

# A Performance Study of LFSCK

## Bottlenecks and Potentials

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# Outline

- Motivation
- Background
- Performance Study Methods
- Results and Analysis
- Potentials
- Conclusions and Plans

# Motivation

- Lustre is one of the mostly used high performance parallel file systems on the market
- It is deployed on large and complex systems
  - ANL's coming Exascale machine: Aurora<sup>[1]</sup>
  - Planned to use Lustre file system
  - File System Capacity: > 150 petabytes
  - File System Throughput: > 1 terabyte/s
- Complex systems may fail for various of reasons
  - Management issues
  - Unknown software/hardware bugs
- Extreme failure cases may lead to data loss
  - LFSCK

[1] ANL Aurora Programs: <https://www.alcf.anl.gov/programs/aurora-esp>

- The tool to check and fix problems of Lustre file system
  - `lfsck`: early version of file system checker for Lustre
    - Utilizes databases created via `e2fsck` (`e2fsprogs`)
    - Very slow due to large database size
  - LFSCCK: online Lustre file system checker
    - Available in Lustre 2.x
    - Lustre 2.4: can verify and repair the directory FID-in-dirent and LinkEA consistency
    - Lustre 2.6: can verify and repair MDT-OST file layout consistency
    - Lustre 2.7: can support verify and repair inconsistencies between multiple MDTs
    - ...
- Current Status
  - LFSCCK 1 -> LFSCCK 1.5 -> LFSCCK 2 -> LFSCCK 3 -> LFSCCK 4
  - LFSCCK-4: LFSCCK performance enhancement
    - LU-5820: evaluation: linkEA verification history in RAM performance

# LFSCCK - Issues

- LFSCCK is a powerful tool to combat failures and inconsistencies in Lustre file system
- However, it is not perfect yet...
- Functionality
  - It might not be able to identify nor fix failures of Lustre
  - Please check our study in [ICS'18]

Node(s) Affected	Fault Models	LFSCCK	WikiR	WikiW-async	WikiW-sync
MGS	a-DevFail	normal	✓	✓	✓
	b-Inconsist	normal	✓	✓	✓
	c-Network	normal	✓	✓	✓
MDS	a-DevFail	<b>Invalid</b>	hang	hang	hang
	a-DevFail (v2.10)	<b>I/O err</b>	I/O err	I/O err	I/O err
	b-Inconsist	normal	✓	✓	✓
	c-Network	<b>I/O err</b>	hang	hang	hang
OSS#2	a-DevFail	<b>hang</b>	hang	hang	hang
	a-DevFail (v2.10)	normal	✓	✓	✓
	b-Inconsist	<b>reboot</b>	corrupt	hang	hang
	c-Network	<b>hang</b>	hang	hang	hang
three OSSes	a-DevFail	<b>hang</b>	hang	hang	hang
	a-DevFail (v2.10)	normal	✓	hang	hang
	b-Inconsist	<b>reboot</b>	corrupt	hang	hang
	c-Network	<b>hang</b>	hang	hang	hang
MDS + OSS#2	a-DevFail	<b>Invalid</b>	hang	hang	hang
	a-DevFail (v2.10)	<b>I/O err</b>	I/O err	I/O err	I/O err
	b-Inconsist	<b>reboot</b>	corrupt	hang	hang
	c-Network	<b>I/O err</b>	hang	hang	hang
	c-Network (v2.10)	<b>hang</b>	hang	hang	hang

**Table 2: Response of LFSCCK and post-LFSCCK Workloads.** The first column shows where the faults are injected. The second column shows the fault models applied (v2.10 means applying to Lustre v2.10). "normal": LFSCCK finishes normally; "reboot": at least one OSS node is forced to reboot; "Invalid": an "Invalid Argument" error; "I/O err": an "Input/Output error"; "hang": cannot finish within one hour; "✓": complete w/o error; "corrupt": checksum mismatch. The bold font highlights the unexpected response of LFSCCK.

[ICS'18] PFAULT: A General Framework for Analyzing the Reliability of High-Performance Parallel File Systems. Jinrui Cao, Om Rameshwar Gatla, Mai Zheng, **Dong Dai**, Vidya Eswarappa, Yan Mu and Yong Chen. Accepted to appear in the proceedings of the 32nd ACM/SIGARCH International Conference on Supercomputing (ICS'18) 2018.

# LFSCCK - Issues

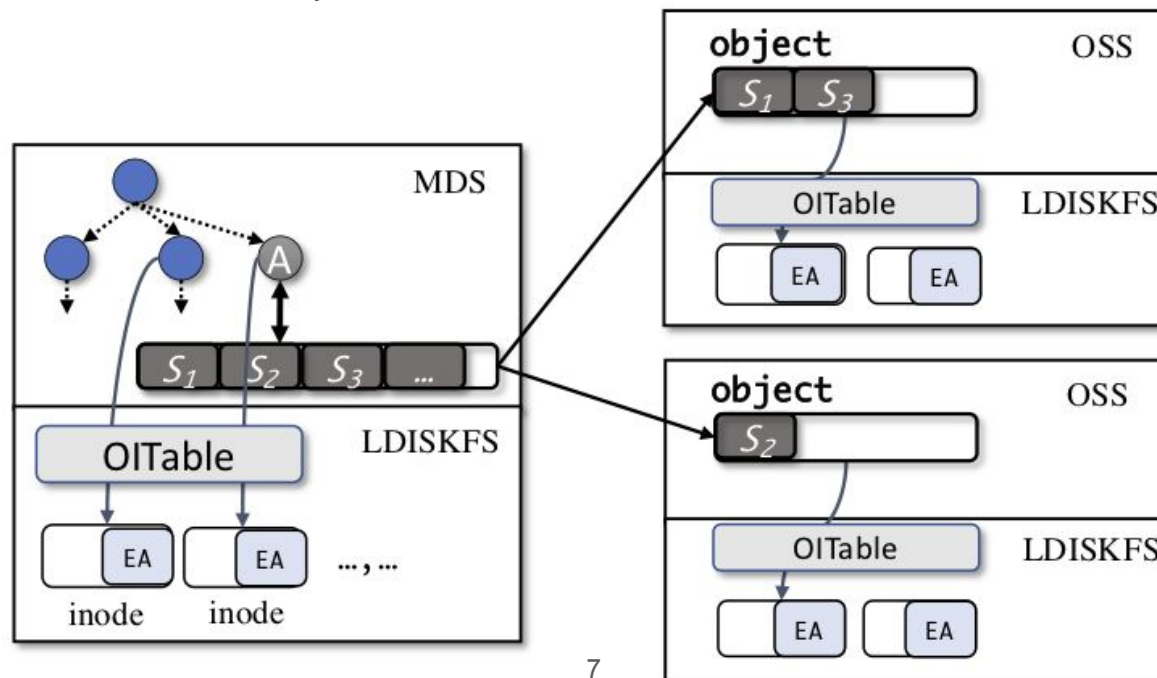
- LFSCCK is a powerful tool to combat failures and inconsistencies in Lustre file system
- However, it is not perfect yet...
- Performance
  - Sometimes, we observed that running file system checking could be still slow, especially on a well-aged Lustre
  - The slowness might come from e2fsck or LFSCCK itself. But, we focus on LFSCCK in this study

