.81E-07	124.71	2.95	5.33E-05	2.76E-08	151.19	2.63

Comparable compression ratios with related works
(SZ_T and ZFP_T)

Data quality (Cont'd)

Visualization of decompressed dark matter density dataset (slice 200) at the compression ratio of 2.75.

SZ series has a better visual quality than ZFP does.
SZ_P (both mode A and B) lead to satisfied visual quality!





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Conclusion

- We accelerate the SZ compressor for point-wise relative error bound control by designing a table-lookup method.
- We control the error bound strictly by an in-depth analysis of mapping relation between predicted value and quantization factor.
- Experiments show that 1.2x ~ 1.5x on compression speed and 1.3x ~ 3.0x on decompression speed, compared with SZ 2.1.



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Thank you

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