*“... experienced a major power outage Sunday, Jan. 3 and another set of outages Tuesday, Jan. 5 that occured while file system were being recovered from the first outage. As a result, there were major losses of important parts of the file systems for .... Lustre areas...”*

*- Quoted from the email about the failure*

# LFSCK Basis

- Lustre
  - Metadata Server (MDS)
  - Object Storage Server (OSS)
- Lustre Metadata Management
  - Mapping metadata between global entity and local inode
  - POSIX namespace metadata
  - Distributed data layout metadata



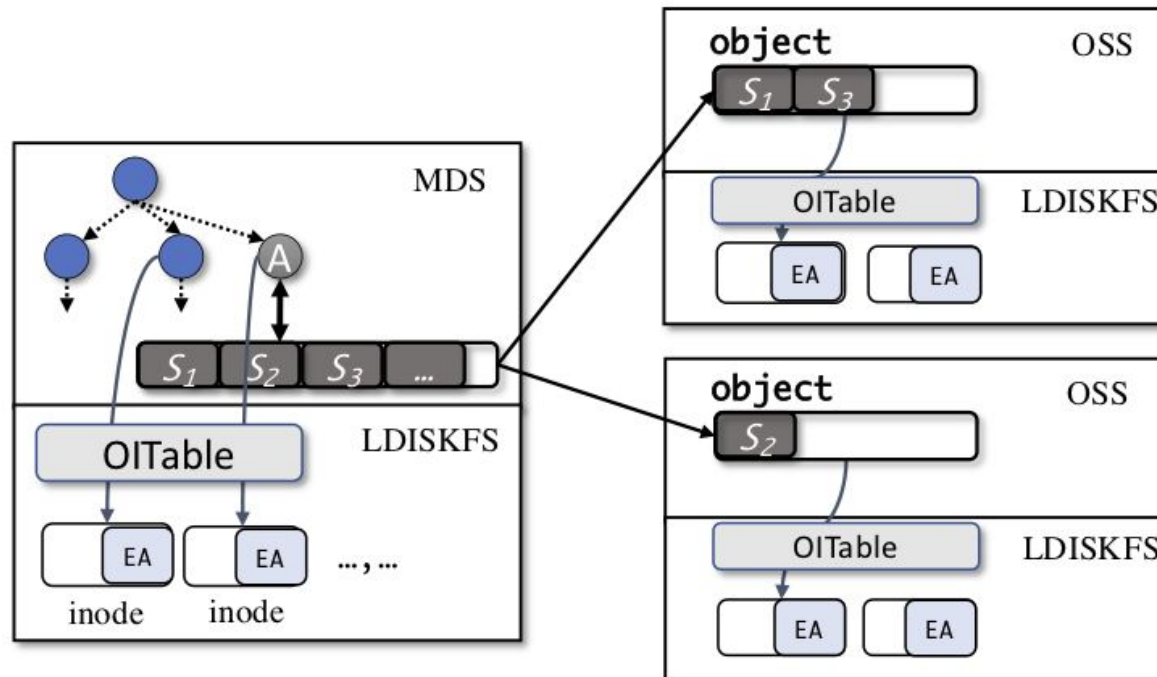
# LFSCK Basis Cont'

- Lustre stores these metadata in more than one places
  - a. So that LFSCK can cross-check them
  - b. For example, global file -> local inode; local inode also has global fid
- LFSCK is a two-stage procedure
  - a. MDS drives the OSS nodes to conduct checking and potential local fix
  - b. Once finish, all OSS nodes will start the second stage to resolve orphan and missing objects detected in the first stage.
- We focus on the first stage
  - a. as most of the time, the number of inconsistencies in a healthy Lustre should be minimal;
    - => indicating a short second stage
  - b. while, we still need to run LFSCK oftenly to guarantee a healthy Lustre;
    - => indicating its importance regarding the performance



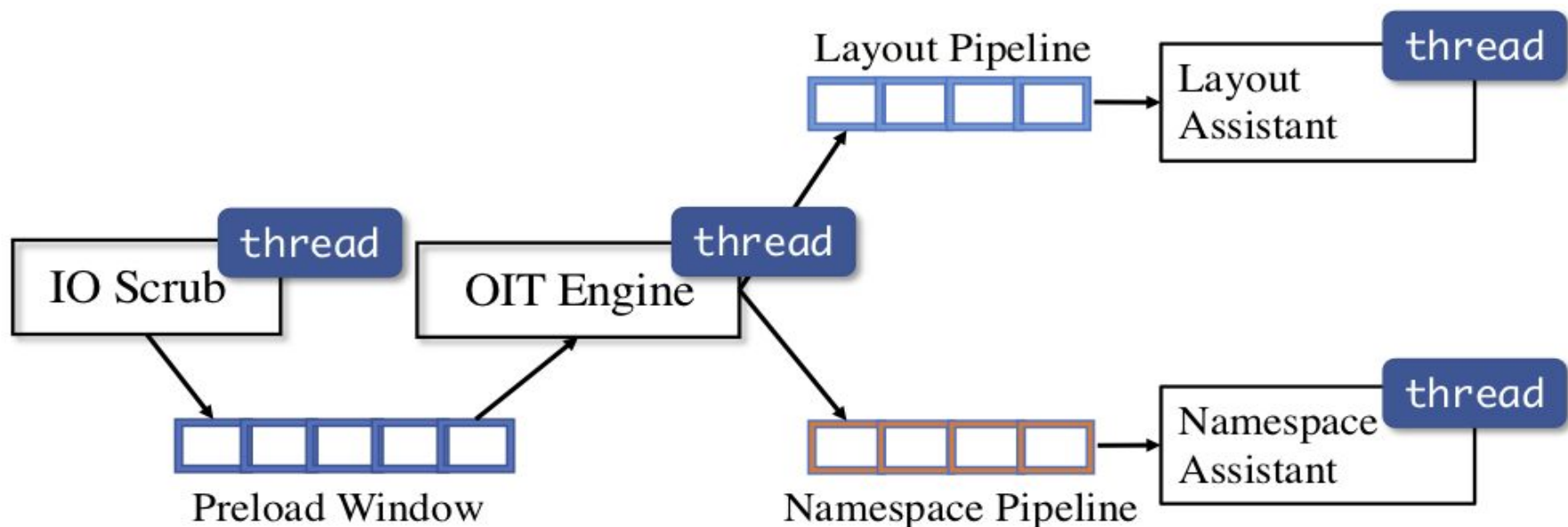
# LFSCK Details

- Three essential checks in the first stage
  - For mapping metadata
    - `FID.local_inode.ext_attr(fid) == FID`
  - For namespace metadata
    - `FID.local_inode.ext_attr(linkEA) == parent.FID`
  - For data layout metadata
    - `FID.Stripe(i).OSS.local_inode.ext_attr(fid) == FID`



# LFSCK Details Cont'

- LFSCK is implemented as several kernel threads connected with kernel buffers.
  - OIscrub sequentially scans the local disks
  - OIT Engine is the driving engine of LFSCK. It checks OITable
  - OIT Engine and IOscrub share buffer, sized `SCRUB_WINDOW_SIZE`
  - Layout/Namespace assistant threads will check Layout and Namespace metadata respectively
  - The buffer sizes are `LFSCK_ASYNC_WIN_DEFAULT`



# Performance Study Methods

- Evaluation Platform

- Our testbed is built on NSF Cloudblab platform
- Small cluster with 1 MDS node, 8 OSS nodes, and 4 client nodes.
- Each node is a c220g1 node
  - Two E5-2630 CPU
  - 128 GB Memory
  - 1.2 TB 10K RPM Hard disk
  - 480GB Intel DC3500 SSD
  - 10GB Intel NIC
- CentOS 7.3
- Lustre 2.10.4
- Reproducible using the same Cloudblab profile
  - We will public our profile and all the test scripts



# Performance Study Methods

- Aging Method

- It is non-trivial to age a file system into proper state for performance testing on file system checkers
  - Especially true for large-scale parallel file system
- We leveraged Darshan in this study
  - Darshan is an I/O characterization tool for HPC
  - Developed by ANL team
  - Deployed at various supercomputers
- Public Darshan logs collected from Intrepid
  - Record all I/O activities done by each applications
- Two Key I/O metadata
  - CP\_SIZE\_AT\_OPEN: how large the file was when the application opened it
  - CP\_BYTES\_WRITTEN: how many bytes were written to the file by this application

**DARSHAN**  
HPC I/O Characterization Tool

# Performance Study Methods

- Random paths
  - Darshan logs were anonymized. No full path or name of the file
  - Generate the structure manually, randomly
    - Randomly generate a directory with random (1->10) depth
    - Randomly generate files (1->100) into the same directory
- Reduce data sizes
  - Files in real-world supercomputers are really big
    - There are files that over 1TB each
  - Reduce files that are larger than 1MB \* 8 to be 8MB.
    - Keep the metadata the same
    - Minimize the occupied data size
- Set Stripe Count
  - We set stripe count to be “-1” and stripe size to be “1MB”
  - Maximize the number of stripes
  - Note that, this setting is **not typical** in production system.

# Performance Study Methods

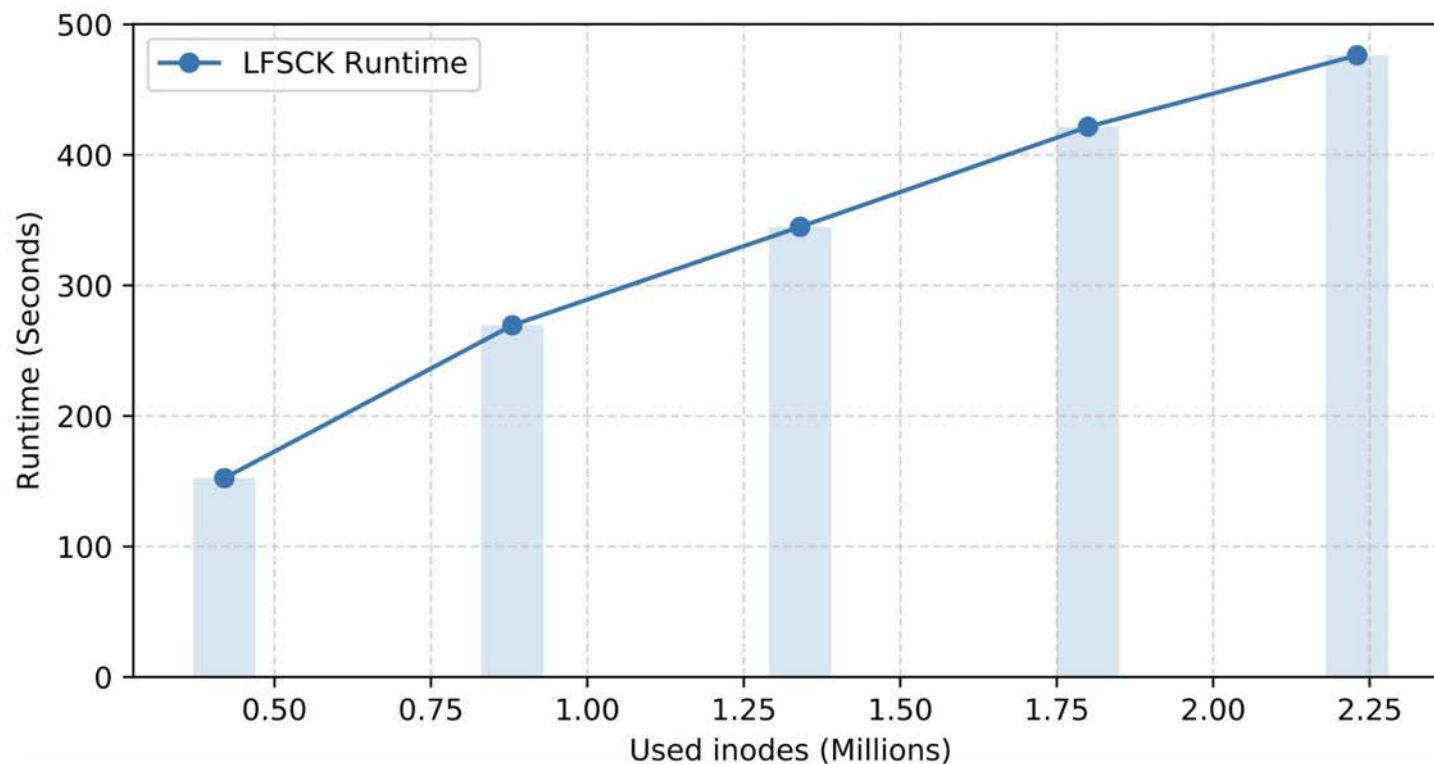
- Monitoring Framework

- The key is to monitor the execution of LFSCK
- We leverage two tools
- **iostat**
  - monitor the storage devices on all MDS and OSS nodes where Lustre was mounted.
  - `iostat -p sdc -d 1`
- **netstate**
  - monitor the utilization of the network cards in the cluster used by Lustre
  - `netstat --interfaces=enp6s0f0 -c 1`

# Results and Analysis

## ● Scalability

- How LFSCK reacts to increasing number of files (inodes) in the system?
- 8 OSS nodes and 1 MDS node
- inode increases from 0.4 million to 2.3 million



# Results and Analysis

- Scalability

- How LFSCK reacts to increasing number of OSS nodes (larger Lustre)?
- The same number of inodes (0.88 Million)
- LFSCK performance with 2, 4, and 8 OSS nodes
- With more OSS nodes, LFSCK takes longer time to finish
- Note that, this results is based on “stripe count = -1”, which stripes files to all the OSSes when possible
  - Production deployment may have better performance
  - For the future, this still can be a problem even for production deployment

TABLE I

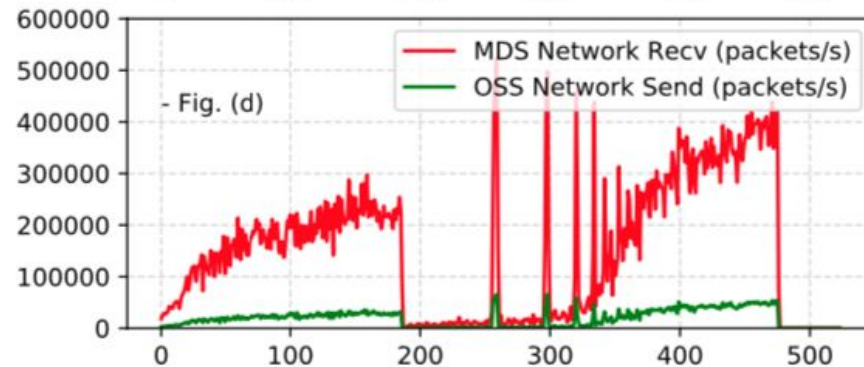
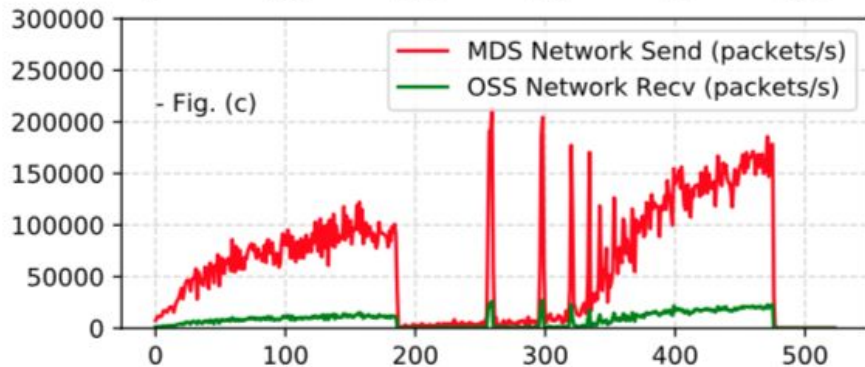
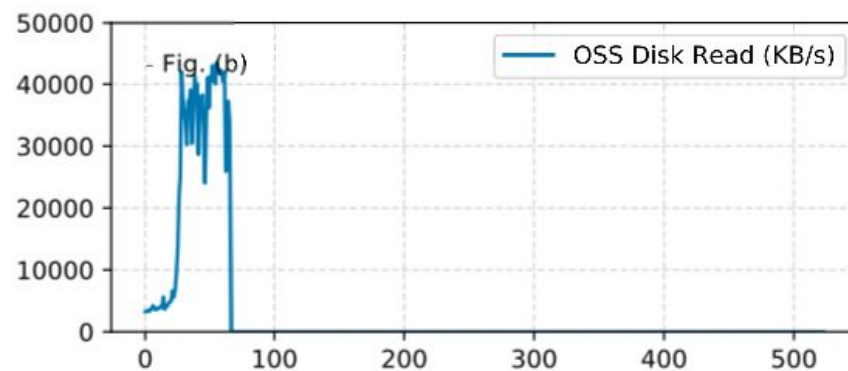
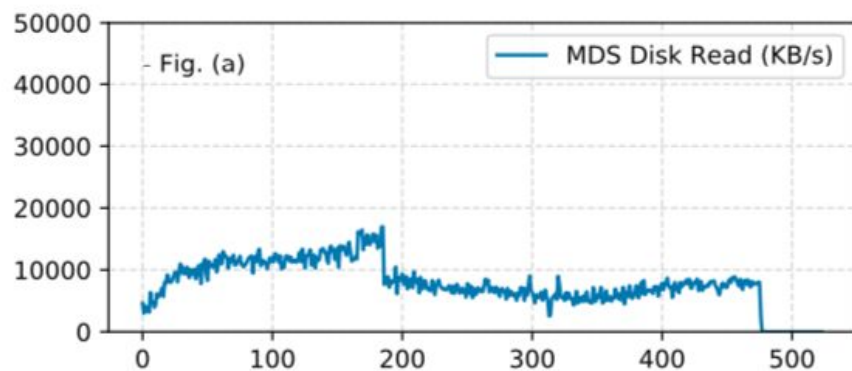
EXECUTION TIME OF LFSCK (IN SECONDS) ON LUSTRE WITH DIFFERENT NUMBER OF OSS NODES

# of OSS Nodes	2-OSS	4-OSS	8-OSS
Execution Time	144.8	170.5 (1.18X)	218.4 (1.50X)



# Results and Analysis

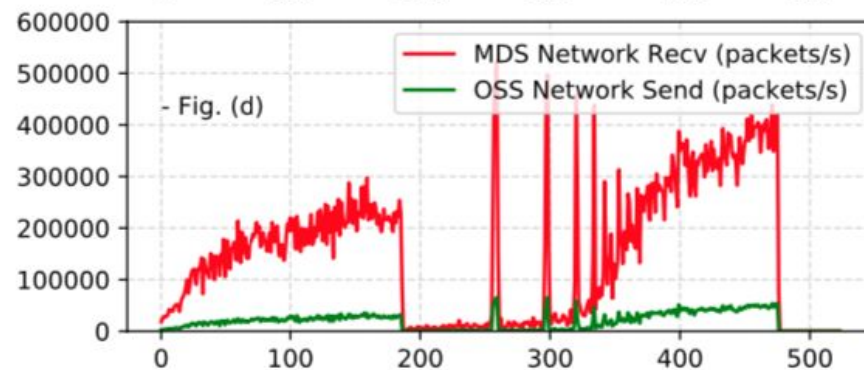
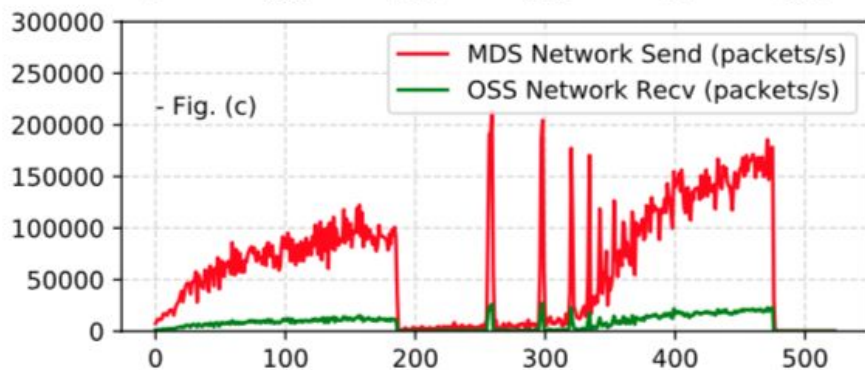
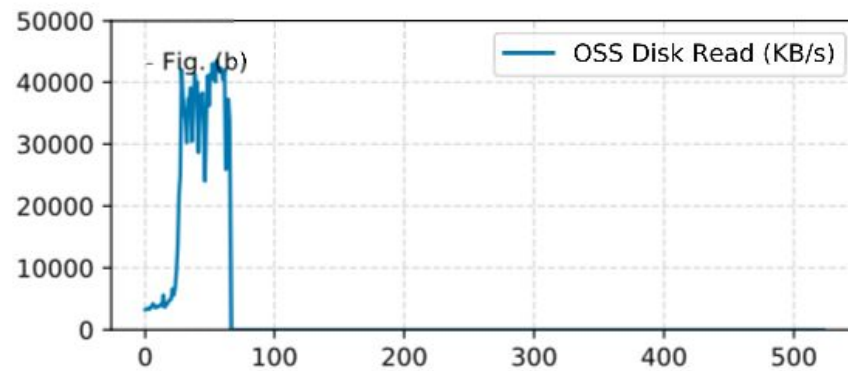
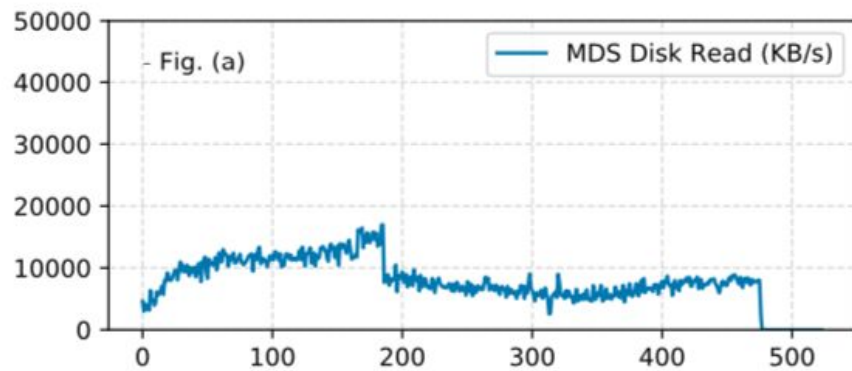
- LFSCK performance results
  - Is there potential performance bottleneck in LFSCK implementation?
  - 8 OSS nodes and 1 MDS node
  - Age Lustre with 2.5 million inodes and 4.8 TB of storage
  - Monitor Disk and Network to see what happened during LFSCK



# Results and Analysis

## ● Analysis

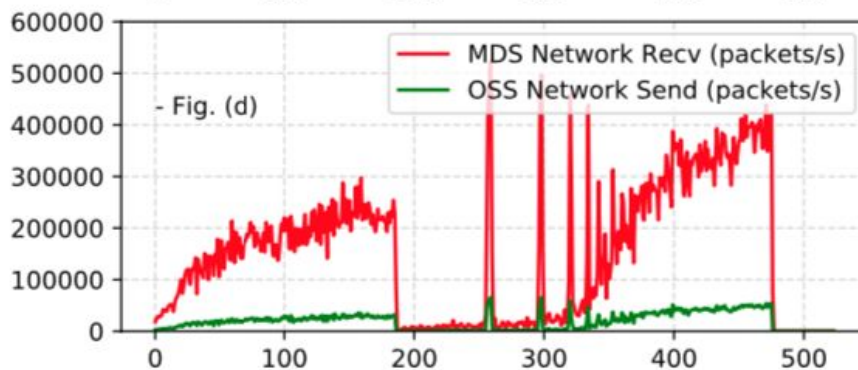
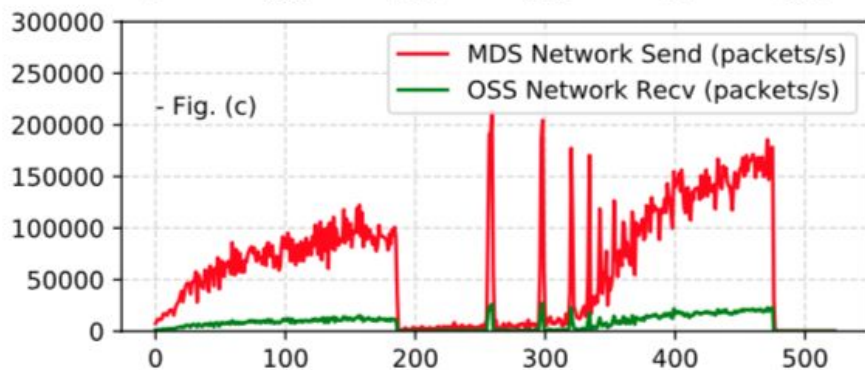
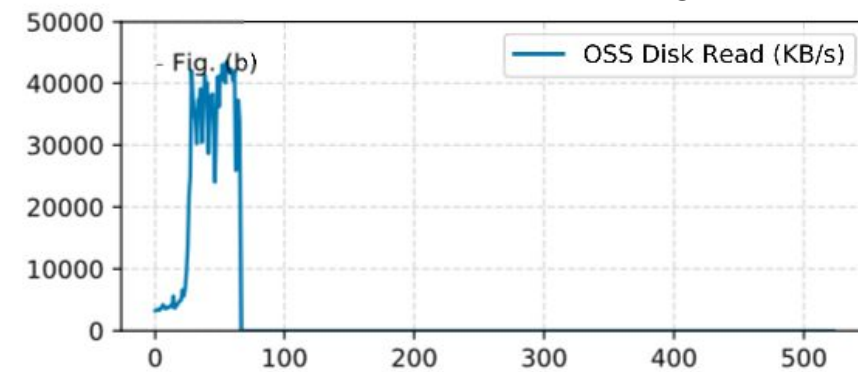
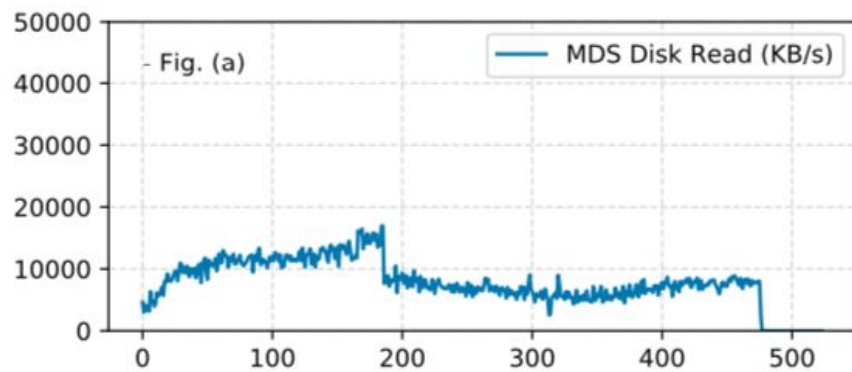
- In MDS and OSS nodes, network and disk are not fully utilized
  - Disk bandwidth is expected to be 100 MB/s
  - Network bandwidth should be more stable and much larger
- It might be because layout checking or namespace checking



# Results and Analysis

## ● Analysis

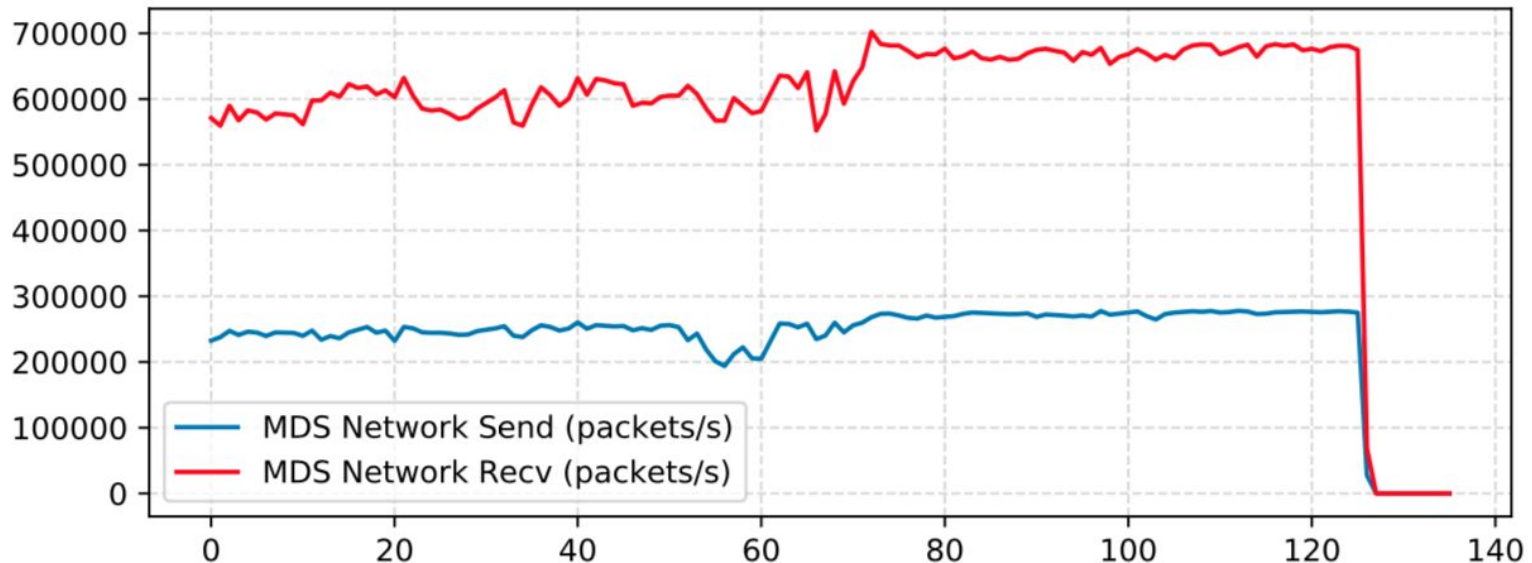
- Look at the network bandwidth closer
  - MDS send around eight times more packages than one OSS receives
  - Plus, MDS receives as many as twice packages as it sends out
- From the network point of view
  - Such a fan-out ratio indicates MDS can be saturated earlier on the receiving side



# Results and Analysis

- Layout Checking Performance

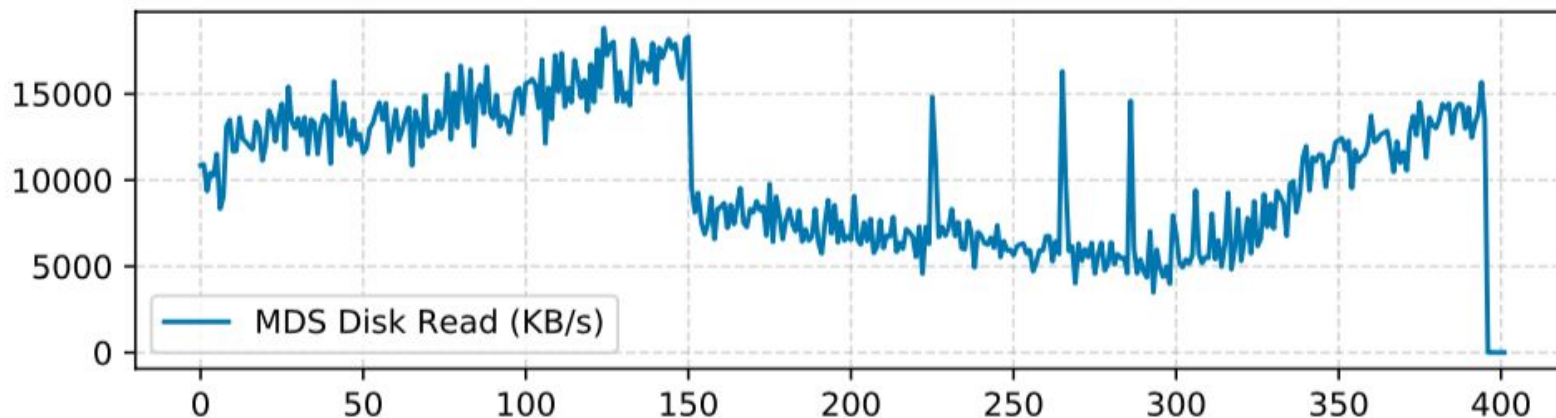
- Idea: all MDS and OSS nodes already buffered metadata in memory
- Then only network between MDS and OSS affects the performance
- LFSCK took around 130 seconds to finish (comparing to 500 seconds)
- Much higher and more stable network band comparing to NOT cached
- MDS receives double packages compared to what it sends
- Layout checking is not the bottleneck in previous example





# Results and Analysis

- Namespace Checking Performance
  - Namespace checking only happens in MDS node, but it can slow down the layout checking. It seems to be the bottleneck of previous test.
  - Verify this by only running namespace checking
  - Only running namespace checking takes 400 second
  - Incurs the similar low disk bandwidth
  - Confirm that namespace checking is the bottleneck.

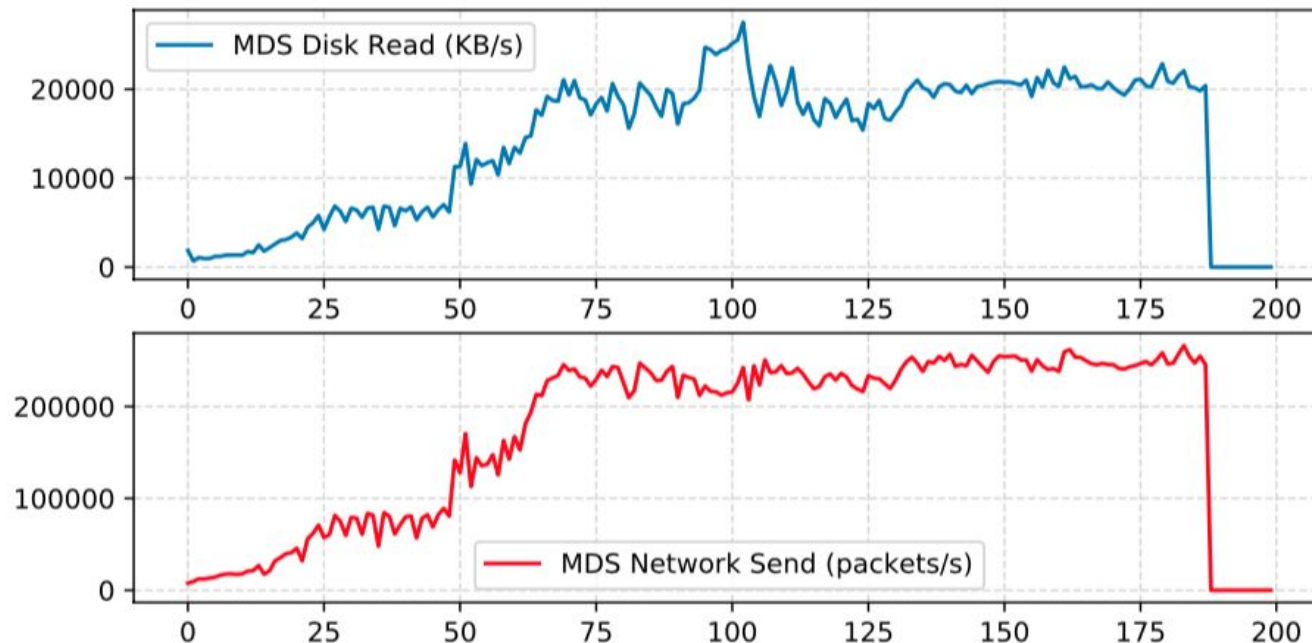


# Potentials

- Previous evaluation shows
  - The tight binding between layout and namespace checking in LFSCK design could generate performance bottleneck
  - Namespace checking is too slow due to the slow disk speed.
    - SSD might be helpful
  - Layout checking is faster
    - But, more OSSes can saturate the network earlier and becomes the bottleneck
  - Plus, things are actually dynamic
    - Namespace checking could be temporarily slow because of a huge directory
    - Layout checking could be temporarily slow because of a congested network
- Decouple the tight binding
  - Changing `SCRUB_WINDOW_SIZE` dynamically
  - Changing `LFSCK_ASYNC_WIN_DEFAULT` dynamically

# Potentials

- Quick demonstration using current implementation
  - Assume layout checking buffers all metadata in memory, indicating an infinite `SCRUB_WINDOW_SIZE` and `LFSCK_ASYNC_WIN_DEFAULT`
  - In this case, LFSCK finishes in 200 seconds
  - Better performance comes from better disk and network bands



# Conclusions and Plans

- Conclusions

- We show that LFSCK still has large room for performance improvement
- The root cause would be the tight binds of various components

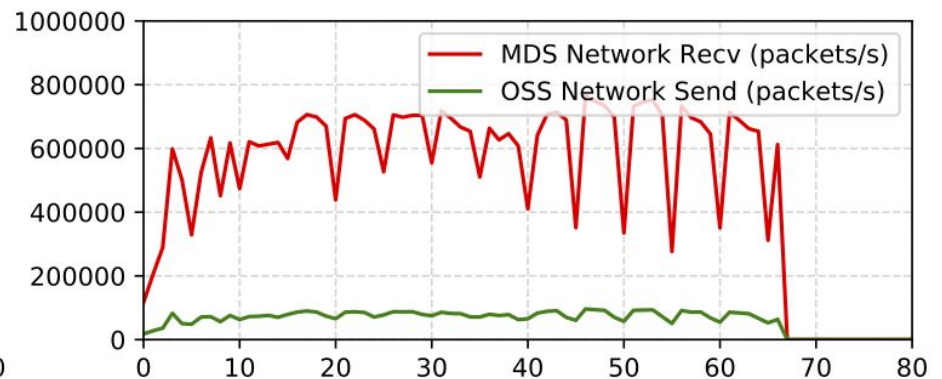
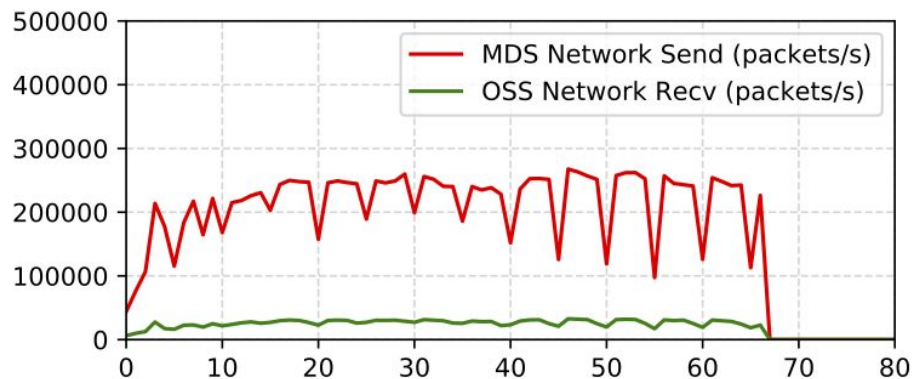
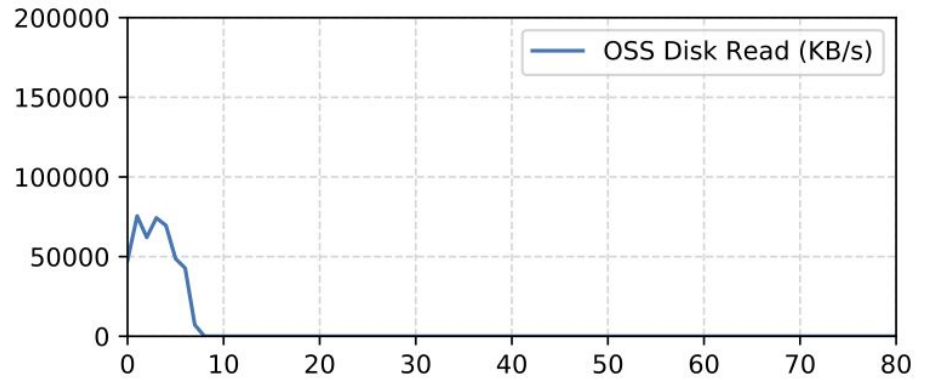
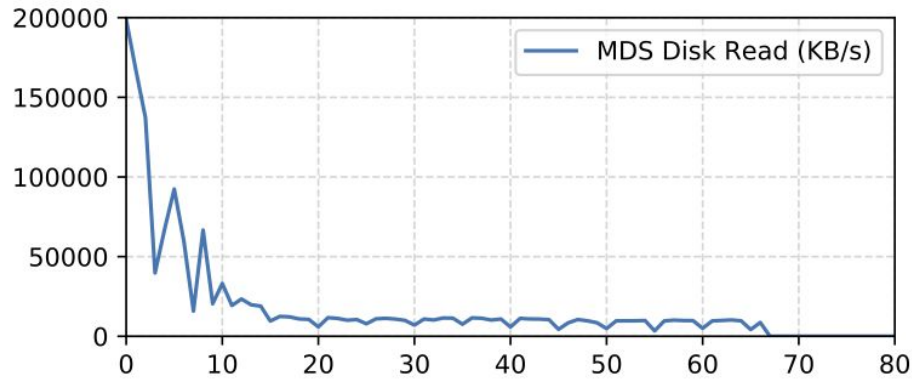
- Current Work and Plans

- We are running the experiments in more realistic settings
  - Using SSD instead of HDD
  - With more realistic stripe count setting
  - With more realistic aging strategy
- We are also implementing the proposed optimizations into LFSCK
  - Dynamic parameters setting in runtime
  - Compression of data packages during LFSCK
- We plan to investigate different strategies to implement LFSCK

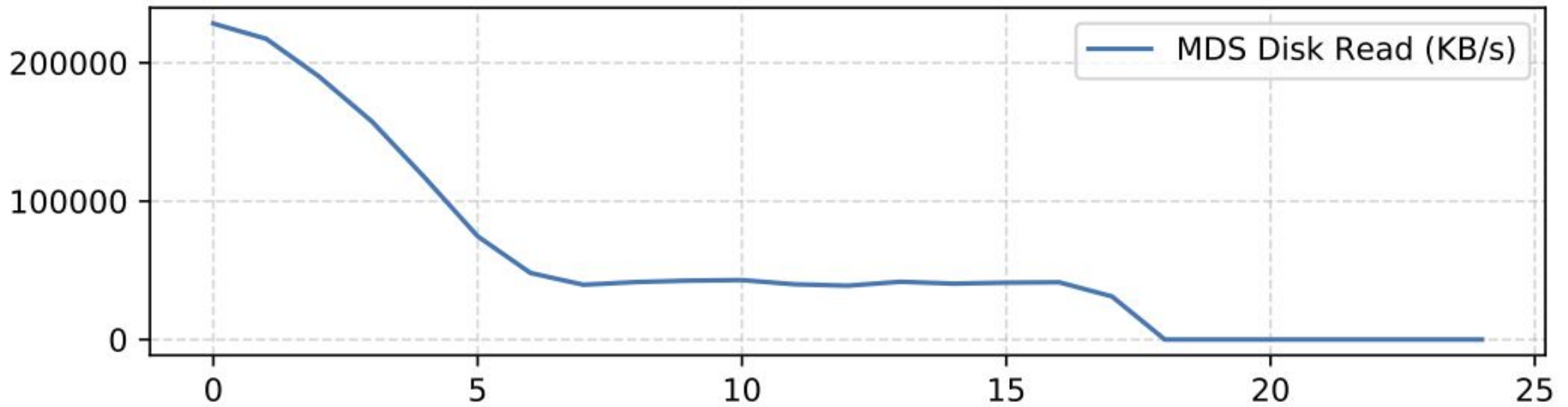


# Questions?

# LFSCK on SSD



## Namespace Checking on SSD



## Layout Checking on SSD

